## CSC 172– Data Structures and Algorithms

Lecture #10 Spring 2018

Please put away all electronic devices

- From now on, no electronic devices allowed during lecture
  - Includes Phone and Laptop
  - Why?
    - For your own good
    - And for others
  - What should I do instead?
    - Take your notebook out and start keeping notes!
  - But, I can multitask!
    - You can just stay at home and watch lecture videos!

#### Announcements

• Lab 5 is out (List ADT)

- Project 2 is out
  - Please finish the project before the spring recess
  - I really do not want you work on the project during the recess
  - The project looks really simple but probably is the hardest of all the projects.

## PROJECT 2 (UR CALCULATOR)

#### Link

<u>http://www.cs.rochester.edu/courses/172/spring</u>
 <u>2018/projects/proj2.pdf</u>



• What will be the output if I enter





• What will be the output if I enter





Result of evaluation should be in double.

• What will be the output if I enter





Assignments won't produce any output

• What will be the output if I enter





## Agenda

- We will talk about two new Data Structures
  - Мар
  - Stack
  - (Both useful for Project 2)
- We will see you how can you use these in Java.
- Later, we will learn how to implement them.
  - Stack: This week
  - Map: After MidTerm

#### Maps

• A Map is an object that maps keys to values.

• A map cannot contain duplicate keys

• Each key can map to at most one value.

## Map stores (key,value)

#### Students HashMap

Кеу	Value
Alice	15
Bob	12
Carol	108
Dave	105

Students.get("Bob") will give 12

System.out.println(students.get("Alice"));
System.out.println(students.get("David"));
System.out.println(students.get("Dave"));

15			
null			
105			



• Stacks

• Well-ballanced expressions

• Infix and postfix expressions

- Well-formed expressions
- stacks
- Infix, postfix

## **STACKS AND APPLICATIONS**

## **STACK IN JAVA**

#### Java Stack

Modifier and Type	Method and Description
boolean	empty()
	Tests if this stack is empty.
E	peek()
	Looks at the object at the top of this stack without removing it from the stack.
E	<b>pop</b> ()
	Removes the object at the top of this stack and returns that object as the value of this function.
Е	<pre>push(E item)</pre>
	Pushes an item onto the top of this stack.
int	<pre>search(Object o)</pre>
	Returns the 1-based position where an object is on this stack.

## Stack: push(obj), pop(), and peek()



#### **Practice problem**

```
public static void main (String[] args)
    Ł
        Stack<Integer> stack = new Stack<Integer>();
        for(int i = 0; i <= 10; i=i+2)</pre>
        {
            stack.push(i);
        }
        System.out.println("Top of stack = "+ stack.peek());
        System.out.println("Popping element = "+ stack.pop());
        System.out.println("Top of stack = "+ stack.peek());
        System.out.println("Popping element = "+ stack.pop());
        System.out.println("Popping element = "+ stack.pop());
        System.out.println("Top of stack = "+ stack.peek());
        stack.push(12);
        System.out.println("Top of stack = "+ stack.peek());
                                 Top of stack = 10
    }
```

```
Popping element = 10
Top of stack = 8
Popping element = 8
Popping element = 6
Top of stack = 4
Top of stack = 12
```

## **Application: Parsing**

Most parsing uses stacks

Examples includes:

- Matching tags in HTML
- In Java or Well-formed expression, matching
  - ( ... )
  - [ ... ]
  - { ... }

## The first example will demonstrate parsing HTML

# We will show how stacks may be used to parse an HTML document

## HTML File

```
<div id="navigation">
 <div class="inner">
 <div id="searcher">
  <form method="get" action="http://www.gnu.org/cgi-bin/estseek.cgi">
   <div><label class="netscape4" for="phrase">Search:</label>
   <input name="phrase" id="phrase" type="text" size="18" accesskey="s"
          value="Why GNU/Linux?" onfocus="this.value=''' />
   <input type="submit" value="Search" /></div><!-- unnamed label -->
  </form>
 </div><!-- /searcher -->
 <ul>
  id="tabPhilosophy"><a href=</li>
                         "/philosophy/philosophy.html">Philosophy</a>
  id="tabLicenses"><a href="/licenses/licenses.html">Licenses</a>
  id="tabEducation"><a href="/education/education.html">Education</a>
  id="tabSoftware"><a href="/software/software.html">Downloads</a>
  id="tabDoc"><a href="/doc/doc.html">Documentation</a>
  id="tabHelp"><a href="/help/help.html">Help&nbsp;GNU</a>
  <a
href="https://www.fsf.org/associate/support freedom?referrer=4052">Join the 
FSF! < /a > 
 </div><!-- /inner -->
```

</div><!-- /navigation -->

#### HTML is made of nested

- opening tags, e.g., <some\_identifier>, and
- matching closing tags, e.g., </some\_identifier>
  - <html>
    - <head>
      - <title>Hello</title>
    - </head>
    - <body>This appears in the
      - <i>browser</i>.
    - </body>
  - </html>

*Nesting* indicates that any closing tag must match the most <u>recent</u> opening tag

Strategy for parsing HTML:

- read though the HTML linearly
- place the opening tags in a stack
- when a closing tag is encountered, check that it matches what is on top of the stack
  - If yes, take it out. If no, error...
  - [Most Browser knows how to recover]

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>			
---------------	--	--	--

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>
---------------

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>	<head></head>	<title></title>	
---------------	---------------	-----------------	--

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html> <head></head></html>	<title></title>
-----------------------------	-----------------

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html> <body></body></html>	
-----------------------------	--

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>	<body></body>		
---------------	---------------	--	--

<html>

<head><title>Hello</title></head> <body>This appears in the <i>browser</i>.</body>

<html> <body< th=""><th>y&gt;</th><th><i></i></th></body<></html>	y>	<i></i>
-------------------------------------------------------------------	----	---------

<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>	<body></body>		<i></i>
---------------	---------------	--	---------

<html>

<head><title>Hello</title></head>
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<i>browser</i>.</body>

<html> <body></body></html>		
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<html>

<head><title>Hello</title></head>
<body>This appears in the
<i>browser</i>.</body>

<html></html>	<body></body>		
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<html>

<head><title>Hello</title></head>
<body>This appears in the
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<html></html>			
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3.2.5.1

We are finished parsing, and the stack is empty

- Possible errors:
- a closing tag which does not match the opening tag on top of the stack
- a closing tag when the stack is empty
- the stack is not empty at the end of the document

#### **Well-Formed Expressions**

Or "balanced expressions":

- ([this is] { a number } 12345) # well-formed
- ([this is] { a number ) 12345} # not wf
- {[(34+4)/5] + 7}/4 # wf
- {[(34+4)/5} + 7]/4 # not wf

## (Recursive) Definition of WFE

- The empty sequence is well-formed.
- If A and B are well-formed, then the concatenation AB is well-formed
- If A is well-formed, then [A], {A}, and (A) are well-formed.
- How to we check if an expression is WF?
   Use a stack!

## Algorithm for Recognizing WFE

- Read the next delimiter token.
- If it is an open delimeter (i.e. [ ( {)
  - push it into the stack.
- If it is a close delimiter (i.e. ] ) })
  - match it with a corresponding open delimiter in the stack ([with ] and so on).
  - If there is no match  $\rightarrow$  not WF
  - If there is a match, stack.pop() and discard both
- When there is no more token left
  - If the stack is empty  $\rightarrow$  WF
  - If the stack is not empty ightarrow not WF

#### Infix vs Postfix Expressions

- Infix:  $5 + 4 \times 5/2 3$
- Postfix: 5 4 5 \* 2 / + 3 -

- Infix: (5+4)\*5/2 3
- Postfix: 5 4 + 5 \* 2 / 3 -

## **Postfix Expression Evaluation Algorithm**

- Initialize an empty stack
- While (there is still a token to read)
  - read the token t
  - if t is an operand, push it onto the stack
  - if t is an operator,
    - pop two operands from the stack, compute the result (using t) // if there is division by zero, scream
    - push the result back onto the stack // if there is less than two operands, scream
- In the end, if there is one number in the stack, output it.
  - // If there is more than one number in the stack, scream.

#### Example We will see

• Infix:

 $1 - (2 + 3 + (4 - 5 \times 6) \times 7) + 8 \times 9$ 

Postfix (Reverse-Polish Notation):

 $1 \ 2 \ 3 \ + \ 4 \ 5 \ 6 \ \times \ - \ 7 \ \times \ + \ - \ 8 \ 9 \ \times \ +$ 

Evaluate the following reverse-Polish(postfix) expression using a stack:

 $1 \ 2 \ 3 \ + \ 4 \ 5 \ 6 \ \times \ - \ 7 \ \times \ + \ - \ 8 \ 9 \ \times \ +$ 



#### Push 1 onto the stack 1 2 3 + 4 5 6 × $-7 \times + -89 \times +$



#### Push 1 onto the stack 1 2 3 + 4 5 6 × $-7 \times + -89 \times +$



Push 3 onto the stack 1 2 3 + 4 5 6 × -7 × + -8 9 × +



Pop 3 and 2 and push 2 + 3 = 51 2 3 + 4 5 6 × - 7 × + - 8 9 × +



Push 4 onto the stack 1 2 3 + 4 5 6 × -7 × + -8 9 × +



Push 5 onto the stack 1 2 3 + 4 5 6 ×  $-7 \times + -89 \times +$ 



Push 6 onto the stack 1 2 3 + 4 5 6 × - 7 × + - 8 9 × +



## Pop 6 and 5 and push $5 \times 6 = 30$ 1 2 3 + 4 5 6 × - 7 × + - 8 9 × +



## Pop 30 and 4 and push 4 - 30 = -261 2 3 + 4 5 6 × - 7 × + - 8 9 × +



Push 7 onto the stack 1 2 3 + 4 5 6 × -7 × + - 8 9× +



## Pop 7 and -26 and push $-26 \times 7 = -182$ 1 2 3 + 4 5 6 × - 7 × + - 8 9 × +



#### Pop -182 and 5 and push -182 + 5 = -1771 2 3 + 4 5 6 × $-7 \times + -89 \times +$



Pop -177 and 1 and push 1 - (-177) = 178 1 2 3 + 4 5 6 × - 7 × + - 8 9 × +



Push 8 onto the stack 1 2 3 + 4 5 6 ×  $-7 \times + -8 9 \times +$ 



#### Push 1 onto the stack 1 2 3 + 4 5 6 × $-7 \times + -89 \times +$



## Pop 9 and 8 and push 8 × 9 = 72 1 2 3 + 4 5 6 × $-7 \times + -89 \times +$



## Pop 72 and 178 and push 178 + 72 = 2501 2 3 + 4 5 6 × - 7 × + - 8 9 × +



#### How about Infix Expression?

• Shunting yard algorithm

• Convert infix to postfix

• Or, evaluate infix expressions directly

## Rough Idea (Shunting Yard Algorithm)

- Use 2 stacks: an operand stack, an operator stack
- If tok is an operand, push it on operand stack
- Else if tok is one of + \* /
  - while (precedence(tok) ≤ precedence(stack.peek())
    - Evaluate stack.peek()
  - Push tok on top of operator stack
- Else if tok is one of ( [ {
  - Push tok on top of operator stack
- Else if tok is one of ) ] }
  - Evaluate operators on top until ( [ { seen, match up

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