



Performance of HTTP Persistency

TCP per-connection overhead:

- Connection establishment
- Congestion control: slow start

Non-persistent HTTP (1.0)

- At most one object is sent over a TCP connection.
- Pays TCP per-connection overhead for each object.

Persistent HTTP (1.1)

 Multiple objects can be sent over single TCP connection between the browser and web server.

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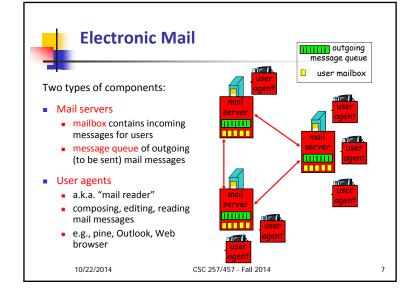
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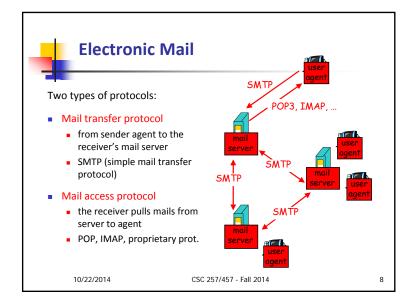
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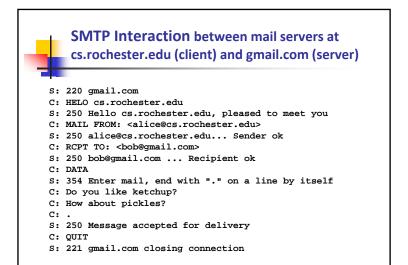
Additional Example Applications

- Remote console
 - · Components: host machine and console client
 - Comm. protocol: SSH
- Electronic mail
 - Components: mail servers and mail readers
 - Comm. protocol: SMTP, POP3, IMAP
- Domain name system (DNS)
- Network file system (NFS)
- Peer-to-peer content sharing (Gnutella, BitTorrent)
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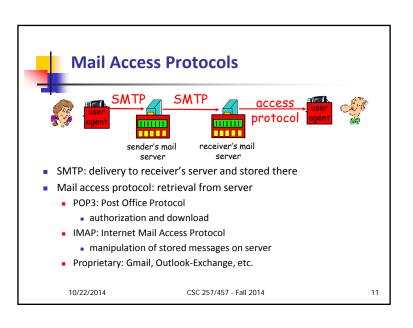


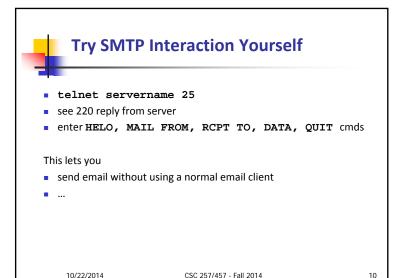




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DNS: Domain Name System

People: multiple identifiers

- SSN good for machine bookkeeping, indexing
- name human friendly, easy to remember

Internet hosts:

- IP address used for addressing, routing on the Internet
- "name", e.g., cycle1.cs.rochester.edu human friendly

Q: map between IP addresses and names?

- DNS query: find the IP address for a given name
- not part of Internet as a network
- ease of use for Internet applications
 - "ssh cycle1.cs.rochester.edu" vs. "ssh 128.151.67.99"
 - http://www.cnn.com vs. http://64.236.24.20

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Domain Name System: Overview

- Distributed database
 - Implemented with collaboration of many name servers distributed all
- DNS queries
 - Name lookup: find the IP address for a given name
 - Others: find mail server for a domain, ipv4/v6 address mapping, ...

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DNS: Decentralized and Hierarchical

- Decentralized
 - a large number of servers/sites work together
- Why not centralize DNS?
 - performance bottleneck
 - single point of failure
 - can only be close to few hosts ⇒ long latency to many others
- Hierarchical/partitioned
 - there are different types of name servers in charge of different mappings
- What if massively replicating it?
 - high overhead of maintaining consistency: hard to make any changes on the data

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Types of Name Servers

- Local name server (defined for a particular group of hosts):
 - each organization often has a local (default) name server
 - used when one of the hosts in its domain initiates a DNS query
 - first stop for a DNS query
- Authoritative name server (defined for a particular group of hosts):
 - maintain "IP address name" mapping for a group of hosts so it can answer DNS query for these hosts
 - used when one of the hosts in its domain is the target of a DNS query
 - last stop for a DNS query (if ever reached)
 - often the same as local name server for an organization
- Each local NS may not know all authoritative NSes
 - something is missing ...

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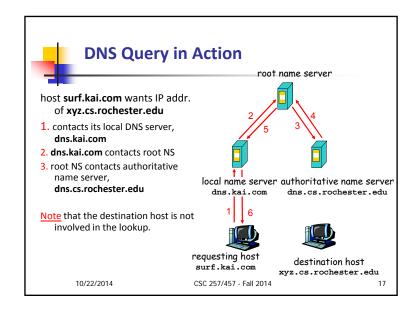
Root Name Servers

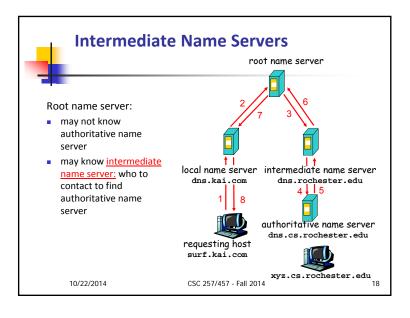
- Root name server (defined globally)
 - bridging between local NSs and authoritative NSes
 - contacted by local name server that cannot resolve name
- What it does:
 - · contacts authoritative name server
 - gets result and returns it to the requesting local name server

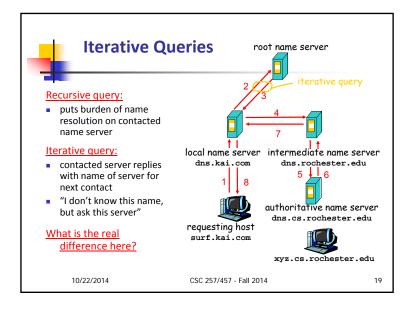


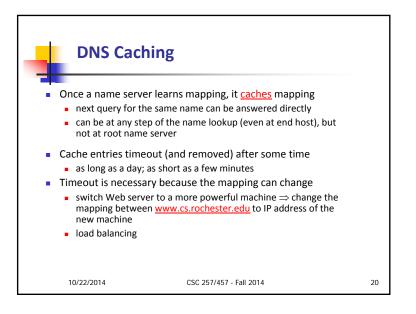
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DNS Properties

- Scalability
 - poor scalability at the few root name servers
 - made more scalable by iterative queries and caching
- Reliability
 - problems at (even a single) root name server(s) can cause big problem
 - again, caching helps improve reliability
- Consistency
 - DNS caching is replication without consistency guarantee
 - only feasible for stable name-IP address mapping, may not be suitable for more dynamic information

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TCP vs. UDP

- We saw applications:
 - Web/HTTP
 - Email/SMTP/POP3/IMAP
 - DNS
- They build on some transport services.
 - TCP or UDP?

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