

Bits, Bytes, and Integers

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Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- Integers
 - Representation: unsigned and signed
 - Conversion, casting

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Byte-Oriented Memory Organization



- Programs Refer to Virtual Addresses
 - Conceptually very large array of bytes
 - Actually implemented with hierarchy of different memory types
 - System provides address space private to particular “process”
 - Program being executed
 - Program can clobber its own data, but not that of others
- Compiler + Run-Time System Control Allocation
 - Where different program objects should be stored
 - All allocation within single virtual address space

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Encoding Byte Values

- Byte = 8 bits
 - Binary 00000000_2 to 11111111_2
 - Decimal: 0_{10} to 255_{10}
 - Hexadecimal 00_{16} to FF_{16}
 - Base 16 number representation
 - Use characters ‘0’ to ‘9’ and ‘A’ to ‘F’
 - Write $FA1D37B_{16}$ in C as
 - $0xFA1D37B$
 - $0xfa1d37b$

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

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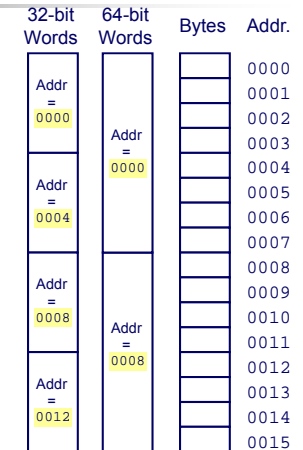
Machine Words

- Machine Has “Word Size”
 - Nominal size of integer-valued data, more importantly, basic unit of internal storage (registers) and computation
 - Including addresses
 - Some machines use 32 bits (4 bytes) addresses
 - Limits addresses to 4GB
 - Too small for memory-intensive applications
 - High-end systems use 64 bits (8 bytes) addresses
 - Potential address space $\approx 1.8 \times 10^{19}$ bytes
 - x86-64 machines support 48-bit addresses: 256 Terabytes
 - Machines support multiple data formats
 - Fractions or multiples of word size
 - Always one or multiples of bytes

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Word-Oriented Memory Organization

- Addresses Specify Byte Locations
 - Address of first byte in word
 - Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)



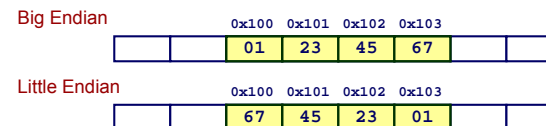
Data Representations

C Data Type	Typical 32-bit	Intel IA32	x86-64
char	1	1	1
short	2	2	2
int	4	4	4
long	4	4	8
long long	8	8	8
float	4	4	4
double	8	8	8
long double	8	10/12	10/16
pointer	4	4	8

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Byte Ordering

- How should bytes within a multi-byte word be ordered in memory?
- Big Endian (Sun, PPC Mac, Internet)
 - Most significant byte has lowest address
- Little Endian (Intel x86)
 - Least significant byte has lowest address
- Example
 - Variable x has 4-byte representation 0x01234567
 - Address given by &x is 0x100



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Reading Byte-Reversed Listings

- Disassembly
 - Text representation of binary machine code
 - Generated by program that reads the machine code
- Example Fragment

Address	Instruction Code	Assembly Rendition
8048365:	5b	pop %ebx
8048366:	81 c3 ab 12 00 00	add \$0x12ab,%ebx
804836c:	83 bb 28 00 00 00 00	cmpl \$0x0,0x28(%ebx)

- Deciphering Numbers

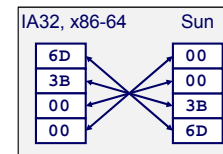
- Value: 0x12ab
- Pad to 32 bits: 0x000012ab
- Split into bytes: 00 00 12 ab
- Reverse: ab 12 00 00

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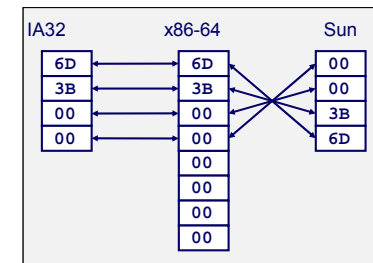
Representing Integers

Decimal: 15213
 Binary: 0011 1011 0110 1101
 Hex: 3 B 6 D

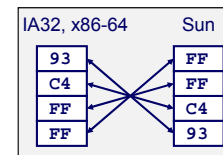
int A = 15213;



long int C = 15213;



int B = -15213;

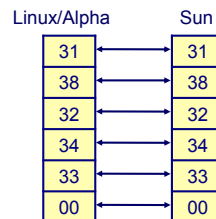
Two's complement representation
(Covered later)

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Representing Strings

- Strings in C
 - Represented by array of characters
 - Each character encoded in ASCII
 - Standard 7-bit encoding of characters
 - Character "0" has code 0x30
 - Digit i has code $0x30+i$
 - String should be null-terminated
 - Final character = 0
- Compatibility
 - Byte ordering not an issue

char S[6] = "18243";



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Boolean Algebra

- Algebraic representation of logic
 - Encode "True" as 1 and "False" as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

$\&$	0	1
0	0	0
1	0	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

$ $	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

\sim	0	1
0	1	0
1	0	1

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

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Boolean Algebra on Bit Vectors

- Operate on Bit Vectors
 - Operations applied bitwise

01101001	01101001	01101001	01101001
$\& 01010101$	$ 01010101$	$\wedge 01010101$	~ 01010101
01000001	01111101	00111100	10101010

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Representing & Manipulating Sets

Representation

- Width w bit vector represents subsets of $\{0, \dots, w-1\}$
- $a_j = 1$ if $j \in A$
 - 01101001 $\{0, 3, 5, 6\}$
 - 76543210
 - 01010101 $\{0, 2, 4, 6\}$
 - 76543210

Operations

- $\&$ Intersection 01000001 $\{0, 6\}$
- $|$ Union 01111101 $\{0, 2, 3, 4, 5, 6\}$
- \wedge Symmetric difference 00111100 $\{2, 3, 4, 5\}$
- \sim Complement 10101010 $\{1, 3, 5, 7\}$

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Bit-Level Operations in C

- Operations $\&$, $|$, \sim , \wedge Available in C
 - Apply to any basic data type
 - long, int, short, char, unsigned
 - View arguments as bit vectors
 - Arguments applied bit-wise
- Examples
 - $\sim 0x41 = 0xBE$
 - $\sim 01000001_2 = 10111110_2$
 - $\sim 0x00 = 0xFF$
 - $\sim 00000000_2 = 11111111_2$
 - $0x69 \& 0x55 = 0x41$
 - $01101001_2 \& 01010101_2 = 01000001_2$
 - $0x69 | 0x55 = 0x7D$
 - $01101001_2 | 01010101_2 = 01111101_2$

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Contrast: Logic Operations in C

- Contrast to Logical Operators
 - &&, ||, !
 - View 0 as "False"
 - Anything nonzero as "True"
 - Always return 0 or 1
 - Early termination
- Examples (char data type)
 - !0x41 = 0x00
 - !0x00 = 0x01
 - !!0x41 = 0x01
 - 0x69 && 0x55 = 0x01
 - 0x69 || 0x55 = 0x01
 - p && *p (avoids null pointer access)

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Shift Operations

- Left Shift: $x \ll y$
 - Shift bit-vector x left y positions
 - Throw away extra bits on left
 - Fill with 0's on right
- Right Shift: $x \gg y$
 - Shift bit-vector x right y positions
 - Throw away extra bits on right
 - Logical shift
 - Fill with 0's on left
 - Arithmetic shift
 - Replicate most significant bit on right

Argument x	01100010
<< 3	00010000
Log. >> 2	00011000
Arith. >> 2	00011000

Argument x	10100010
<< 3	00010000
Log. >> 2	00101000
Arith. >> 2	11101000

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Encoding Integers

Unsigned

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

Two's Complement

$$B2T(X) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

```
short int x = 15213;
short int y = -15213;
```

Sign Bit

- C short 2 bytes long

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
y	-15213	C4 93	11000100 10010011

- Sign Bit
 - For 2's complement, most significant bit indicates sign
 - 0 for nonnegative, 1 for negative
- What is the bit representation of -1?

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Encoding Example (Cont.)

x = 15213: 00111011 01101101
y = -15213: 11000100 10010011

Weight	15213	-15213
1	1	1
2	0	1
4	1	0
8	1	0
16	0	1
32	1	0
64	1	0
128	0	1
256	1	0
512	1	0
1024	0	1
2048	1	0
4096	1	0
8192	1	0
16384	0	1
-32768	0	1
Sum	15213	-15213

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Numeric Ranges

Unsigned Values

- $UMin = 0$
000...0
- $UMax = 2^w - 1$
111...1

Two's Complement Values

- $TMin = -2^{w-1}$
100...0
- $TMax = 2^{w-1} - 1$
011...1

Values for $W = 16$

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000

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Values for Different Word Sizes

	W			
	8	16	32	64
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615
TMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

Observations

- $|TMin| = TMax + 1$
 - Asymmetric range
- $UMax = 2 * TMax + 1$

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Unsigned & Signed Numeric Values

X	B2U(X)	B2T(X)
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

Equivalence

- Same encodings for nonnegative values

Uniqueness

- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding

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Disclaimer

These slides were adapted from the CMU course slides provided along with the textbook of "Computer Systems: A programmer's Perspective" by Bryant and O'Hallaron. The slides are intended for the sole purpose of teaching the computer organization course at the University of Rochester.

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