CSC 261/461 – Database Systems Lecture 2

Spring 2018

Agenda

- Database System Concepts and Architecture
- 2. SQL introduction & schema definitions
 - ACTIVITY: Table creation
- 3. Basic single-table queries
 - ACTIVITY: Single-table queries!
- 4. Multi-table queries
 - ACTIVITY: Multi-table queries!

Table Schemas

• The schema of a table is the table name, its attributes, and their types:

```
Product(Pname: string, Price: float, Category: string, Manufacturer: string)
```

• A key is an attribute whose values are unique; we underline a key

```
Product(<u>Pname</u>: string, Price: float, Category: string, <u>Manufacturer</u>: string)
```

Database Schema vs. Database State

• Database State:

- Refers to the *content* of a database at a moment in time.

• Valid State:

- A state that satisfies the structure and constraints of the database.

Database Schema vs. Database State (continued)

- Distinction
 - The *database schema* changes very infrequently.
 - The *database state* changes every time the database is updated.
- Schema is also called intension.
- State is also called extension.

Example of a Database Schema

STUDENT

Name Student_number Class Major

Figure 2.1

Schema diagram for the database in Figure 1.2.

COURSE

| Course name | Course_number | Credit hours | Department |
|-------------|---------------|--------------|------------|
| | | | |

PREREQUISITE

| Course_number | Prerequisite_number |
|---------------|---------------------|
|---------------|---------------------|

SECTION

| Section_identifier | Course_number | Semester | Year | Instructor | |
|--------------------|---------------|----------|------|------------|--|
|--------------------|---------------|----------|------|------------|--|

GRADE_REPORT

Student_number | Section_identifier | Grade

Example of a database state

COURSE

| Course_name | Course_number | Credit_hours | Department |
|---------------------------|---------------|--------------|------------|
| Intro to Computer Science | CS1310 | 4 | CS |
| Data Structures | CS3320 | 4 | CS |
| Discrete Mathematics | MATH2410 | 3 | MATH |
| Database | CS3380 | 3 | CS |

SECTION

| Section_identifier | Course_number | Semester | Year | Instructor |
|--------------------|---------------|----------|------|------------|
| 85 | MATH2410 | Fall | 04 | King |
| 92 | CS1310 | Fall | 04 | Anderson |
| 102 | CS3320 | Spring | 05 | Knuth |
| 112 | MATH2410 | Fall | 05 | Chang |
| 119 | CS1310 | Fall | 05 | Anderson |
| 135 | CS3380 | Fall | 05 | Stone |

GRADE_REPORT

| Student_number | Section_identifier | Grade |
|----------------|--------------------|-------|
| 17 | 112 | В |
| 17 | 119 | С |
| 8 | 85 | Α |
| 8 | 92 | Α |
| 8 | 102 | В |
| 8 | 135 | A |

PREREQUISITE

Figure 1.2
A database that stores student and course information.

Course_number Prerequisite_number

CS3380 CS3320

MATH2410

CS3320 CS1310

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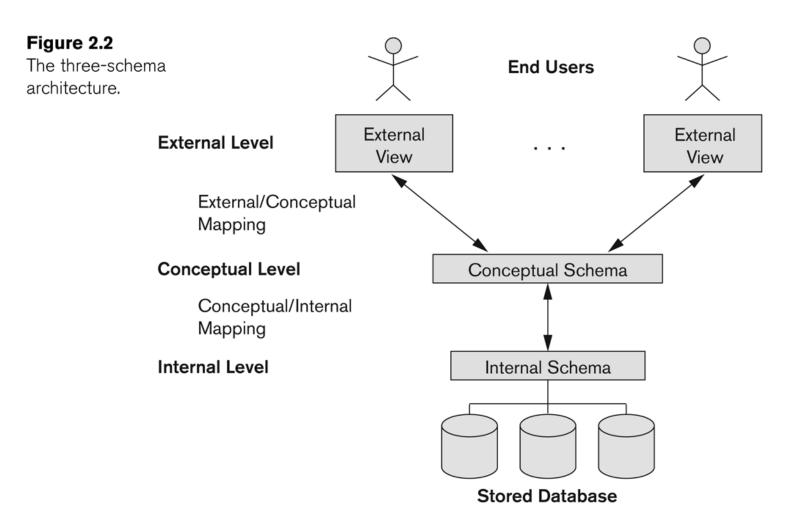
Three-Schema Architecture

- Proposed to support DBMS characteristics of:
 - Program-data independence.
 - Support of multiple views of the data.
- Not explicitly used in commercial DBMS products, but has been useful in explaining database system organization

Three-Schema Architecture

- Defines DBMS schemas at *three* levels:
 - Internal schema at the internal level to describe physical storage structures and access paths (e.g indexes).
 - Typically uses a physical data model.
 - Conceptual schema at the conceptual level to describe the structure and constraints for the whole database for a community of users.
 - Uses a conceptual or an implementation data model.
 - External schemas at the external level to describe the various user views.
 - Usually uses the same data model as the conceptual schema.

The three-schema architecture



Three-Schema Architecture

- Mappings among schema levels are needed to transform requests and data.
 - Programs refer to an external schema, and are mapped by the DBMS to the internal schema for execution.
 - Data extracted from the internal DBMS level is reformatted to match the user's external view (e.g. formatting the results of an SQL query for display in a Web page)

Data Independence

• Logical Data Independence:

- The capacity to change the conceptual schema without having to change the external schemas and their associated application programs.

• Physical Data Independence:

- The capacity to change the internal schema without having to change the conceptual schema.
- For example, the internal schema may be changed when certain file structures are reorganized or new indexes are created to improve database performance

Data Independence (continued)

- When a schema at a lower level is changed, only the mappings between this schema and higher-level schemas need to be changed in a DBMS that fully supports data independence.
- The higher-level schemas themselves are unchanged.
 - Hence, the application programs need not be changed since they refer to the external schemas.

DBMS Languages

- Data Definition Language (DDL)
- Data Manipulation Language (DML)
 - High-Level or Non-procedural Languages: These include the relational language SQL
 - May be used in a standalone way or may be embedded in a programming language
 - Low Level or Procedural Languages:
 - These must be embedded in a programming language

DBMS Languages

• Data Definition Language (DDL):

- Used by the DBA and database designers to specify the conceptual schema of a database.
- In many DBMSs, the DDL is also used to define internal and external schemas (views).

DBMS Languages

- Data Manipulation Language (DML):
 - Used to specify database retrievals and updates
 - DML commands (data sublanguage) can be *embedded* in a general-purpose programming language (host language), such as COBOL, C, C++, or Java.
 - A library of functions can also be provided to access the DBMS from a programming language
 - Alternatively, stand-alone DML commands can be applied directly (called a *query language*).

Types of DML

• High Level or Non-procedural Language:

- For example, the SQL relational language
- Are "set"-oriented and specify what data to retrieve rather than how to retrieve it.
- Also called **declarative** languages.

• Low Level or Procedural Language:

- Retrieve data one record-at-a-time;
- Constructs such as looping are needed to retrieve multiple records, along with positioning pointers.

DBMS Interfaces

- Stand-alone query language interfaces
 - Example: Entering SQL queries at the DBMS interactive SQL interface (e.g. SQL*Plus in ORACLE)
- Programmer interfaces for embedding DML in programming languages
- User-friendly interfaces
 - Menu-based, forms-based, graphics-based, etc.
- Mobile Interfaces
 - interfaces allowing users to perform transactions using mobile apps

User-Friendly DBMS Interfaces

- Menu-based (Web-based)
 - popular for browsing on the web
- Forms-based
 - designed for naïve users used to filling in entries on a form
- Graphics-based
 - Point and Click, Drag and Drop, etc.
 - Specifying a query on a schema diagram
- Natural language
 - requests in written English
- Combinations of the above:
 - For example, both menus and forms used extensively in Web database interfaces

Other DBMS Interfaces

- Natural language: free text as a query
- Speech : Input query and Output response
- Web Browser with keyword search
- Parametric interfaces, e.g., bank tellers using function keys.
- Interfaces for the DBA:
 - Creating user accounts, granting authorizations
 - Setting system parameters
 - Changing schemas or access paths

Typical DBMS Component Modules

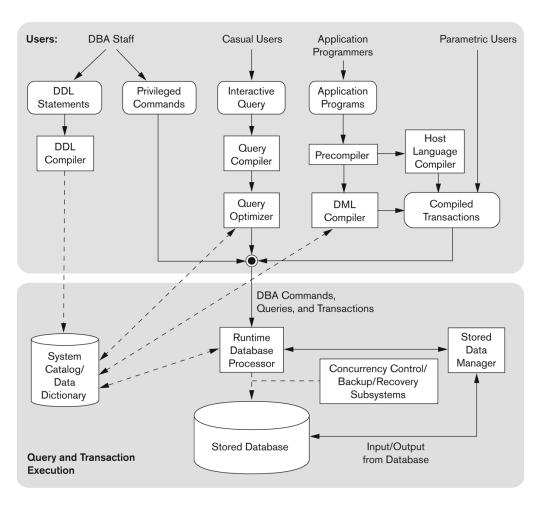


Figure 2.3

Component modules of a DBMS and their interactions.

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1. SQL INTRODUCTION & DEFINITIONS

What you will learn about in this section

- 1. What is SQL?
- 2. Basic schema definitions
- 3. Keys & constraints intro
- 4. ACTIVITY: CREATE TABLE statements

Basic SQL

SQL Introduction

- SQL is a standard language for querying and manipulating data
- SQL stands for Structured Query Language
- SQL is a very high-level programming language
 - This works because it is optimized well!

Basic SQL

- SQL language
 - Considered one of the major reasons for the commercial success of relational databases
- SQL
 - The origin of SQL is relational predicate calculus called tuple calculus which was proposed initially as the language SQUARE.
 - SQL Actually comes from the word "SEQUEL (Structured English Query Language)"
 - Original term used in the paper: "SEQUELTO SQUARE" by Chamberlin and Boyce. IBM could not copyright that term, so they abbreviated to SQL and copyrighted the term SQL.
 - Now popularly known as "Structured Query language".

SQL Data Definition, Data Types, Standards

- Terminology:
 - Table, row, and column used for relational model terms relation, tuple, and attribute
- CREATE statement
 - Main SQL command for data definition

SQL is a...

- Data Definition Language (DDL)
 - Define relational *schemata*
 - Create/alter/delete tables and their attributes
- Data Manipulation Language (DML)
 - Insert/delete/modify tuples in tables
 - Query one or more tables discussed next!

Product

| PName | Price | Manufacturer |
|-------------|----------|--------------|
| Gizmo | \$19.99 | GizmoWorks |
| Powergizmo | \$29.99 | GizmoWorks |
| SingleTouch | \$149.99 | Canon |
| MultiTouch | \$203.99 | Hitachi |

A <u>relation</u> or <u>table</u> is a multiset of tuples having the attributes specified by the schema

Let's break this definition down

Product

| PName | Price | Manufacturer |
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| Powergizmo | \$29.99 | GizmoWorks |
| SingleTouch | \$149.99 | Canon |
| MultiTouch | \$203.99 | Hitachi |

A <u>multiset</u> is an unordered list (or: a set with multiple duplicate instances allowed)

List: [1, 1, 2, 3] Set: {1, 2, 3} Multiset: {1, 1, 2, 3}

i.e. no *next()*, etc. methods!

Product

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|-------------|----------|--------------|
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An <u>attribute</u> (or <u>column</u>) is a typed data entry present in each tuple in the relation

NB: Attributes must have an <u>atomic</u> type in standard SQL, i.e. not a list, set, etc.

Product

| PName | Price | Manufacturer |
|-------------|----------|--------------|
| Gizmo | \$19.99 | GizmoWorks |
| Powergizmo | \$29.99 | GizmoWorks |
| SingleTouch | \$149.99 | Canon |
| MultiTouch | \$203.99 | Hitachi |

Also referred to sometimes as a <u>record</u>

A <u>tuple</u> or <u>row</u> is a single entry in the table having the attributes specified by the schema

Product

| PName | Price | Manufacturer |
|-------------|----------|--------------|
| Gizmo | \$19.99 | GizmoWorks |
| Powergizmo | \$29.99 | GizmoWorks |
| SingleTouch | \$149.99 | Canon |
| MultiTouch | \$203.99 | Hitachi |

The number of tuples is the **cardinality** of the relation

The number of attributes is the **arity** of the relation

Data Types in SQL

- Atomic types:
 - Characters: CHAR(20), VARCHAR(50)
 - Numbers: INT, BIGINT, SMALLINT, FLOAT
 - Others: MONEY, DATETIME, ...

- Every attribute must have an atomic type
 - Hence tables are flat

Key constraints

A **key** is a **minimal subset of attributes** that acts as a unique identifier for tuples in a relation

- A key is an implicit constraint on which tuples can be in the relation
 - i.e. if two tuples agree on the values of the key, then they must be the same tuple!

```
Students(sid:string, name:string, gpa:
float)
```

- 1. Which would you select as a key?
- 2. Can we have more than one key?

NULL and NOT NULL

- To say "don't know the value" we use NULL
 - NULL has (sometimes painful) semantics, more detail later

Students(sid:string, name:string, gpa:
float)

| sid | name | gpa |
|-----|------|------|
| 123 | Bob | 3.9 |
| 143 | Jim | NULL |

Say, Jim just enrolled in his first class.

In SQL, we may constrain a column to be NOT NULL, e.g., "name" in this table

General Constraints

- We can actually specify arbitrary assertions
 - E.g. "There cannot be 25 people in the DB class"
- In practice, we don't specify many such constraints. Why?
 - -<u>Performance!</u>

Whenever we do something ugly (or avoid doing something convenient) it's for the sake of performance

Summary of Schema Information

- Schema and Constraints are how databases understand the semantics (meaning) of data
- They are also useful for optimization
- SQL supports general constraints:
 - Keys and foreign keys are most important
 - We'll give you a chance to write the others

ACTIVITY

Acknowledgement

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