

# **Simultaneous Diophantine Approximation with Excluded Primes**

**László Babai**  
**Daniel Štefankovič**

# Dirichlet (1842)

## Simultaneous Diophantine Approximation

**Given reals**  $\alpha_1, \alpha_2, \dots, \alpha_n, Q$

**$\exists$  integers**  $r_1, \dots, r_n$  **and**  $q$

**such that**  $q \leq Q$  **and**

$$|\alpha_i q - r_i| \leq Q^{-1/n} \quad \mathbf{for\ all\ } i$$

$$|\alpha_i Q - p_i| \leq 1/2 \quad \mathbf{trivial}$$

# Simultaneous Diophantine Approximation with an excluded prime

Given reals  $\alpha_1, \alpha_2, \dots, \alpha_n$  prime  $p$

$\exists$ ? integers  $r_1, \dots, r_n$  and  $q$

such that  $\gcd(p, q) = 1$  and

$$|q\alpha_i - r_i| \leq \varepsilon \text{ for all } i$$

# Simultaneous diophantine $\varepsilon$ -approximation excluding $p$

**Not always possible**

**Example**  $p = 3$

**If**

$$\alpha_1 = 1/3$$

**then**

$$\varepsilon \geq |q\alpha_1 - r_1| = |q/3 - r_1| \geq 1/3$$

# Simultaneous diophantine $\varepsilon$ -approximation excluding $p$

**obstacle with 2 variables**

**If**

$$\alpha_1 + 2\alpha_2 = 1/p$$

**then**

$$\varepsilon \geq |q\alpha_1 - r_1|$$

$$\varepsilon \geq |q\alpha_2 - r_2|$$

$$3\varepsilon \geq |q(\alpha_1 + 2\alpha_2) - (r_1 + 2r_2)| \geq 1/p$$

# Simultaneous diophantine $\varepsilon$ -approximation excluding $p$

## general obstacle

**If**

$$b_1\alpha_1 + b_2\alpha_2 + \dots + b_n\alpha_n = 1/p + t$$

**then**

$$\varepsilon \sum |b_i| \geq 1/p$$

# Simultaneous diophantine $\varepsilon$ -approximation excluding $p$

## Theorem:

**If there is no  $\varepsilon$ -approximation excluding  $p$  then there exists an obstacle with**

$$\sum |b_i| \leq n^{3/2} / \varepsilon$$

## Kronecker's theorem ( ):

**Arbitrarily good approximation excluding  $p$  possible IFF no obstacle.**

**Simultaneous diophantine  $\varepsilon$ -approximation  
excluding  $p$**

**obstacle with**

$$\sum |b_i| \leq n^{3/2} / \varepsilon$$

**necessary to prevent  $\varepsilon$ -approximation  
excluding  $p$**

**sufficient to prevent  $\frac{\varepsilon}{pn^{3/2}}$ -approximation  
excluding  $p$**



**Motivating example**

**Shrinking by stretching**

# Motivating example

set  $A \subseteq (\mathbb{Z} / m\mathbb{Z})$

**arc length of A**

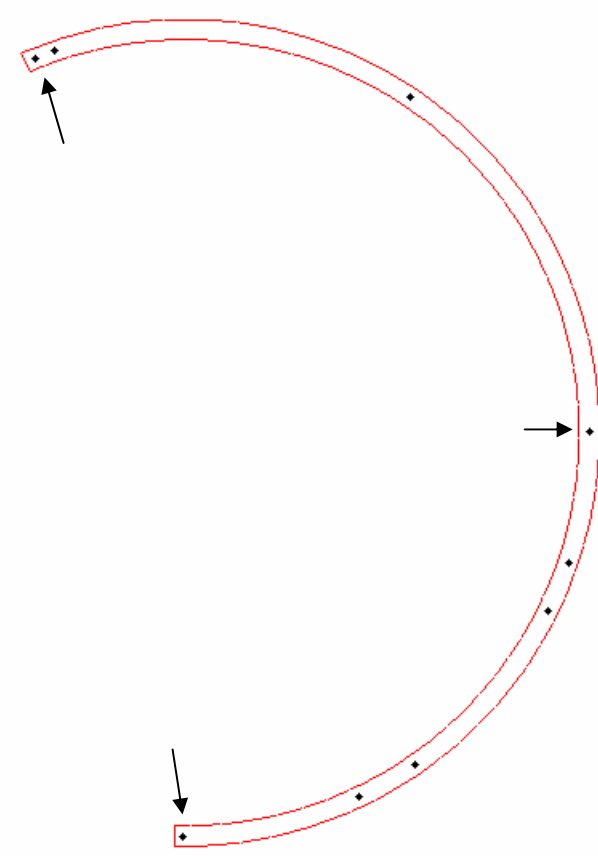
$$\max_{a \in A} |a(\bmod m)|$$

**stretching by  $x$**

$$a \mapsto ax \bmod m$$

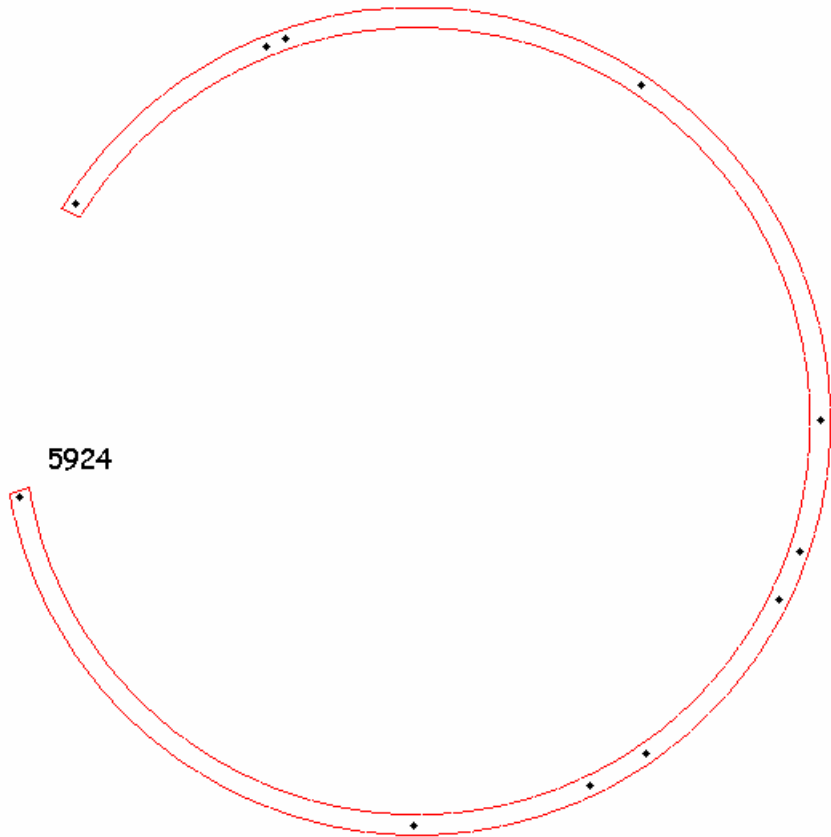
$$Ax = \{ax \mid a \in A\}$$

$$\gcd(x, m) = 1$$



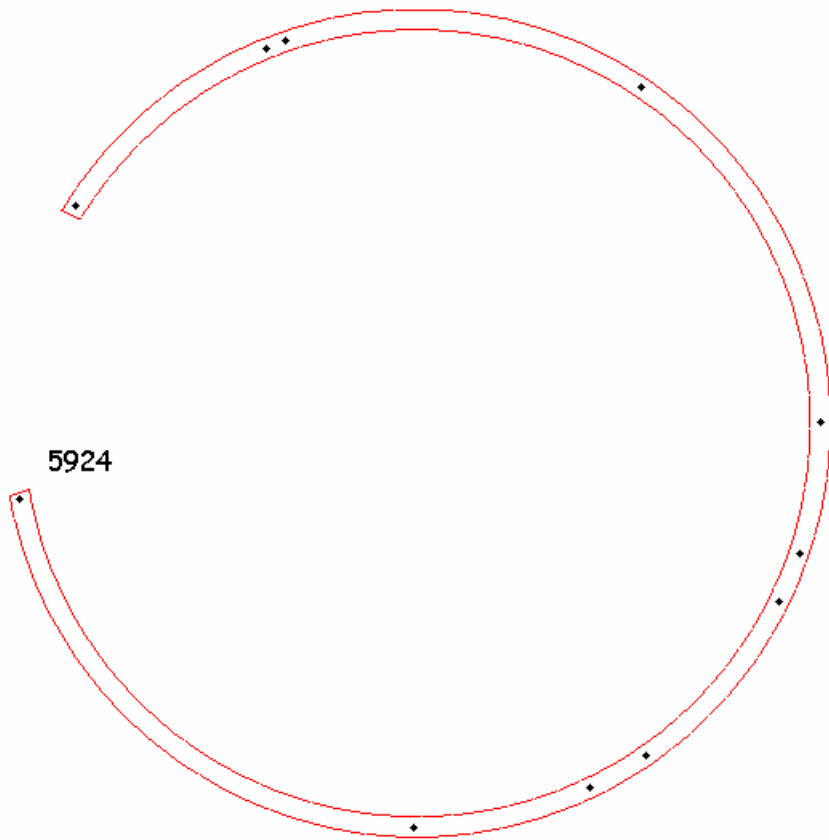
# Example of the motivating example

**A = 11-th roots of unity mod 11177**

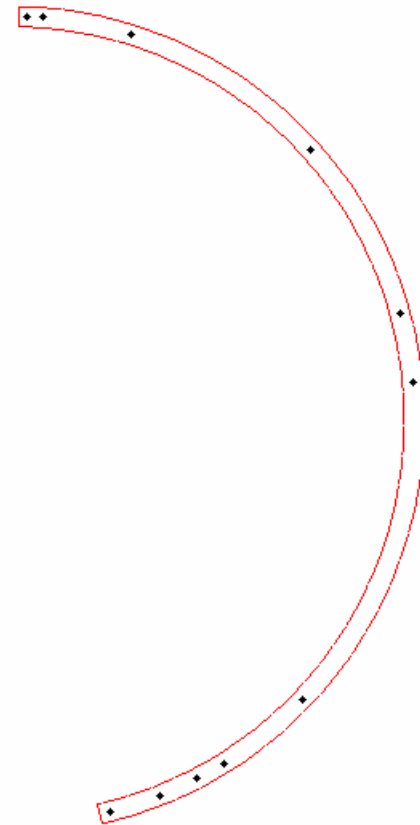


# Example of the motivating example

**A = 11-th roots of unity mod 11177**



**168**



# Shrinking modulo a prime

**If  $m$  a prime**

**then**

**every small set can be **shrunk****

# Shrinking modulo a prime

$m$  a prime

$$d = |A|$$

there exists  $x$  such that  
arc-length of  $Ax \leq m^{1-1/d}$

**proof:**

$$\frac{a_1}{m}, \dots, \frac{a_d}{m}$$

$$Q := m - 1$$

**Dirichlet**

$$\exists q; 0 < q \leq Q$$

$$x := q$$

$$|q\alpha_i - p_i| \leq \frac{1}{Q^{1/n}}$$

# Shrinking modulo any number

$m$  ~~a prime~~  $\longrightarrow$  every small set can be shrunk

?

# Shrinking modulo any number

~~$m$  a prime~~  $\longrightarrow$  ~~every small set can be shrunk~~

$$m = 2^k$$

$$A = \{1, 1 + 2^{k-1}\}$$

**If**

$$\gcd(x, m) = 1$$

**then the arc-length of  $Ax$**

$$\geq 2^{k-2}$$



# Where does the proof break?

$$m = 2^k$$

**proof:**

$$\frac{a_1}{m}, \dots, \frac{a_d}{m}$$

$$Q := m - 1$$

**Dirichlet**

$$\exists q; 0 < q \leq Q$$

$$x := q$$

$$|q\alpha_i - p_i| \leq \frac{1}{Q^{1/n}}$$

# Where does the proof break?

$$m = 2^k$$

**need:**

**approximation excluding 2**

**proof:**

$$\frac{a_1}{m}, \dots, \frac{a_d}{m}$$

$$Q := m - 1$$

**Dirichlet**

$$\exists q; \textcircled{0 < q \leq Q}$$

$$x := q$$

$$|q\alpha_i - p_i| \leq \frac{1}{Q^{1/n}}$$

# Shrinking cyclotomic classes

$m$  ~~a prime~~  $\longrightarrow$  ~~every small set can be shrunk~~

**set of interest – cyclotomic class**  
(i.e. the set of  $r$ -th roots of unity mod  $m$ )

- **locally testable codes**
- **diameter of Cayley graphs**
- **Waring problem mod  $p^k$**
- **intersection conditions modulo  $p^k$**

# Shrinking cyclotomic classes

**cyclotomic class**

**can be shrunk**

# Shrinking cyclotomic classes

**cyclotomic class**

**can be shrunk**

**Show that there is no small obstacle!**

## **Theorem:**

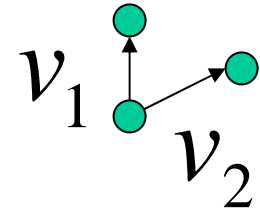
**If there is no  $\varepsilon$ -approximation excluding  $p$  then there exists an obstacle with**

$$\sum |b_i| \leq n^{3/2} / \varepsilon$$

# Lattice

$$v_1, \dots, v_n \in \mathbb{R}^n$$

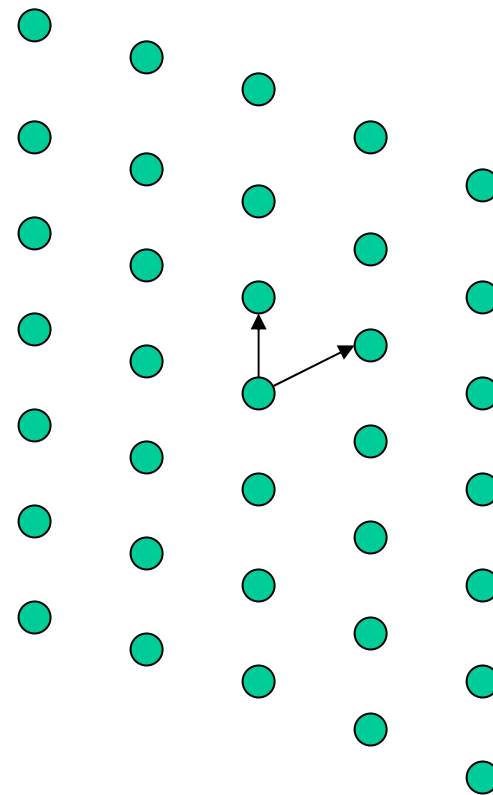
**linearly independent**



# Lattice

$$v_1, \dots, v_n \in \mathbb{R}^n$$

$$v_1\mathbb{Z} + \dots + v_n\mathbb{Z}$$

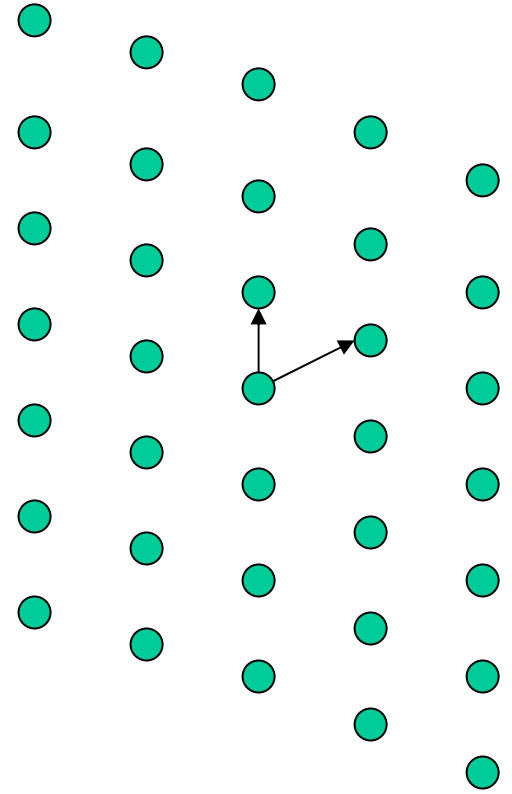




# Lattice

$$v_1, \dots, v_n \in \mathbb{R}^n$$

$$v_1\mathbb{Z} + \dots + v_n\mathbb{Z}$$



# Dual lattice

$$L^* = \{u \mid (\forall v \in L)v^T u \in \mathbb{Z}\}$$

# Banaszczyk's technique (1992)

**gaussian weight of a set**

$$\rho(A) = \sum_{x \in A} e^{-\pi \|x\|^2}$$

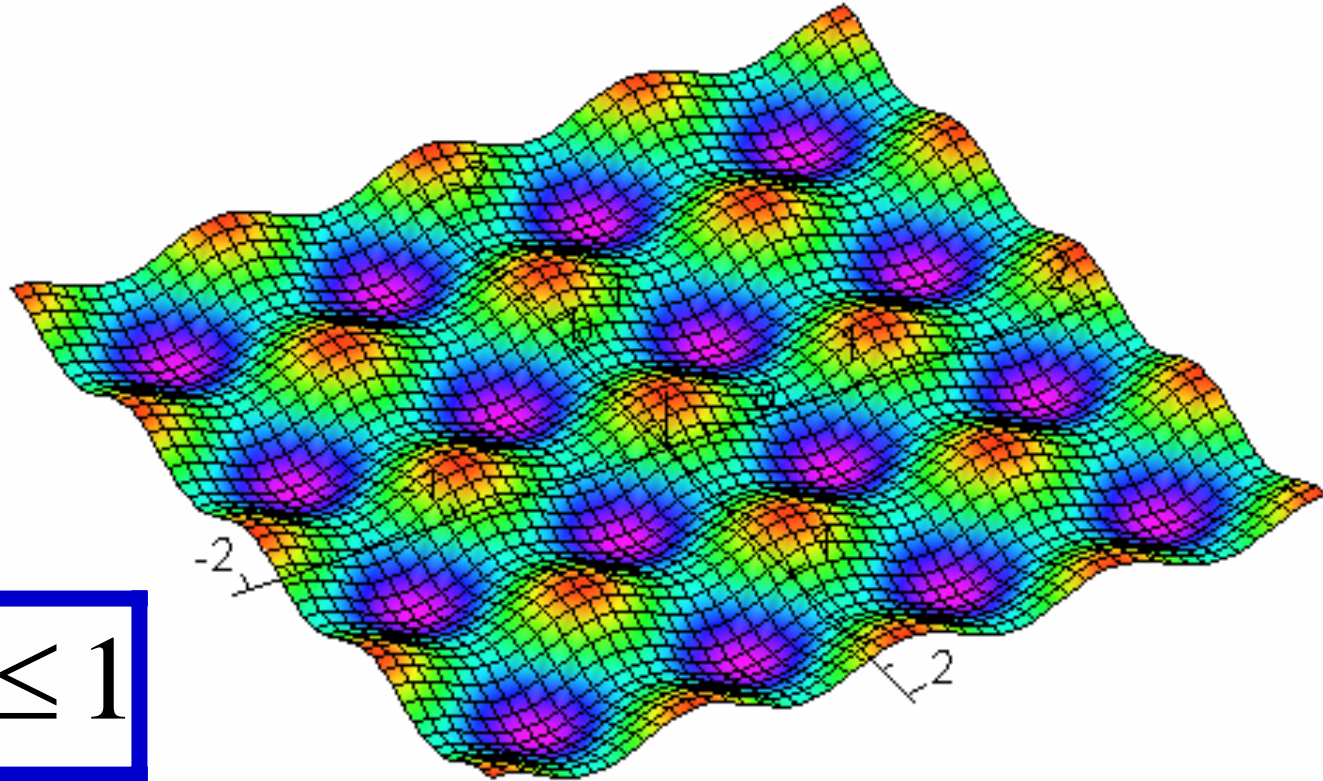
**mass displacement function of lattice**

$$\phi_L(x) = \rho(L + x) / \rho(L)$$

# Banasczyk's technique (1992)

mass displacement function of lattice

$$\phi_L(x) = \rho(L+x) / \rho(L)$$



**properties:**

$$0 \leq \phi_L(x) \leq 1$$

$$\text{dist}(x, L) \geq \sqrt{n} \quad \Rightarrow \quad \phi_L(x) \leq 1/4$$

# Banaszczyk's technique (1992)

**discrete measure**

$$\sigma_L(A) = \rho(L \cap A) / \rho(L)$$

**relationship between the discrete measure and the mass displacement function of the dual**

$$\hat{\sigma}_L(x) = \phi_{L^*}(x)$$

$$\hat{\sigma}_L(x) = \frac{1}{\rho(L)} \sum_{y \in L} \exp(-\pi \|y\|^2) \exp(2\pi i y^T x)$$

# Banasczyk's technique (1992)

discrete measure defined by the lattice

$$\sigma_L(A) = \rho(L \cap A) / \rho(L)$$

$$\hat{\sigma}_L(x) = \phi_{L^*}(x)$$

$$\frac{1}{\rho(L)} \sum_{\|x\| \leq s}^* \quad \frac{1}{\rho(L)} \sum_{\|x\| > s}^*$$

$$\hat{\sigma}_L(x) = \frac{1}{\rho(L)} \sum_{y \in L} \exp(-\pi \|y\|^2) \exp(2\pi i y^T x)$$

# Banasczyk's technique (1992)

$$\alpha_1, \alpha_2, \alpha_3 \longrightarrow \begin{pmatrix} 1 & 0 & 0 & \alpha_1 \\ 0 & 1 & 0 & \alpha_2 \\ 0 & 0 & 1 & \alpha_3 \\ 0 & 0 & 0 & \nu \end{pmatrix} \sqrt{n} / \varepsilon$$

**there is no short vector  $w \in L$   
with coefficient of the  
last column  $\not\equiv 0 \pmod{p}$**

# Banasczyk's technique (1992)

there is no short vector  $w \in L$   
with coefficient of the  
last column  $\not\equiv 0 \pmod{p}$

$$\hat{\sigma}_L(u) \geq 1/2$$

$$\phi_{L^*}(u) \geq 1/2$$

$$\text{dist}(u, L^*) \leq \sqrt{n}$$

**obstacle**

$$u := \frac{\varepsilon}{pv\sqrt{n}} e_{n+1}$$

**QED**

**Lovász (1982)**

## **Simultaneous Diophantine Approximation**

**Given rationals**  $\alpha_1, \alpha_2, \dots, \alpha_n, Q$

**can find in polynomial time**

**integers**  $p_1, \dots, p_n$   $0 < q \leq Q$

$$\left| q\alpha_i - p_i \right| \leq \frac{2^{n^2}}{Q^{1/n}} \quad \text{for all } i$$

**Factoring polynomials with rational coefficients.**



# Simultaneous diophantine $\varepsilon$ -approximation excluding $p$ - algorithmic

Given rationals  $\alpha_1, \alpha_2, \dots, \alpha_n$ , prime  $p$

can find in **polynomial time**

$2C_{n+1}p\varepsilon$ -approximation excluding  $p$

where  $\varepsilon$  is smallest such that there exists  $\varepsilon$ -approximation excluding  $p$

$$C_n = 4\sqrt{n}2^{n/2}$$



# Excluding prime and bounding denominator

**If there is no  $\varepsilon$ -approximation  
excluding  $p$  with  $q \leq Q$   
then there exists an**

**approximate obstacle with**

$$\sum |b_i| \leq n^{3/2} / \varepsilon$$

$$b_1 \alpha_1 + b_2 \alpha_2 + \dots + b_n \alpha_n = 1/p + t + \kappa$$

$$|\kappa| \leq n/Q$$

# Excluding prime and bounding denominator

## the obstacle

**necessary to prevent  $\varepsilon$ -approximation  
excluding  $p$  with  $q \leq Q$**

**sufficient to prevent  
 $\varepsilon / (2n^{3/2} p)$ -approximation  
excluding  $p$  with  $q \leq Q / (2pn)$**

# Excluding several primes

**If there is no  $\varepsilon$ -approximation  
excluding  $p_1, \dots, p_k$   
then there exists**

**obstacle with**

$$\sum |b_i| \leq n^{1/2} (\max(n, k)) / \varepsilon$$

$$\sum_{i=1}^n b_i \alpha_i = \sum_{j \in A \subseteq [k]} 1/p_j + t$$

**Show that there is no small obstacle!**

$$m=7^k$$

**\***  
**m**

**primitive 3-rd root of unity**

**know**

$$1 + \omega + \omega^2 \equiv 0 \pmod{7^k}$$

**obstacle**

$$c_0 + c_1\omega = t7^{k-1}, \quad \gcd(t, 7) = 1$$

# Show that there is no small obstacle!

$$1 + \omega + \omega^2 \equiv 0 \pmod{7^k}$$

$$c_0 + c_1\omega = t7^{k-1}, \quad \gcd(t, 7) = 1$$

$$\text{Res}(1 + x + x^2, c_0 + c_1x)$$

$$\begin{bmatrix} 1 & 1 & 1 \\ c_0 & c_1 & 0 \\ 0 & c_0 & c_1 \end{bmatrix}$$

$\neq 0$

**divisible by**  $7^{k-1}$

$$\leq 2(c_0^2 + c_1^2)$$

$$c_1^2 - c_0c_1 + c_0^2$$

$$\varepsilon = \frac{4}{7^{(k-1)/2}} \rightarrow$$

**There is g with all  
3-rd roots**

$$[-(4\sqrt{7})m^{1/2}, (4\sqrt{7})m^{1/2}]$$

# Dual lattice

$$\left( \begin{array}{cccc} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ -\frac{\alpha_1}{\nu} & -\frac{\alpha_2}{\nu} & -\frac{\alpha_3}{\nu} & \frac{1}{\nu} \end{array} \right) \varepsilon / \sqrt{n}$$



# **Algebraic integers?**

**possible that a small integer  
combination with small coefficients  
is doubly exponentially close to  $1/p$**