# CSC2/458 Parallel and Distributed Systems A Case Study of a High-speed GPU NFA Implementation

Sreepathi Pai April 20, 2022

URCS

Irregular Parallelism

Reducing Data Movement

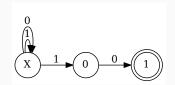
Improving Compute Utilization

Irregular Parallelism

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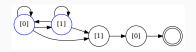
Improving Compute Utilization

- Recall CSC173
- NFA are machines that recognize regular languages
- Are "non-deterministic"
  - can be in multiple states at the same time
- Easy to parallelize
  - unlike DFA



# Homogenous Non-Deterministic Finite Automata

- Traditionally, NFA match characters on their edges
- *Glushkov* NFAs or homogeneous NFAs match on the nodes instead
  - Also don't have  $\epsilon$  transitions
- All children of a node are "activated" on a match
- Only activated nodes try to match a character
- Traditional NFAs can be converted to Homogeneous NFAs losslessly



- Invented by Micron, the memory company
- Never made into a product, AFAIK
- Execution engine for Homogeneous NFAs
  - A non-von Neumann processor
- Ran at 133MHz, about 7.5ns per input symbol
  - Roughly 133MB/s

Regular expression: .\*cat turns into NFA:

and is mapped to the AP as:

Property	AP	GPU
Parallelism	49,152	163,840
Bandwidth	133MB/s	768GB/s

This is going to be easy.

- Obi Wan Kenobi, Revenge of the Sith

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Property	Regular	Irregular
Task Decomposition	Early	Late
Task Placement	Static	Dynamic
Task Activity	Static	Dynamic
Data Access Patterns	Fixed	Input-dependent
Communication Patterns	Fixed	Input-dependent
Example	MM	Sparse MM

(MM - Matrix Multiply)

What shall we parallelize on?

- Each input symbol "must" be processed serially
- Multiple states can be active at the same time
- Pick states as tasks?

How do we map tasks to GPU hardware?

- There can be thousands of states
- Each thread gets one NFA state?
- But each input byte must be processed serially
  - Can't split up an NFA across thread blocks
  - Otherwise would require global synchronization

- Each task must access:
  - input symbol (1 byte)
  - match set (1 byte in 32 bytes, depending on input symbol)
- If symbol matches:
  - children of current node must be activated (1 bit/child state)

What issues can we see in this scheme?

- Many threads are idle most of the time
  - Not all NFA states are active on every input symbol
- No useful work done
- One could use a *worklist* instead to track active states
  - Dynamic assignment of tasks to threads

- 114 KB/second
  - Ouch!

[according to the paper]

Irregular Parallelism

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- NewTran (pron. "neutron")
- Match Set Compression

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- Fixed assignment for highly active states
  - Profiling
  - BFS "depth"
  - Hotstart
- Dynamic assignment for low activity states
  - Worklist

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- Did we handily beat the AP?
- What kind of performance modeling is needed here?

- NFAs are graph-based computations
  - Not instruction-based programs
- Do you know of other graph-based computations?