

## Principles of Network Security

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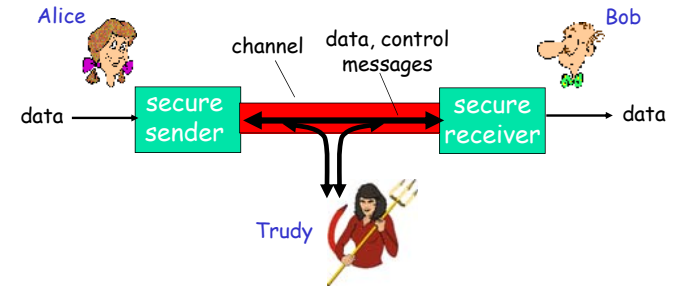
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## The Network Security Model

- Bob and Alice want to communicate “securely”.
- Trudy (the adversary) has access to the channel.



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## Who might Bob and Alice be?

- Web browser/server for electronic transactions (e.g., on-line purchases/banking)
- DNS servers
- Routers exchanging routing table updates
- ... well, *real-life* Bobs and Alices!

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## What can an adversary do?

- **Eavesdrop**: understand the content of messages
- Actively **changing** messages
- **Impersonation**: fake (spoof) identity
- **Denial of service**: prevent service from being used by others (e.g., by overloading resources)

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## What is Network Security?

- Confidentiality:** only sender, intended receiver should "understand" message contents.
- Authentication:** sender, receiver want to confirm identity of each other.
- Message Integrity:** sender, receiver want to ensure message not altered (in transit, or afterwards).
- Access and Availability:** services must be accessible and available to (and only to) legitimate users.

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## Principles of Network Security

- Confidentiality: cryptography
- Authentication
- Integrity
- Key distribution and certification

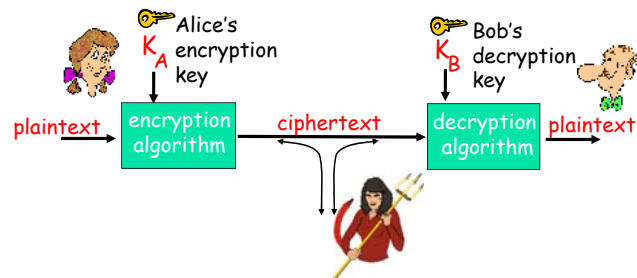
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## The Language of Cryptography

First goal of cryptography: confidentiality.



- **Symmetric key** crypto: encryption and decryption keys are identical. (both are **secret**)
- **Public key** crypto: encryption key is **public**, decryption key is **secret**.

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## Symmetric Key Cryptography: Monoalphabetic Cipher

**Monoalphabetic cipher:** substitute one letter for another.

plaintext:   abcdefghijklmnopqrstuvwxyz  
 ciphertext:   mnbvcxzasdfghjklpoiuytrewq

Example:   Plaintext: bob. i love you. alice  
               ciphertext: nkn. s gktc wky. mgsbc

Q1: How hard to break this simple cipher?

- ☐ brute force?
- ☐ other?

Q2: How to make it more difficult to break?

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## Symmetric Key Cryptography: DES

### DES: Data Encryption Standard

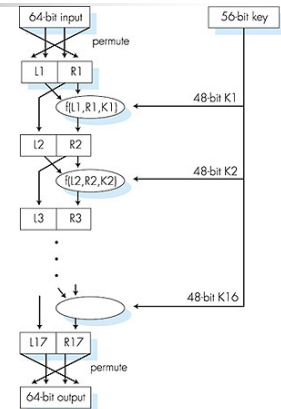
- US encryption standard [NIST 1993]
- 56-bit symmetric key, 64-bit plaintext input
- encryption**: initial permutation  $\Rightarrow$  16 "rounds", each using different 48 bits of key  $\Rightarrow$  final permutation
- decryption**: reverse operation using the same key

### How secure is DES?

- DES Challenge (1999): 56-bit-key-encrypted phrase decrypted (brute force) in 22 hours 15 minutes

### Making DES more secure:

- use three keys sequentially (3-DES)
- use more bits



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## AES: Advanced Encryption Standard

- Newer (Nov. 2001) symmetric-key NIST standard, replacing DES
- Processes data in 128 bit blocks
- 128, 192, or 256 bit keys
- Brute force decryption (try each key) taking 1 sec on DES, would take 149 trillion years for 128-bit AES
- Cost
  - 128-bit AES takes a few hundred microseconds to encrypt 4KB data on modern mobile processors
- Block cypher

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## Stream Cypher

- Much faster than block cypher like AES
- Fixed-size seed key  $K$  expanded (through some kind of shuffled bit shifts) to an infinite key stream  $C(K)$
- For data  $X$ , it uses a portion of the key stream (equal length to data  $X$ ), then xor the data to produce encrypted data
- Same key cannot be used twice, otherwise
  - $\text{Encrypted}(X) = X \text{ xor } C(K)$
  - $\text{Encrypted}(Y) = Y \text{ xor } C(K)$
  - $\Rightarrow \text{Encrypted}(X) \text{ xor } \text{Encrypted}(Y) = X \text{ xor } Y$
  - $\Rightarrow$  If you know  $X$  (or  $Y$ ), then you know the other
  - $\Rightarrow$  Even if you know neither, you can guess  $X/Y$  quite well from  $X \text{ xor } Y$  if  $X$  and  $Y$  are in a natural language

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## Stream Cypher

- How to guarantee the same key is not reused?
  - A shared master key known secretly by the two sides
  - An initialization vector, communicated publicly for each session, augment the master key to create the seed key and key stream for the session
- WEP wireless encryption
  - Employ an initialization vector of 24-bit, but insufficient (doesn't take long to observe sessions with the same initialization vector)

Sources: [http://en.wikipedia.org/wiki/Stream\\_cipher\\_attack](http://en.wikipedia.org/wiki/Stream_cipher_attack)

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## Public Key Cryptography

### Symmetric key cryptography

- requires sender, receiver know shared secret key
- Q: how to agree on key in the first place? (particularly difficult if Trudy is eavesdropping on all communication)

### Public key cryptography

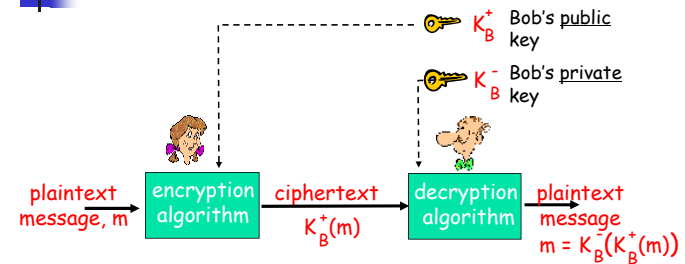
- encryption key is different from decryption key
- encryption key is **public**, known to **everyone**, also called **public key**
- decryption key is **secret**, known only to **receiver**, also called **private key**

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## Public Key Cryptography



Principle for choosing the public/private key pair:  
One should not be able to derive the private key from the public key.

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## Public Key Cryptography: RSA

(Ron Rivest, Adi Shamir and Len Adleman)

- Choosing keys:
  - Choose two large prime numbers  $p, q$ . (e.g., 1024 bits each)
  - Compute  $n = pq$ ,  $z = (p-1)(q-1)$
  - Choose  $e$  (with  $e < n$ ) that has no common factors with  $z$
  - Choose  $d$  such that  $ed-1$  is exactly divisible by  $z$
  - Public key is  $(n, e)$ . Private key is  $(n, d)$
- To encrypt a message,  $m$  ( $< n$ ): do  $c = m^e \bmod n$
- To decrypt a received ciphertext,  $c$ : do  $m = c^d \bmod n$
- Reason: for any  $m$  (relatively prime with  $n$ )
  - $m^z \bmod n = 1$ ; therefore  $m^{ed-1} \bmod n = 1$
- Another property:  $(m^d \bmod n)^e \bmod n = m$
- RSA is much slower than the symmetric key cryptos (block and certainly stream cyphers)

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