CSC 261/461 – Database Systems Lecture 3

Spring 2018

Announcements

- 1. Start forming project teams.
 - Maximum Size of a team: 4
 - It's a team project so you are not allowed to work alone!
- 2. Project 1 (milestone 1) will be released soon (this weekend).
- 3. We will post a problem set this week.
 - Useful for Quiz 2
- 4. (Optional) workshops starting from Tomorrow

Today's Lecture

- 1. Single Table Query
- 2. Multi Table Query
- 3. Aggregation & GROUP BY
- 4. Advanced SQL

SINGLE-TABLE QUERIES

What you will learn about in this section

- 1. The SFW query
- 2. Other useful operators: LIKE, DISTINCT, ORDER BY
- 3. ACTIVITY: Single-table queries

SQL Query

Basic form (there are many many more bells and whistles)

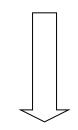
Call this a **SFW** query.

Simple SQL Query: Selection

Selection is the operation of filtering a relation's tuples on some condition

| PName | Price | Category | Manufacturer |
|-------------|----------|-------------|--------------|
| Gizmo | \$19.99 | Gadgets | GizmoWorks |
| Powergizmo | \$29.99 | Gadgets | GizmoWorks |
| SingleTouch | \$149.99 | Photography | Canon |
| MultiTouch | \$203.99 | Household | Hitachi |

```
SELECT *
FROM Product
WHERE Category = 'Gadgets'
```



| PName | Price | Category | Manufacturer |
|------------|---------|----------|--------------|
| Gizmo | \$19.99 | Gadgets | GizmoWorks |
| Powergizmo | \$29.99 | Gadgets | GizmoWorks |

Simple SQL Query: Projection

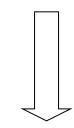
Projection is the operation of producing an output table with tuples that have a subset of their prior attributes

| PName | Price | Category | Manufacturer |
|-------------|----------|-------------|--------------|
| Gizmo | \$19.99 | Gadgets | GizmoWorks |
| Powergizmo | \$29.99 | Gadgets | GizmoWorks |
| SingleTouch | \$149.99 | Photography | Canon |
| MultiTouch | \$203.99 | Household | Hitachi |

SELECT Pname, Price, Manufacturer

FROM Product

WHERE Category = 'Gadgets'



| PName | Price | Manufacturer |
|------------|---------|--------------|
| Gizmo | \$19.99 | GizmoWorks |
| Powergizmo | \$29.99 | GizmoWorks |

Notation

Input schema

Product(<u>PName</u>, Price, Category, <u>Manfacturer</u>)

SELECT Pname, Price, Manufacturer

FROM Product

WHERE Category = 'Gadgets'

Output schema

Answer(PName, Price, Manfacturer)

A Few Details

- SQL commands are case insensitive:
 - Same: SELECT, Select, select
- Table or relation names may be case insensitive:
 - Depends on OS. For Unix, it's case sensitive.
 - Different: Product, product

- Values are not:
 - <u>Different:</u> 'Seattle', 'seattle'
- Use single quotes for constants:
 - 'abc' yes
 - "abc" no

LIKE: Simple String Pattern Matching

```
SELECT *
FROM Products
WHERE PName LIKE '%gizmo%'
```

- s LIKE p: pattern matching on strings
- p may contain two special symbols:
 - % = any sequence of characters
 - _ = any single character

DISTINCT: Eliminating Duplicates

SELECT DISTINCT Category FROM Product



Category

Gadgets

Photography

Household

Versus

SELECT Category FROM Product



Category

Gadgets

Gadgets

Photography

Household

ORDER BY: Sorting the Results

SELECT PName, Price, Manufacturer

FROM Product

WHERE Category='gizmo' AND Price > 50

ORDER BY Price, PName

Ties are broken by the second attribute on the ORDER BY list, etc.

Ordering is ascending, unless you specify the DESC keyword.

3. MULTI-TABLE QUERIES

What you will learn about in this section

1. Foreign key constraints

2. Joins: basics

3. Joins: SQL semantics

4. ACTIVITY: Multi-table queries

Foreign Key constraints

• Suppose we have the following schema:

```
Students(sid: string, name: string, gpa: float)
Enrolled(student_id: string, cid: string, grade: string)
```

- And we want to impose the following constraint:
 - 'Only bona fide students may enroll in courses' i.e. a student must appear in the Students table to enroll in a class

| Stude | nts | | Enrolled | | | |
|-------|------|-----|---------------|------------|-----|-------|
| sid | name | gpa | | student_id | cid | grade |
| 101 | Bob | 3.2 | | 123 | 564 | Α |
| 123 | Mary | 3.8 | \rightarrow | 123 | 537 | A+ |

student_id alone is not a key- what is?

We say that student_id is a **foreign key** that refers to Students

Declaring Foreign Keys

Foreign Keys and update operations

```
Students(<u>sid</u>: string, name: string, gpa: float)
Enrolled(<u>student_id</u>: string, <u>cid</u>: string, grade: string)
```

- What if we insert a tuple into Enrolled, but no corresponding student?
 - INSERT is rejected (foreign keys are <u>constraints</u>)!

Keys and Foreign Keys

Company

| <u>CName</u> | StockPrice | Country |
|----------------|------------|---------|
| GizmoWork s | 25 | USA |
| Canon | 65 | Japan |
| Hitachi | 15 | Japan |

Company:

Primary Key: CName

Foreign Key: None

Product:

Primary Key: PName

Foreign Key: Manufacturer references CName

Product

| <u>PName</u> | Price | Category | Manufacturer |
|--------------|----------|-------------|--------------|
| Gizmo | \$19.99 | Gadgets | GizmoWorks |
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What is a foreign key vs. a primary key here?

```
Product(<u>PName</u>, Price, Category, Manufacturer)
Company(<u>CName</u>, StockPrice, Country)
```

Ex: Find all products under \$200 manufactured in Japan; return their names and prices.

Note: we will often omit attribute types in schema definitions for brevity, but assume attributes are always atomic types

```
Product(<u>PName</u>, Price, Category, Manufacturer)
Company(<u>CName</u>, StockPrice, Country)
```

Ex: Find all products under \$200 manufactured in Japan; return their names and prices.

```
SELECT PName, Price
FROM Product, Company
WHERE Manufacturer = CName
AND Country='Japan'
AND Price <= 200
```

A join between tables returns all unique combinations of their tuples which meet some specified join condition

```
Product(<u>PName</u>, Price, Category, Manufacturer)
Company(<u>CName</u>, StockPrice, Country)
```

Several equivalent ways to write a basic join in SQL:

A few more later on...

Product

| PName | Price | Category | Manuf |
|-------------|-------|-------------|---------|
| Gizmo | \$19 | Gadgets | GWorks |
| Powergizmo | \$29 | Gadgets | GWorks |
| SingleTouch | \$149 | Photography | Canon |
| MultiTouch | \$203 | Household | Hitachi |

| | | Company |
|---------|-------|---------|
| Cname | Stock | Country |
| GWorks | 25 | USA |
| Canon | 65 | Japan |
| Hitachi | 15 | Japan |
| • | | |

Company



| PName | Price | |
|-------------|----------|--|
| SingleTouch | \$149.99 | |

Tuple Variable Ambiguity in Multi-Table

```
Person(<u>name</u>, address, worksfor)
Company(<u>name</u>, address)
```

```
SELECT DISTINCT name, address
FROM Person, Company
WHERE worksfor = name
```

Which "address" does this refer to?

Which "name"s??

Tuple Variable Ambiguity in Multi-Table

```
Person(<u>name</u>, address, worksfor)
Company(<u>name</u>, address)
```

Both equivalent ways to resolve variable ambiguity

```
SELECT DISTINCT Person.name, Person.address
FROM Person, Company
WHERE Person.worksfor = Company.name
```

```
SELECT DISTINCT p.name, p.address
FROM Person p, Company c
WHERE p.worksfor = c.name
```

Acknowledgement

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- Many of these slides are taken from cs145 course offered by Stanford University.
- Thanks to YouTube, especially to <u>Dr. Daniel Soper</u> for his useful videos.