"Homework problem sessions" are in CSB 209, 7:30-8:30pm on Oct. 1 (Wednesday), Oct. 15 (Wednesday); held by Sean Esterkin.

Homework is Thursday, Oct. 2 (there will be additional homework due Thursday Oct. 16).

"Exam problem session" is in CSB 209, 6:00pm-7:00pm on Oct. 22 (Wednesday); held by the instructor.

EXAM #2 will be on Thursday, Oct. 23.

## 1 <u>Homework</u> - solve and turn in

**3.1** (due Thursday, Oct 2)[10 points] Let G = (V, E) be a directed graph (digraph) given in the following adjacency-list representation: for each vertex  $v \in V$  we have a linked list of out-neighbors of v. Assume that  $V = \{1, \ldots, n\}$ ; let m := |E|.

- 1. Write pseudocode for a procedure which outputs an adjacency-list representation of the reverse digraph (i.e., G with each edge reversed). The procedure should run in time O(m + n).
- 2. Write pseudocode for a procedure which outputs the adjacency-list representation of G in which the outneighbors of each vertex are listed in the increasing order. The procedure should run in time O(m+n).
- 3. Write pseudocode for a procedure which checks if G is undirected (i.e., the reverse of every  $e \in E$  is also in E). The procedure should run in time O(m+n).

**3.2** (due Thursday, Oct 2)[10 points] Let G = (V, E) be an undirected graph; let n := |V|, m := |E|. A 2-coloring of G is an assignment  $\phi : V \to \{1, 2\}$  (that is, each vertex gets one of two colors) such that for all  $\{u, v\} \in E$  we have  $\phi(u) \neq \phi(v)$  (that is, endpoints of each edge get different colors). Give an O(m + n) algorithm which finds either a 2-coloring of a graph or an odd cycle.

**3.3** (due Thursday, Oct 2)[10 points] We have a road network consisting of one-way streets. There are a few proposals for a new one-way road to be built. You want to vote for the proposal that would result in the shortest distance between your home and your office. How can you, quickly, figure out the proposal to vote for? We now formalize the problem.

Let G = (V, E, w) be a directed graph with positive edge weights; let n := |V|, m := |E|. Let  $s, t \in V$  be two distinct vertices (s is your home, t is your office). There are k proposals  $(a_1, b_1, c_1), \ldots, (a_k, b_k, c_k)$ , where  $a_i, b_i \in V$  and  $c_i \in \mathbb{R}_+$   $((a_i, b_i, c_i)$  proposes connecting location  $a_i$  to location  $b_i$  with a road of length  $c_i$ ).

Let  $G_i$  be the graph obtained from G by adding edge  $(a_i, b_i)$  with weight  $c_i$  (note that  $G_i$  has m + 1 edges (for each  $i \in \{1, \ldots, k\}$ )). We want to find which of  $G_1, \ldots, G_k$  has the shortest distance from s to t. Describe an  $O(m + n \log n + k)$  algorithm for the problem (no pseudocode is needed, just a clear description of the algorithm; you can use any algorithm that was mentioned in the class or in the book).

## 2 <u>Bonus Homework</u> - solve and turn in

**3.4** (due Thursday, Oct 2)[10 points] Let G = (V, E, w) be an undirected graph with positive edge weights; let m := |E|, n := |V|. Let  $(a_1, b_1, c_1), \ldots, (a_k, b_k, c_k)$  be a collection of triples, where  $a_i, b_i \in V$  and  $c_i \in \mathbb{R}_+$  (for  $i \in \{1, \ldots, k\}$ ). Let  $G_i$  be the graph obtained from G by adding edge  $\{a_i, b_i\}$  with weight  $c_i$ . We want to find the cost of the minimum weight spanning tree of  $G_i$ , for each  $i \in \{1, \ldots, k\}$ . (Thus the output of our procedure is k numbers.)

Describe an  $O(n^2 + k)$  algorithm for the problem. Clearly describe your algorithm in words, give a pseudocode, and argue why the running time is  $O(n^2 + k)$  (if you have a faster algorithm that is, of course, OK). You can use any algorithm that was mentioned in the class or in the book (without writing the pseudocode for that algorithm).

## 3 Additional problems from the book

Try to solve the following problems. A few of them <u>will be</u> on the quiz. We will go over the ones that you choose in the problem sessions.

- $\bullet \ \ 3.1, \ 3.2, \ 3.3, \ 3.4, \ 3.6, \ 3.7, \ 3.8, \ 3.9, \ 3.11, \ 3.12, \ 3.13, \ 3.15, \ 3.18, \ 3.19, \ 3.21, \ 3.22, \ 3.23, \ 3.24, \ 3.25, \ 3.26, \ 3.27, \ 3.26, \ 3.26, \ 3.27, \ 3.26, \ 3.26, \ 3.27, \ 3.26, \ 3.26, \ 3.26, \ 3.26, \ 3.27, \ 3.26,$
- $\bullet \ \ 4.1, \ 4.2, \ 4.3, \ 4.4, \ 4.5, \ 4.7, \ 4.8, \ 4.10, \ 4.11, \ 4.13, \ 4.15, \ 4.17, \ 4.18, \ 4.19, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.21, \ 4.20, \ 4.20, \ 4.20, \ 4.21, \ 4.20,$
- 5.1, 5.2, 5.3, 5.7.