

# CSC 172– Data Structures and Algorithms

Lecture #16

Spring 2018

Please put away all electronic devices





# Announcement

- Extra credit opportunity:
  - Sorting dance
    - [http://www.cs.rochester.edu/courses/172/spring2018/sorting\\_dance.html](http://www.cs.rochester.edu/courses/172/spring2018/sorting_dance.html)
    - (Any sorting algorithm)
    - You need to have 10+ students
      - Multiple workshop teams together?!
    - You have to upload the video to YouTube and provide us the link by **April 23** (include the list of students participated)
    - Also, **2 bonus points** (divided among the number of participants) for the best dance!



# How it will impact

- It will heal 3 pts of your overall project or exam score.
- Example:
- Your Exam Score:
  - 20 out of 35
- Your Project Score:
  - 15 out of 30
- Your overall score =  $20+15 = 35$  out of 65

Sorting Dance will heal by replacing

$= \min(3 * 20/35, 3 * 15/30)$

$= \min(1.71, 1.5)$

$= 1.5$  pts

with 3 pts.

After the dance, your total score for Projects and Exams will be:

$20+15 - 1.5 + 3 = 36.5$

Extremely useful data structure

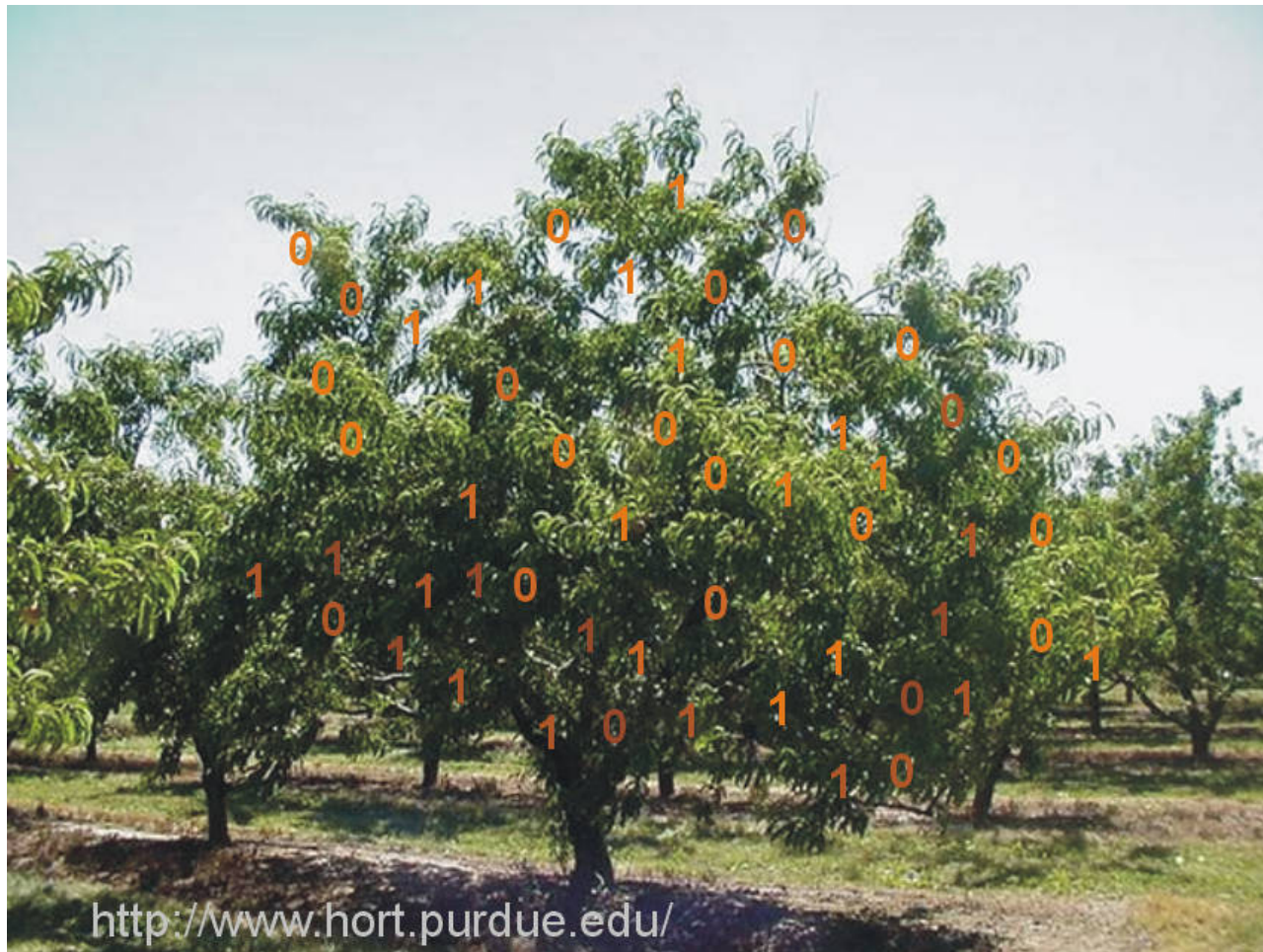
Special cases include

- Huffman tree
- Expression tree
- Decision tree (in machine learning)
- Heap data structure (later lecture)

# BINARY TREES

# Definition

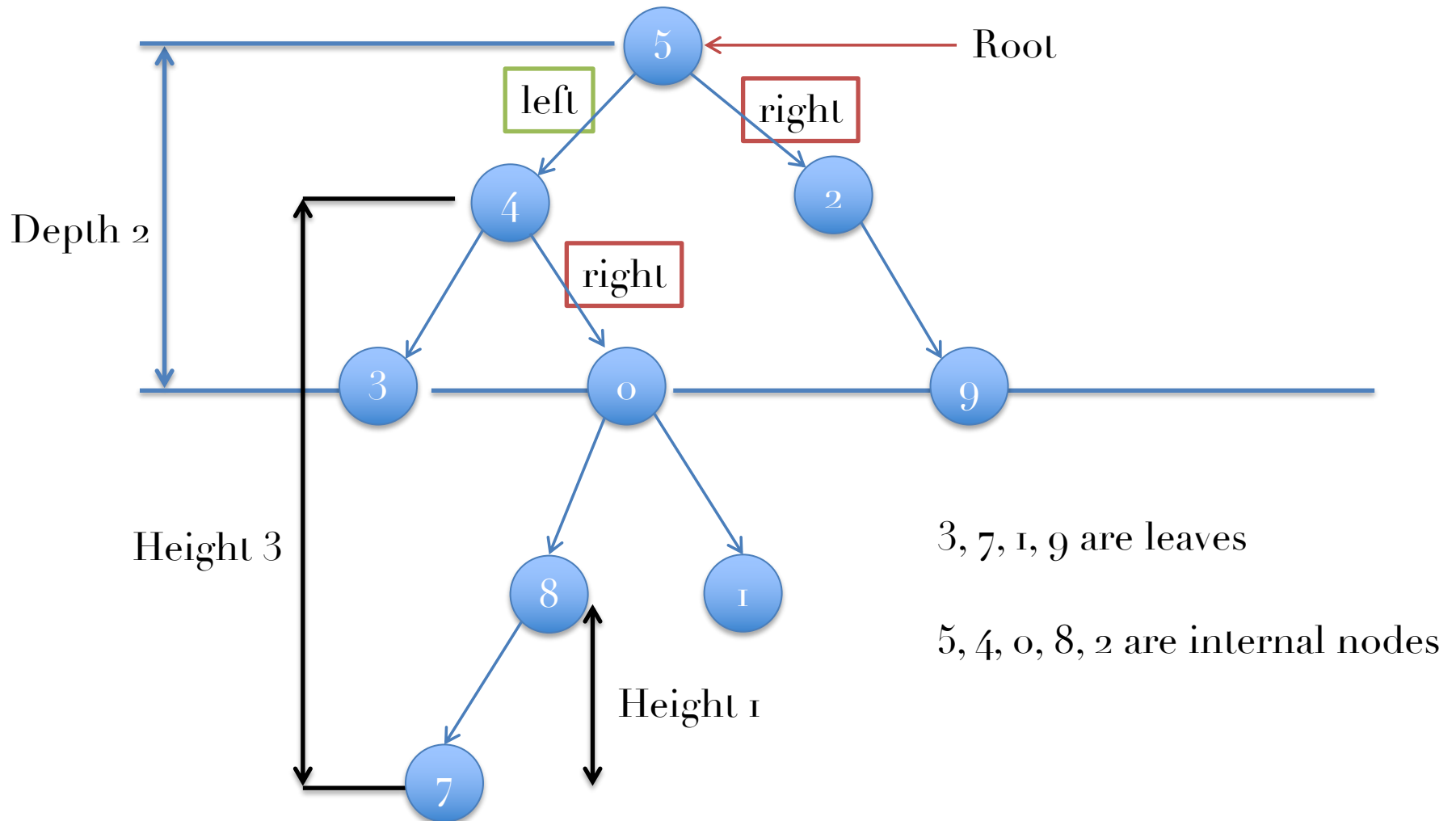
This is not a binary tree:



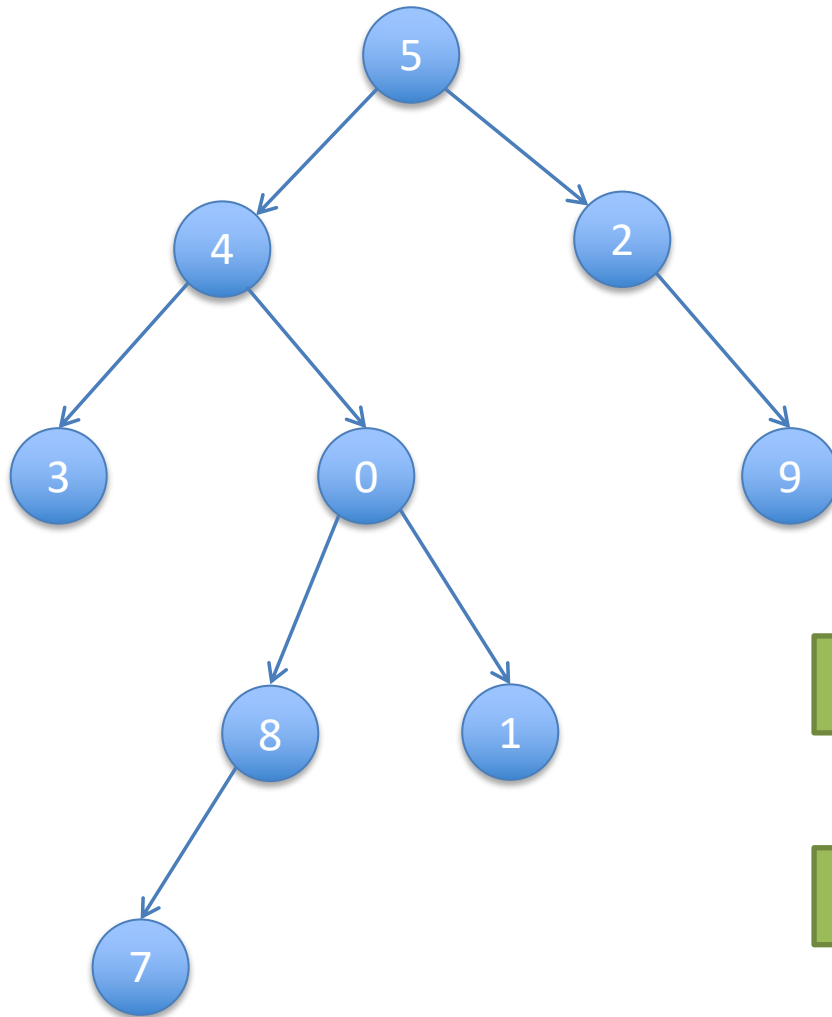
Neither is this



# Binary Trees



# Ancestors and Descendants

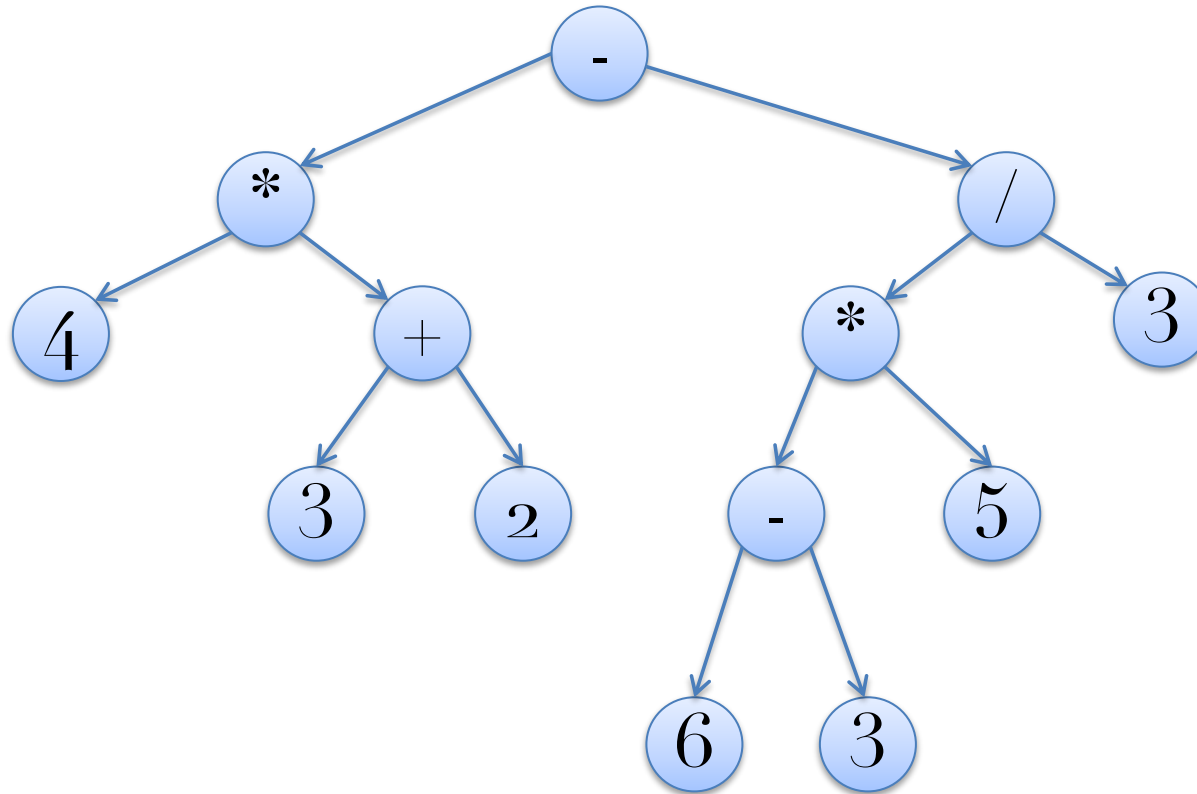


1, 0, 4, 5 are ancestors of 1

0, 8, 1, 7 are descendants of 0

# Application: Expression Trees

$$4 * (3 + 2) - (6 - 3) * 5 / 3$$

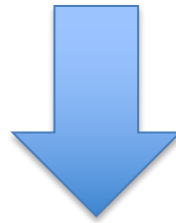




# How to construct Expression Trees?

Infix

$4 * (3 + 2) - (6 - 3) * 5 / 3$



Shunting Yard Algorithm

Postfix

4 3 2 + \* 6 3 - 5 \* 3 / -

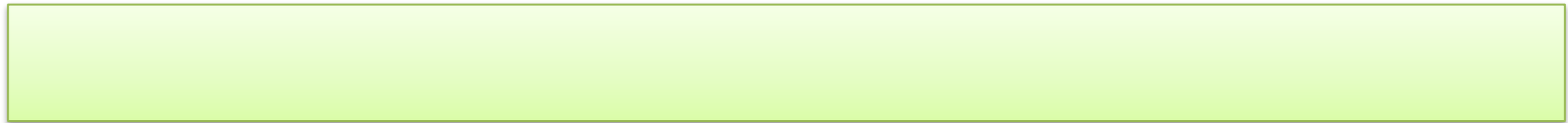


# How to construct Expression Trees?

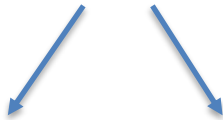
4 3 2 + \* 6 3 - 5 \* 3 / -



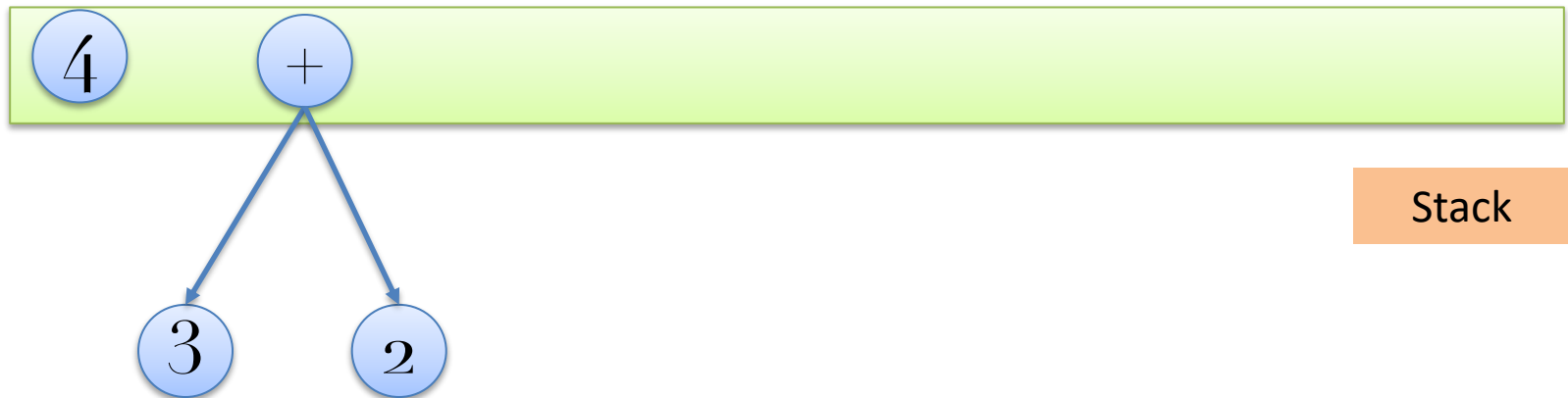
# How to construct Expression Trees?



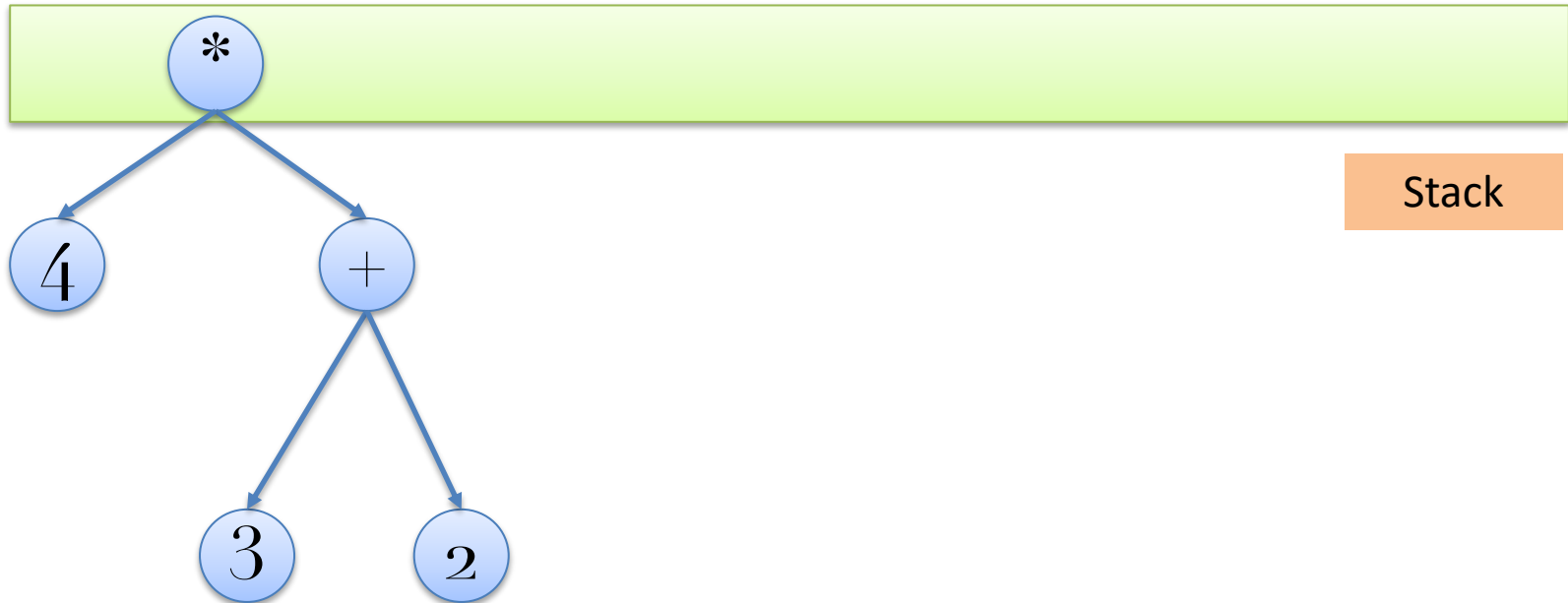
Stack



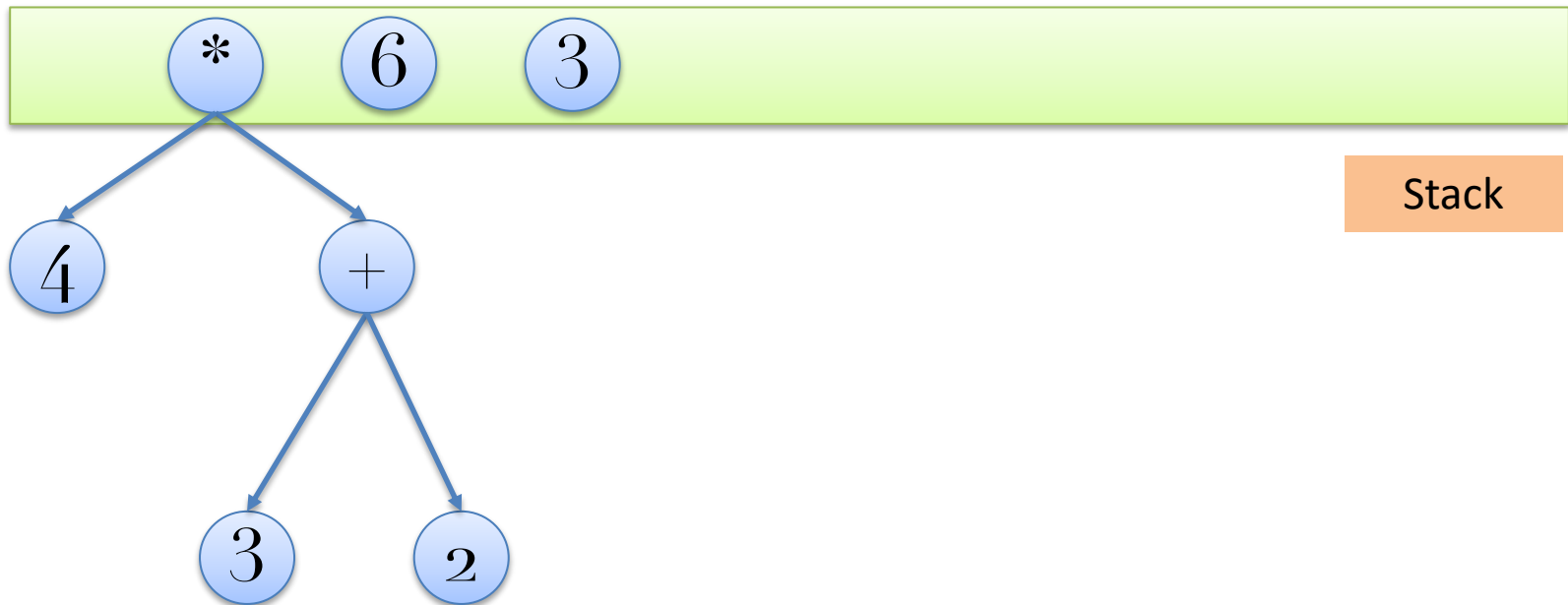
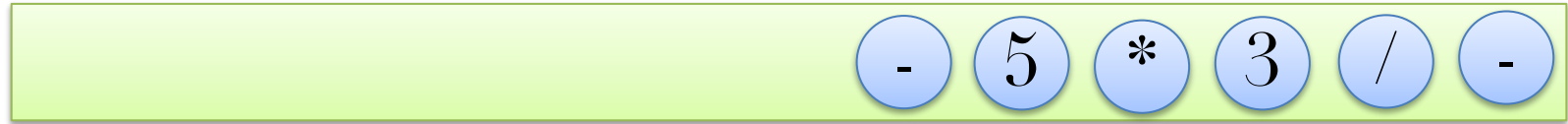
# How to construct Expression Trees?

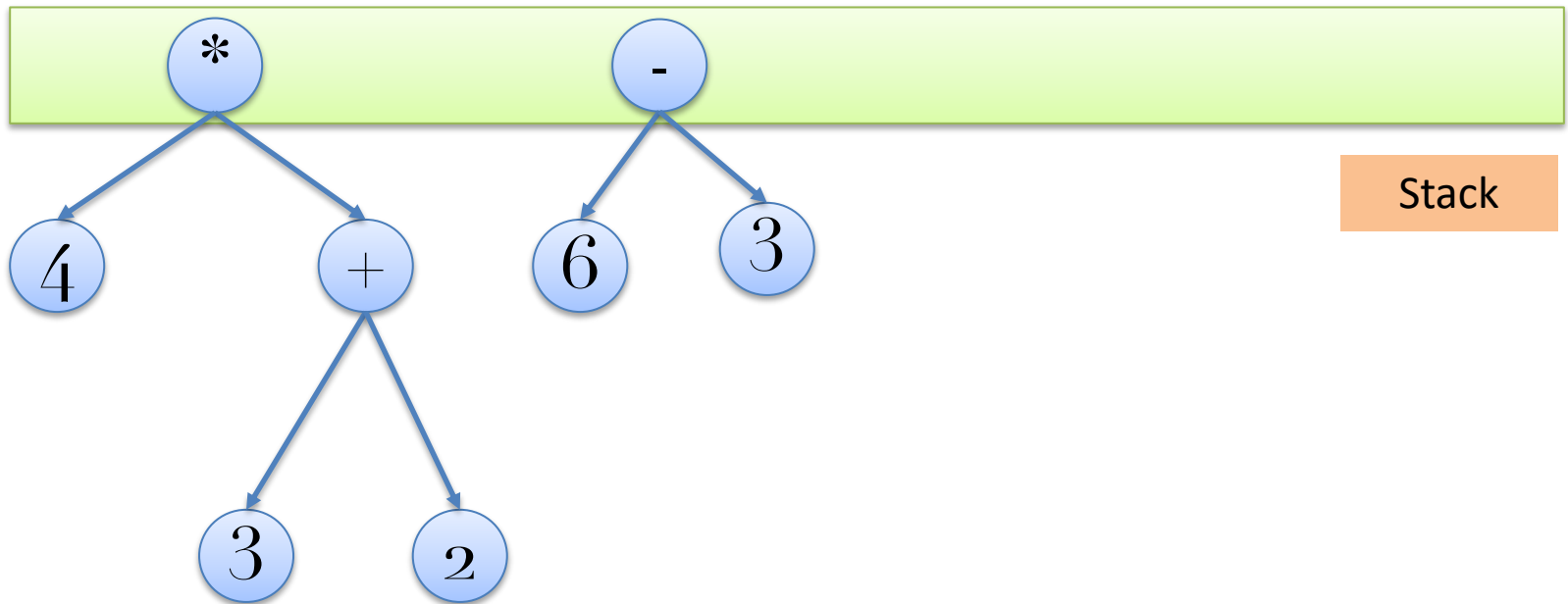
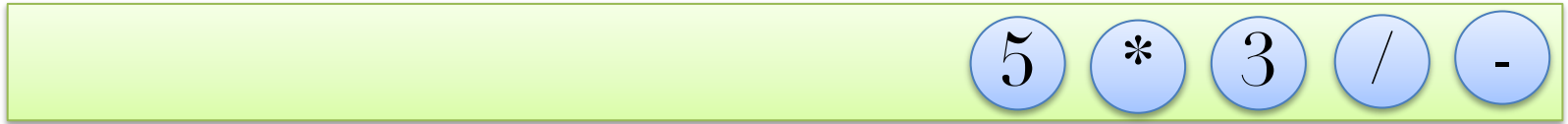


# How to construct Expression Trees?

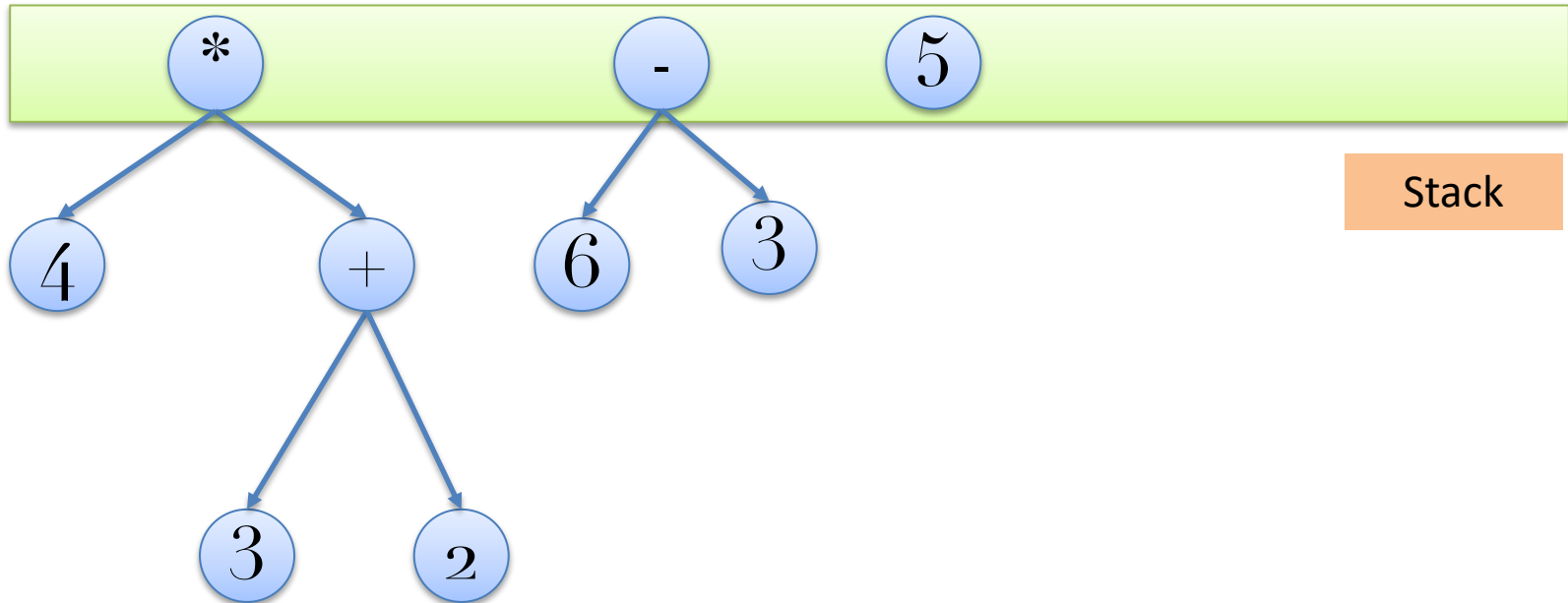
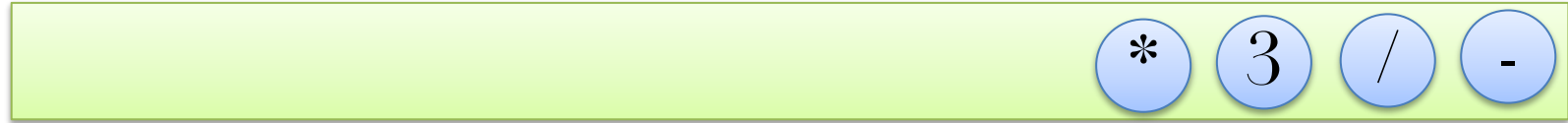


# How to construct Expression Trees?

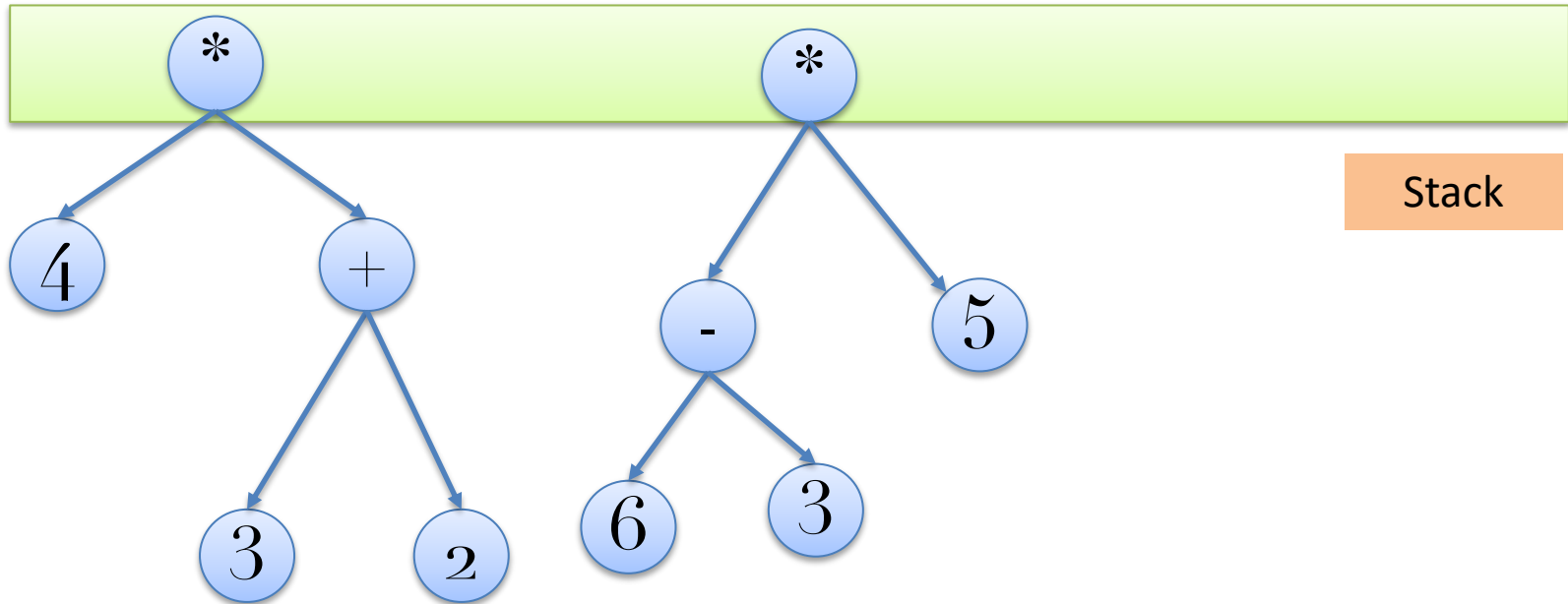
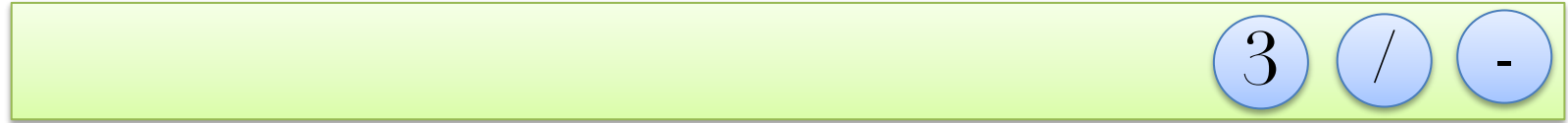




# How to construct Expression Trees?

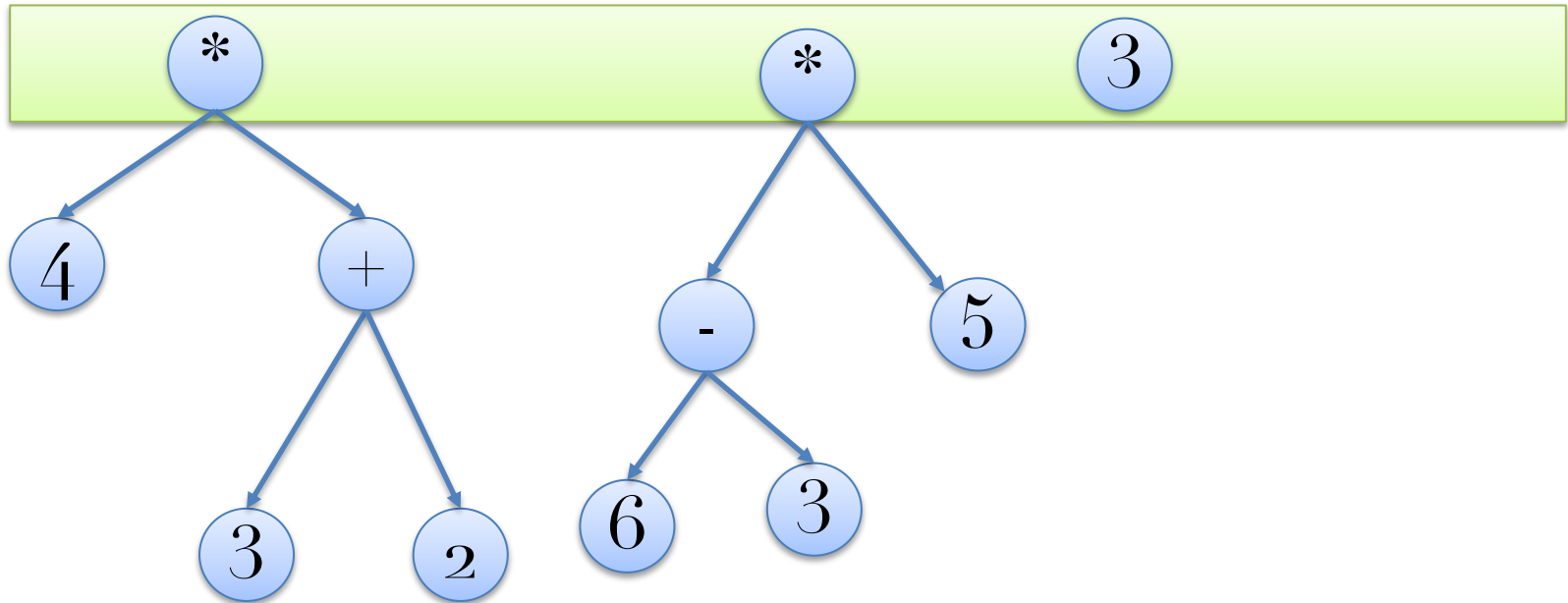
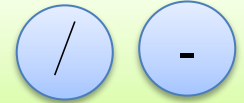


# How to construct Expression Trees?



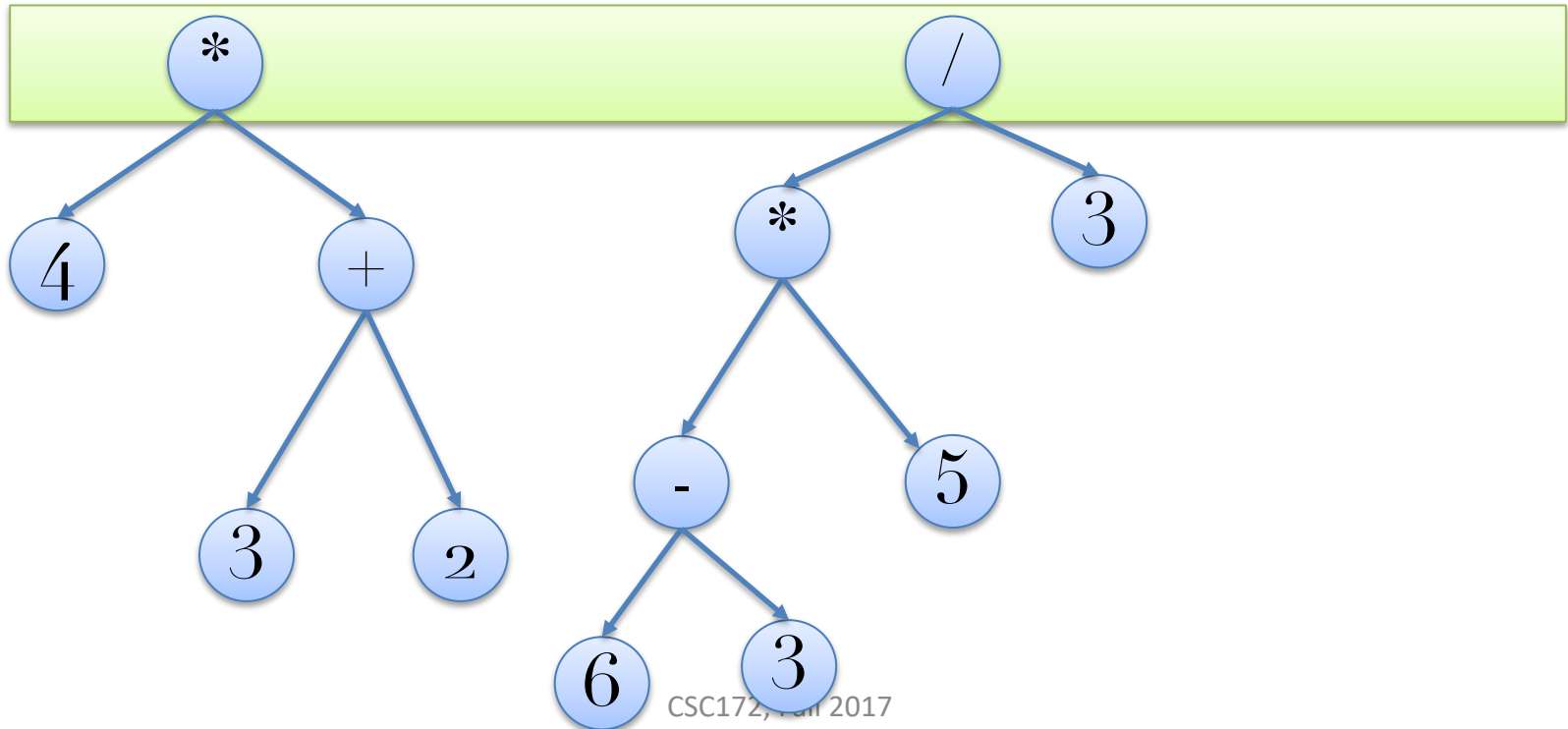


# How to construct Expression Trees?

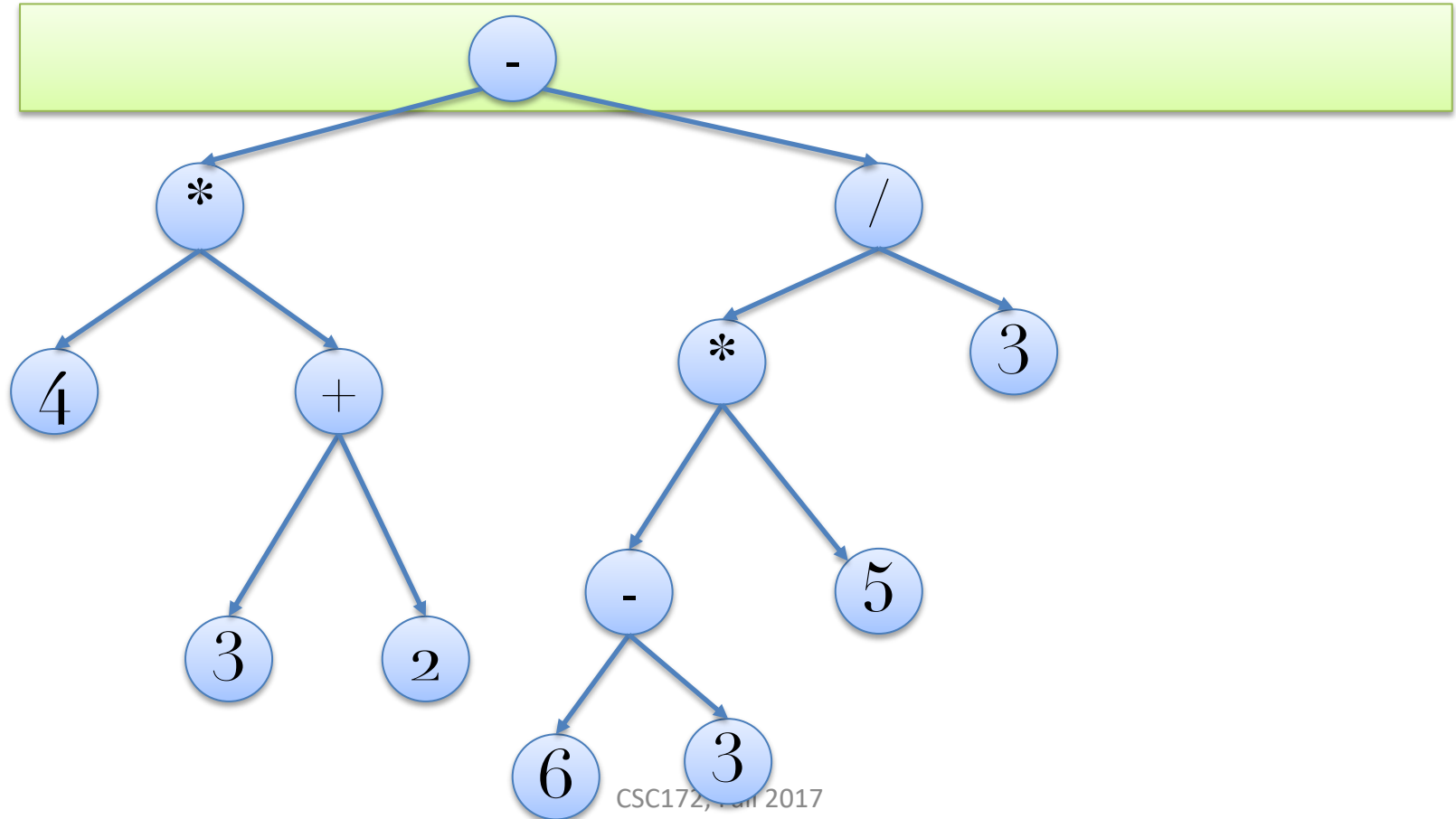


# How to construct Expression Trees?

-

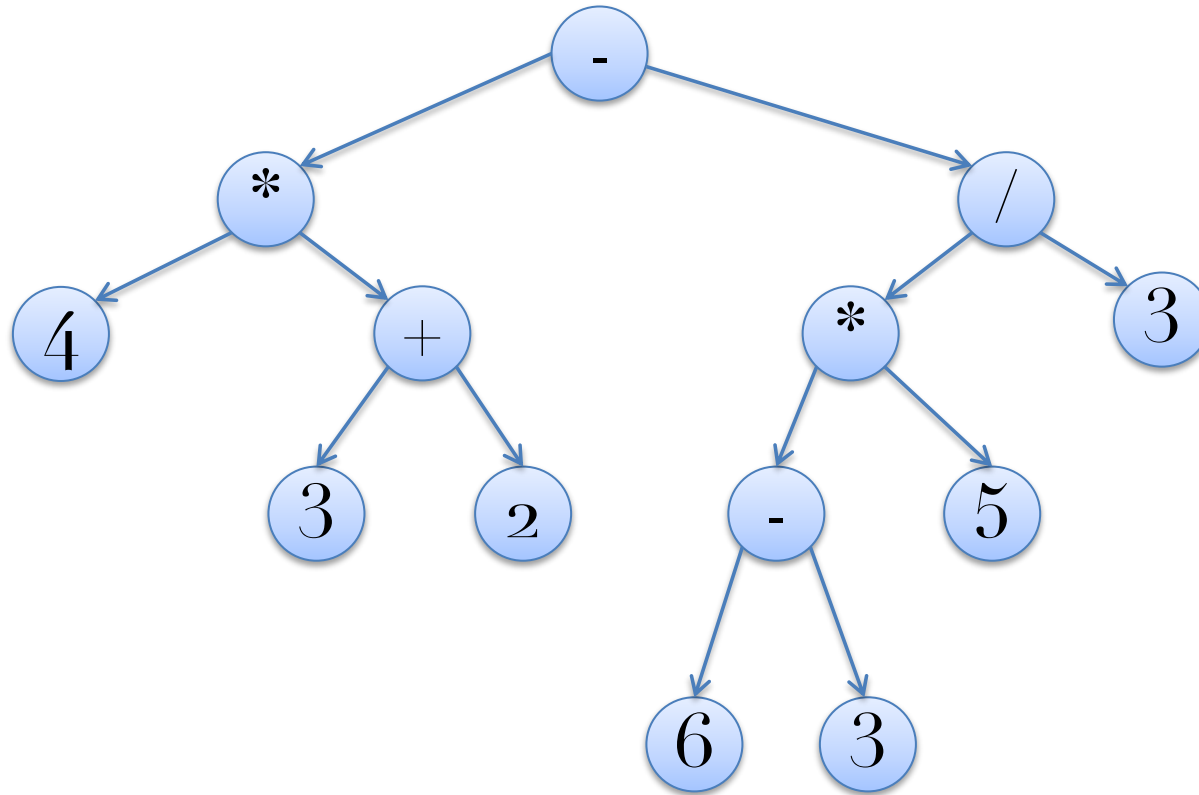


# How to construct Expression Trees?



# Finally!

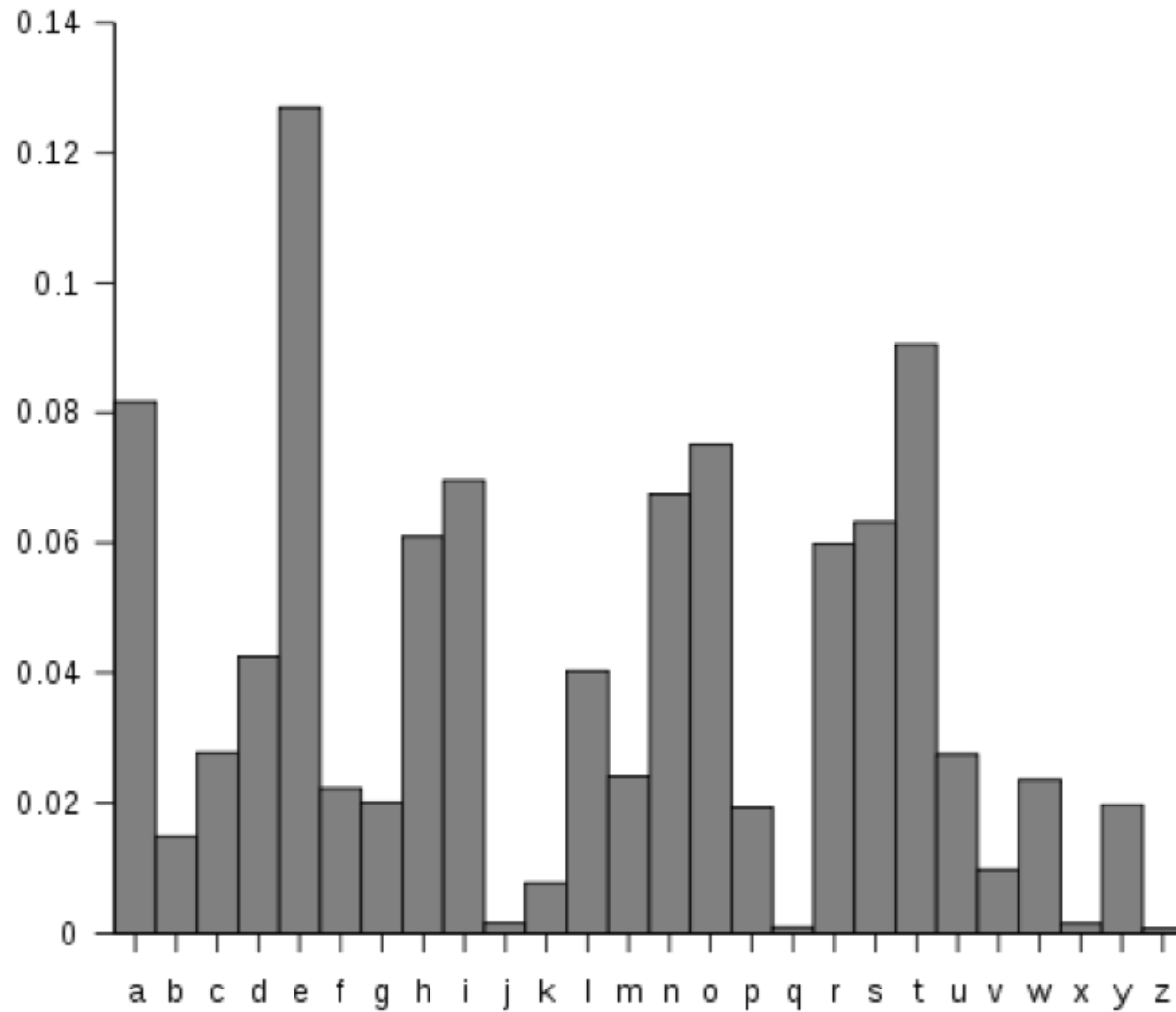
$$4*(3+2) - (6-3)*5/3$$



# Another Application: Character Encoding

- UTF-8 encoding:
  - Each character occupies 8 bits
  - For example, 'A' = 0x41
- A text document with  $10^9$  characters is  $10^9$  bytes long
- But characters were not born equal

# English Character Frequencies



# Variable-Length Encoding: Idea

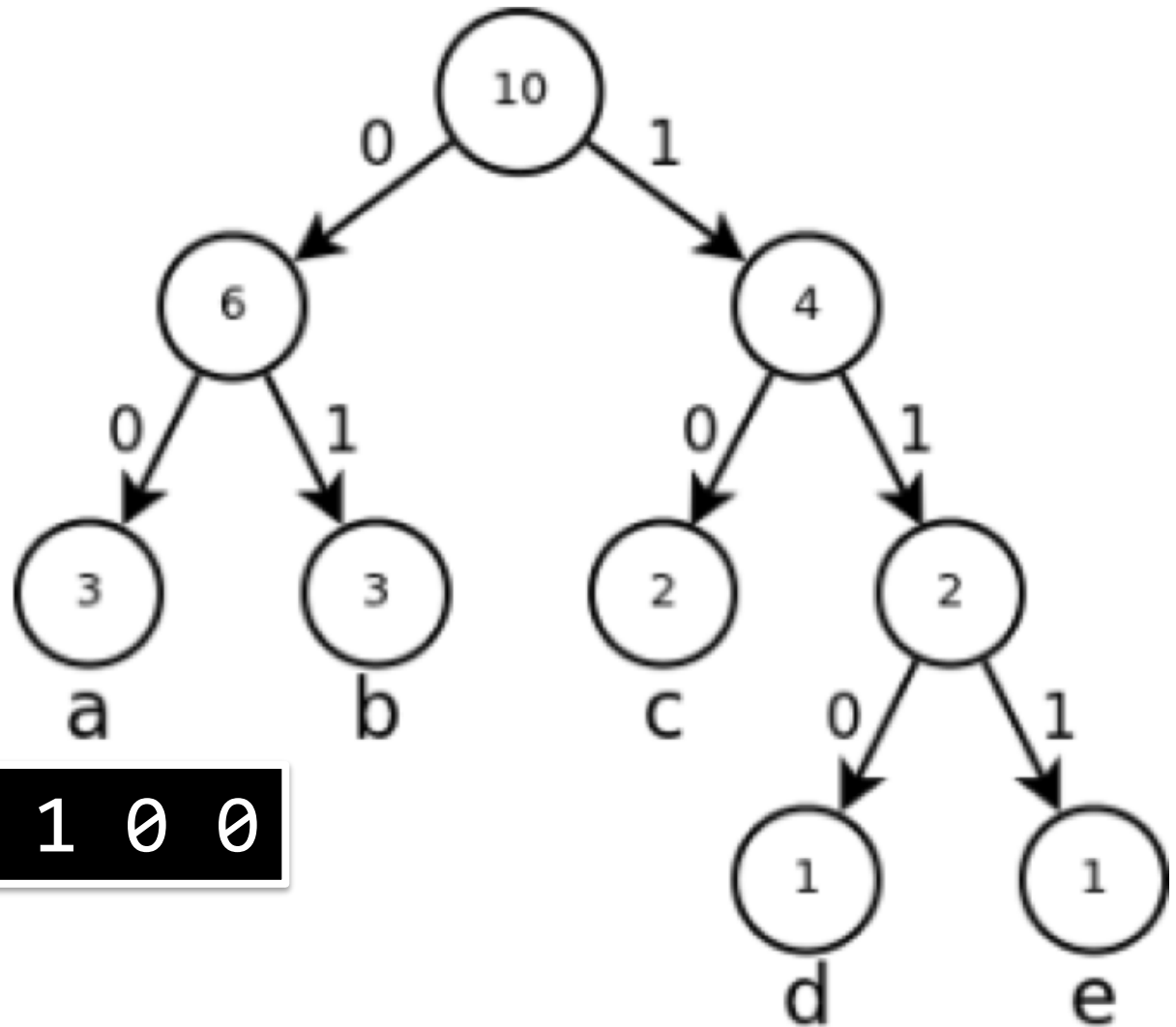
- Encode letter E with fewer bits, say  $b_E$  bits
- Letter J with many more bits, say  $b_J$  bits
- We gain space if

- $$b_E \cdot f_E + b_J \cdot f_J < 8f_E + 8f_J$$

where  $\mathbf{f}$  is the frequency vector

- Problem: how to decode?

# One Solution: Prefix-Free Codes



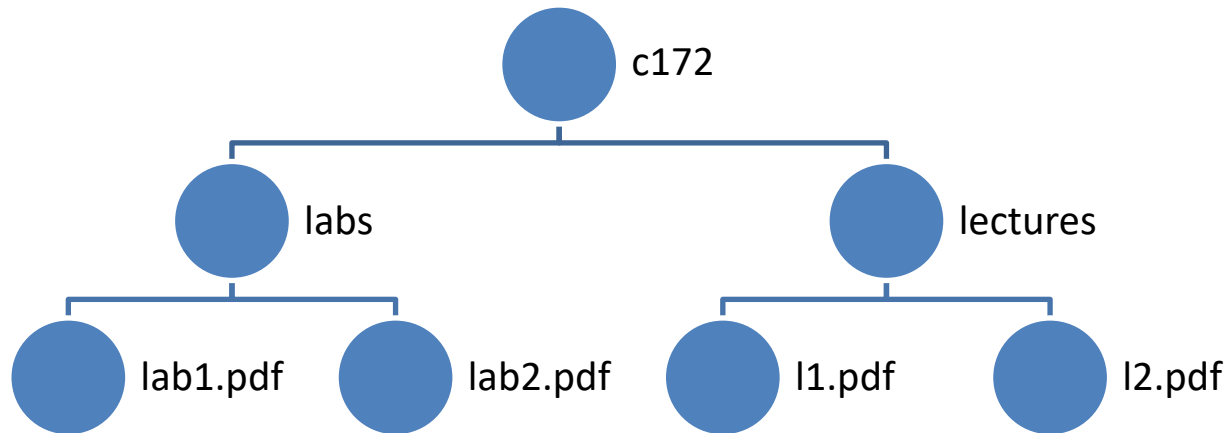
1 0 1 1 1 0 1 0 0

c e b a



# Printing a Hierarchy

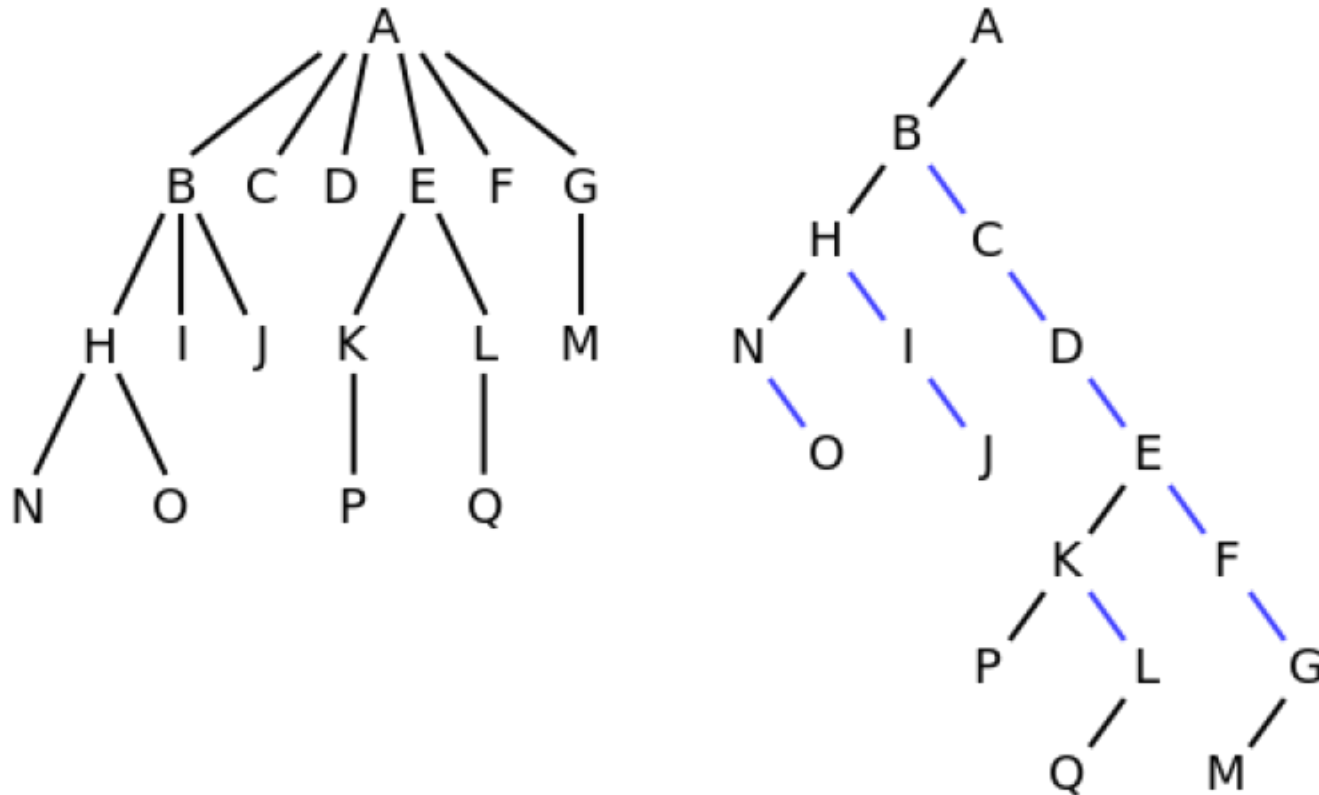
- The easiest way of printing a hierarchy (tree)



```
c172
├── bootstrap.html
├── calendar.html
├── demo.html
├── hide.js
├── images
│   └── textbook.jpg
├── index.html
├── index.html.backup
├── index.html.backup2
├── jq.js
├── labs
│   ├── lab10.pdf
│   ├── lab11.pdf
│   ├── lab1.pdf
│   ├── lab2.pdf
│   ├── lab3.pdf
│   ├── lab4.pdf
│   ├── Lab4.zip
│   ├── lab5.pdf
│   ├── lab6.pdf
│   ├── lab7.pdf
│   ├── lab7.tex
│   ├── Lab7.zip
│   ├── lab8.pdf
│   ├── lab9.pdf
│   ├── URLList.java
│   └── URNode.java
├── Labs.html
└── lectures
    ├── l10.pdf
    ├── l11.pdf
    ├── l12.pdf
    ├── l13.pdf
    ├── l14.pdf
    ├── l15.pdf
    ├── l16.pdf
    ├── l17.pdf
    ├── l18.pdf
    ├── l19.pdf
    ├── l1.pdf
    ├── l20.pdf
    ├── l21.pdf
    ├── l22.pdf
    ├── l23.pdf
    ├── l24.pdf
    ├── l25.pdf
    ├── l26.pdf
    ├── l27.pdf
    └── l2.pdf
```

# WHY ONLY BINARY TREE?

# Any Tree can be “Encoded” as a Binary Tree



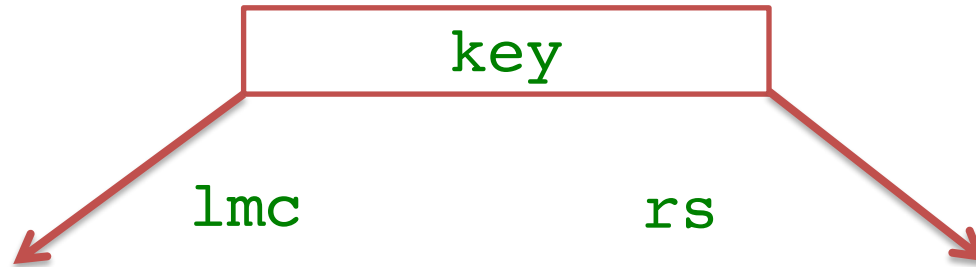
# LMC-RS Representation

- In this representation, every node has two pointers:
  - LMC (Left-most-child)
  - RS (Right Sibling)

# LMC-RS Representation

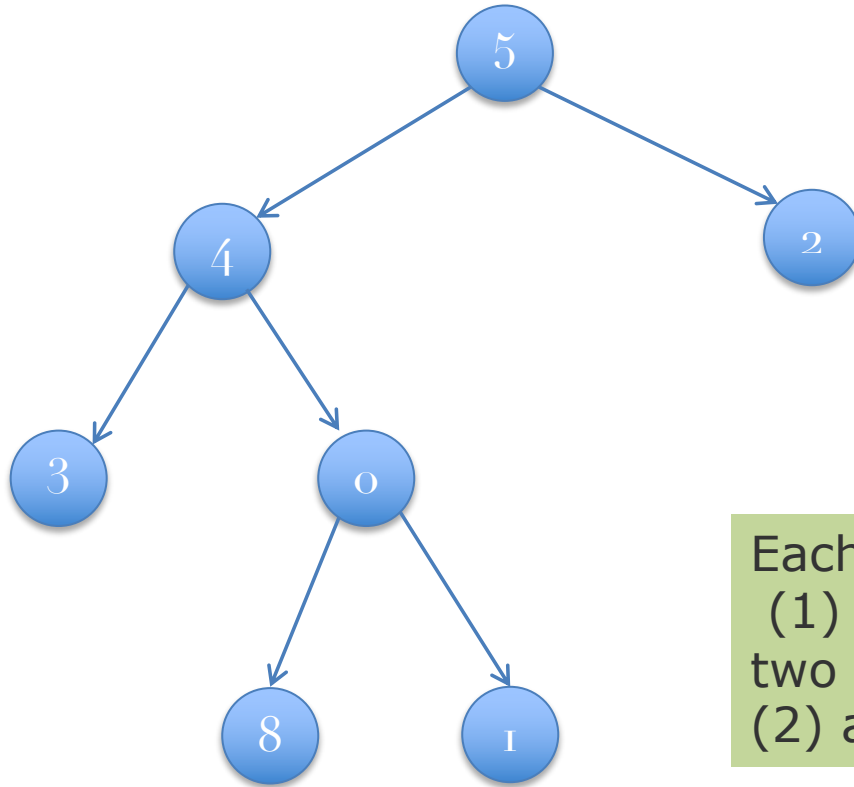
```
public class Node
{
    public int key;
    public Node lmc, rs;

    public Node(int item)
    {
        key = item;
        lmc= rs= null;
    }
}
```



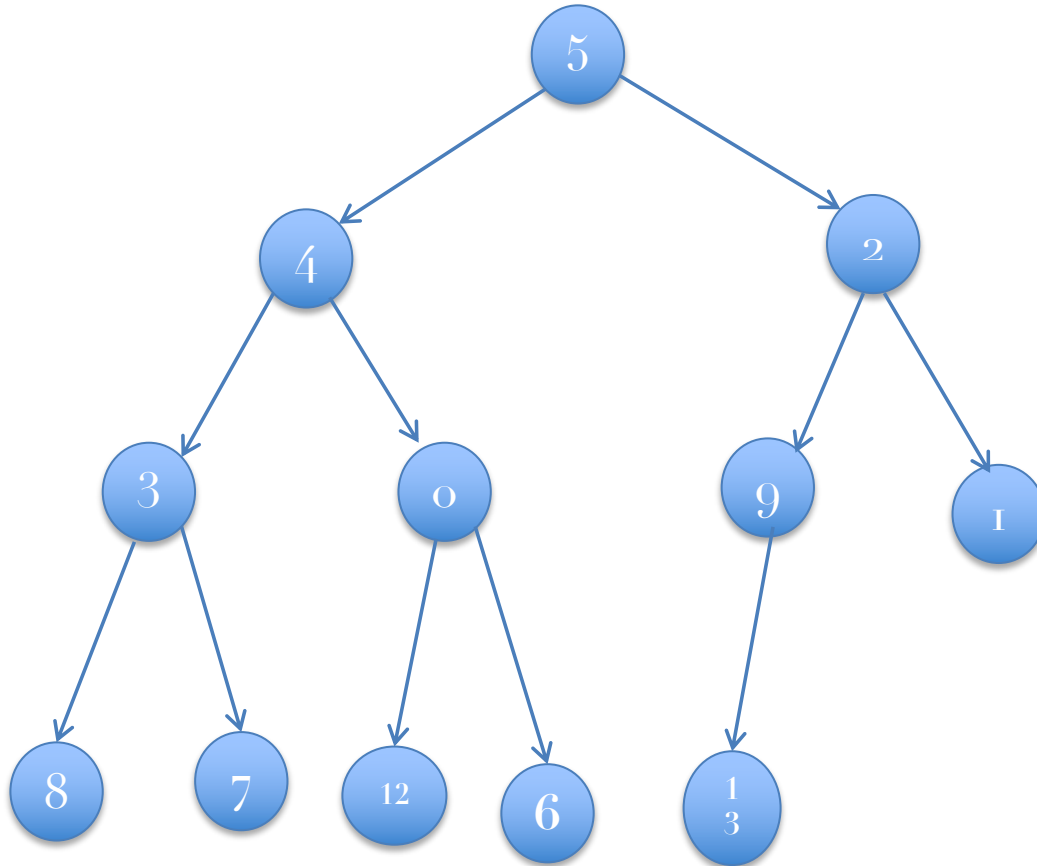
# FULL VS. COMPLETE BINARY TREE

# Full Binary Tree



Each node is either  
(1) an internal node with exactly  
two non-empty children or  
(2) a leaf

# Complete Binary Tree

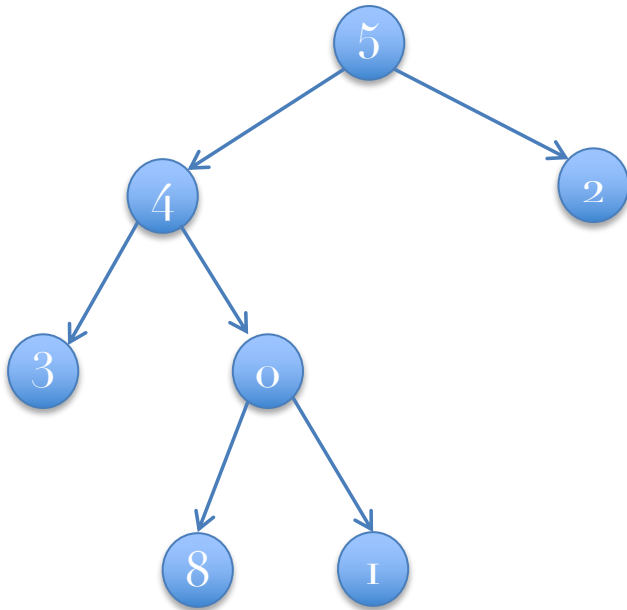


- Has a restricted shape obtained by starting at the root and filling the tree by levels from left to right.
- In the complete binary tree of height  $d$ , all levels except possibly level  $d-1$  are completely full.
- The bottom level has its nodes filled in from the left side.

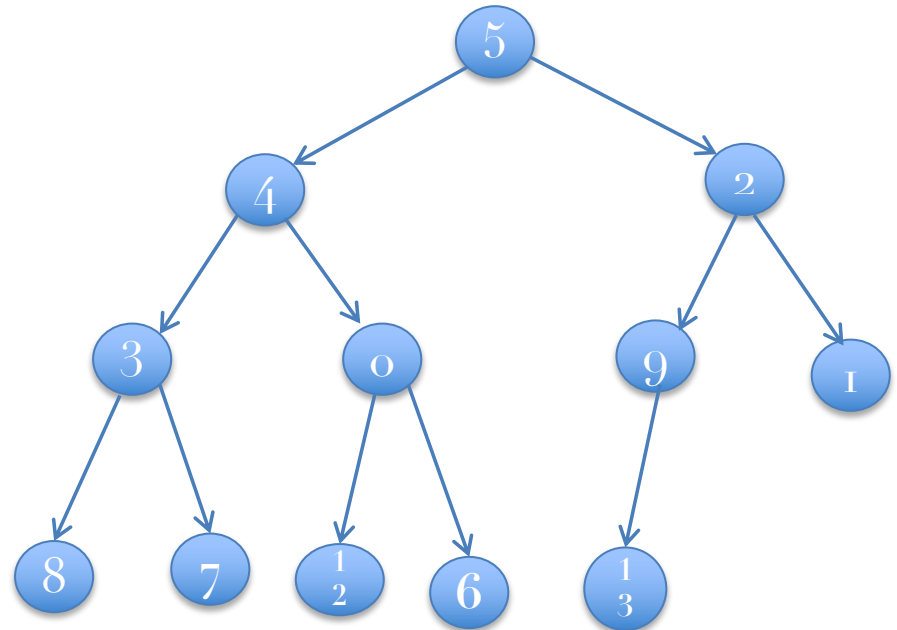


# Let's see the examples again

# Full vs. Complete Binary Tree



Full but not  
complete



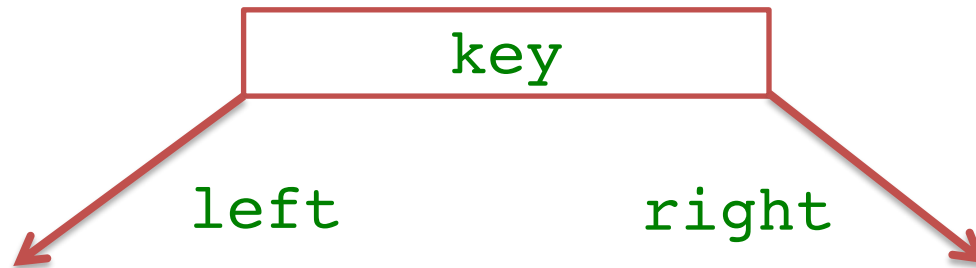
Complete  
but not full

# TREE USING JAVA

# A BTreeNode in Java

```
public class Node
{
    public int key;
    public Node left, right;

    public Node(int item)
    {
        key = item;
        left = right = null;
    }
}
```



There are many ways to traverse a binary tree

- (reverse) In order
- (reverse) Post order
- (reverse) Pre order
- Level order = breadth first

# TREE WALKS/TRAVERSALS

# Inorder Traversal

**Inorder-Traversal**(BTNode root)

- **Inorder-Traversal**(root.left)
- **Visit**(root)
- **Inorder-Traversal**(root.right)

Also called the *(left, node, right)* order

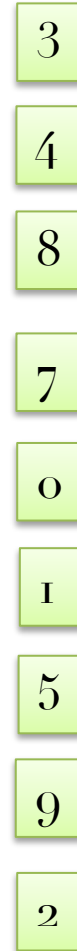
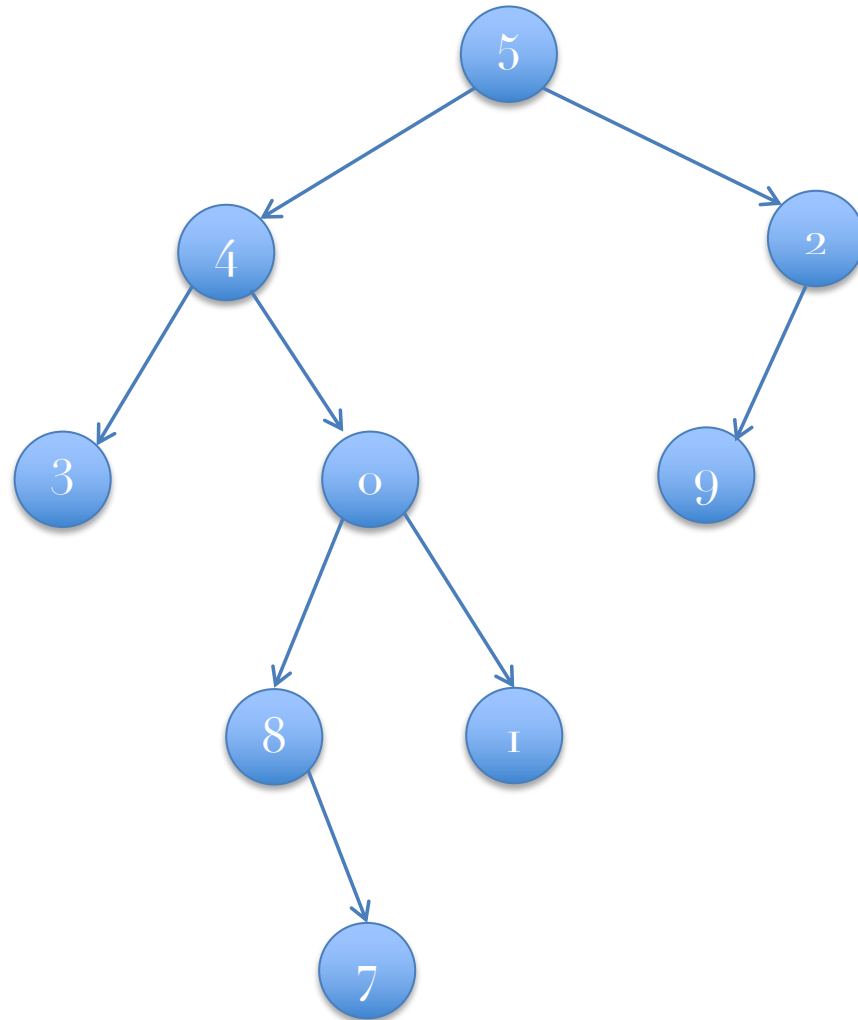
# Inorder Printing in C++

```
void inorder_print(BTNode root)
{
    if (root != null) {
        inorder_print(root.left);
        printNode(root);
        inorder_print(root.right);
    }
}
```



“Visit” the node

# In Picture





# Run Time

- Suppose “visit” takes  $O(1)$ -time, say  $c$  sec
  - $n_l$  = # of nodes on the left sub-tree
  - $n_r$  = # of nodes on the right sub-tree
  - Note:  $n - 1 = n_l + n_r$
- $T(n) = T(n_l) + T(n_r) + c$
- Induction:  $T(n) \leq cn$ , i.e.  $T(n) = O(n)$
- $$\begin{aligned} T(n) &\leq cn_l + cn_r + c \\ &= c(n-1) + c \\ &= cn \end{aligned}$$

# Reverse Inorder Traversal

- `RevInorder-Traversal(root.right)`
- `Visit(root)`
- `RevInorrdor-Traversal(root.left)`

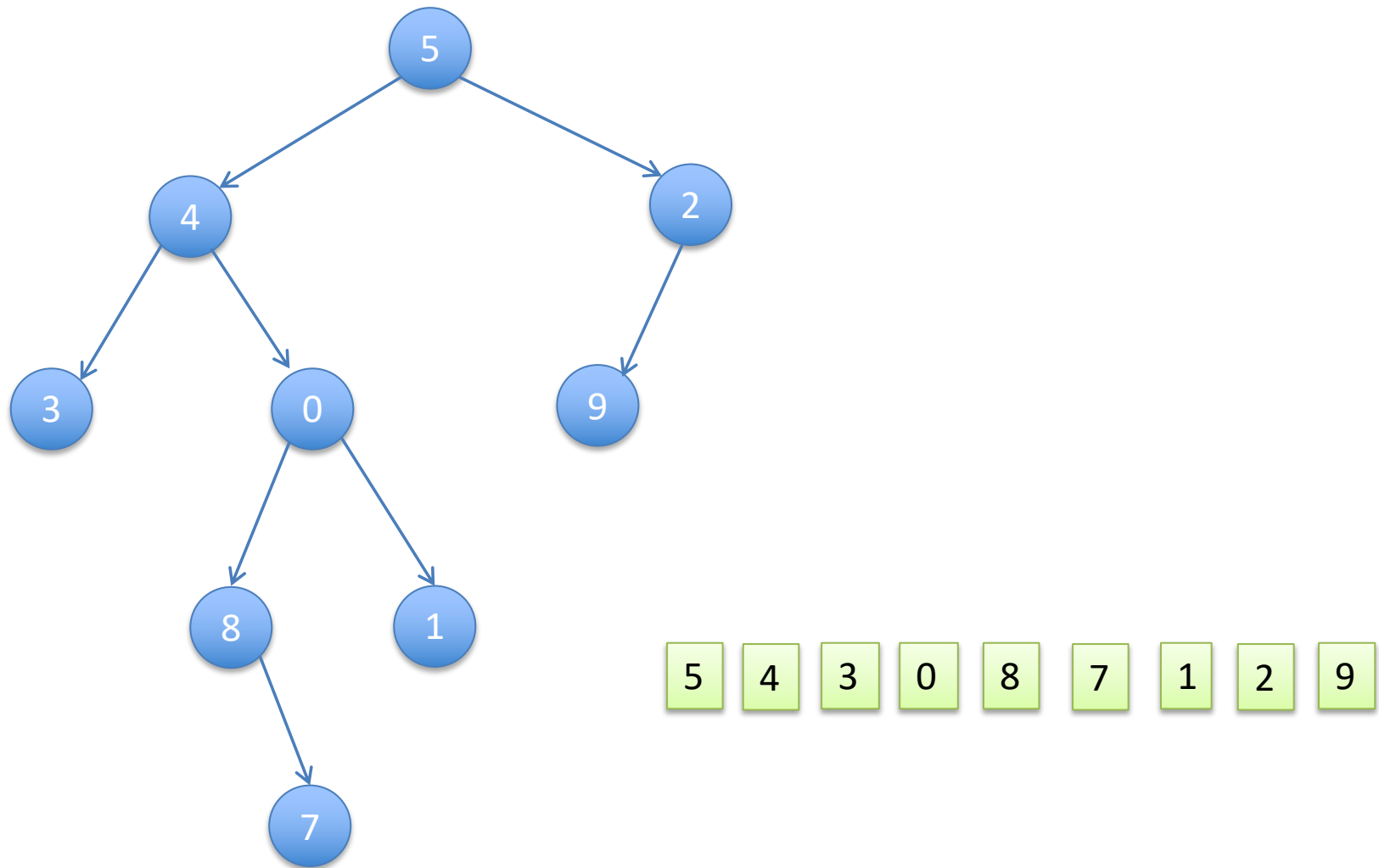
The (right, node, left) order

# The other 4 traversal orders

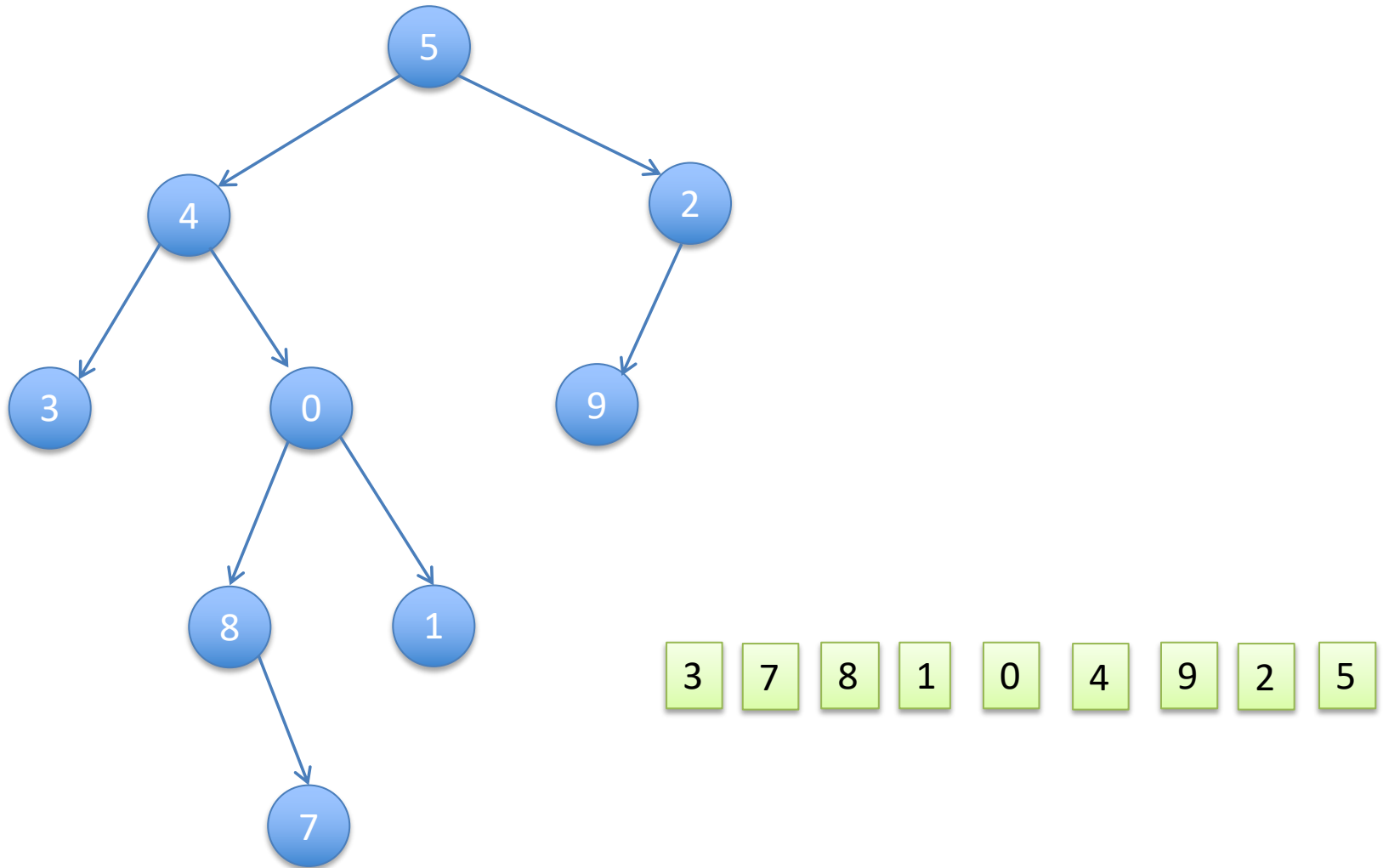
- Preorder: (node, left, right)
- Reverse preorder: (node, right, left)
- Postorder: (left, right, node)
- Reverse postorder: (right, left, node)

We'll talk about level-order later

# What is the preorder output for this tree?



# What is the postorder output for this tree?



# Questions to Ponder

```
void inorder_print(BTNode root) {  
    if (root != NULL) {  
        inorder_print(root.left);  
        printNode(root);  
        inorder_print(root.right);  
    }  
}
```

Write the above routine without the recursive calls?

Use a stack

Don't use a stack

Try this for fun! Not required for  
Exam or Quiz

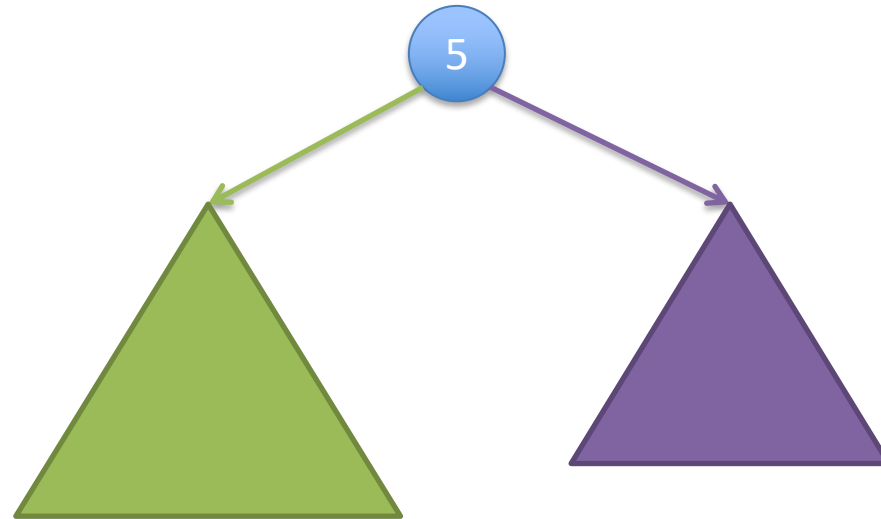
# Reconstruct the tree from inorder+postorder

Inorder

3 4 8 7 0 1 5 9 2

Preorder

5 4 3 0 8 7 1 2 9



# Questions to Ponder

- Can you reconstruct the tree given its postorder and preorder sequences?
- How about inorder and reverse postorder?
- How about other pairs of orders?
- How many trees are there which have the same in/post/pre-order sequence? (suppose keys are distinct)



# Number of trees with a given inorder sequence

Catalan numbers:  
Not required for Exam or Quiz

Catalan numbers

$$C_n = \sum_{i=1}^n C_{i-1} C_{n-i}$$

$$C_0 = 1$$

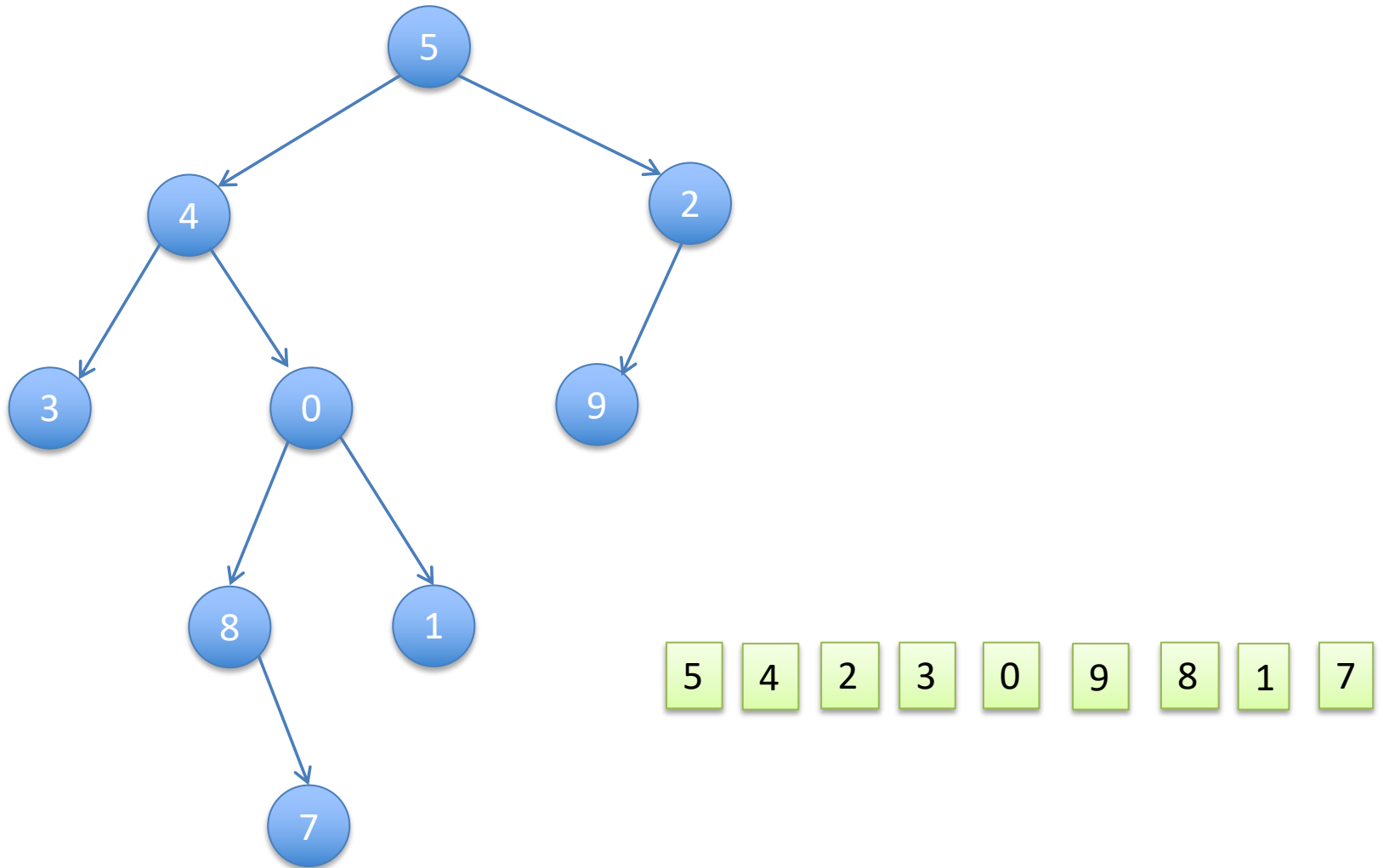
$$C_n = \frac{1}{n+1} \binom{2n}{n} \approx \frac{4^n}{n^{3/2} \sqrt{\pi}}$$

[https://en.wikipedia.org/wiki/Catalan\\_number](https://en.wikipedia.org/wiki/Catalan_number)

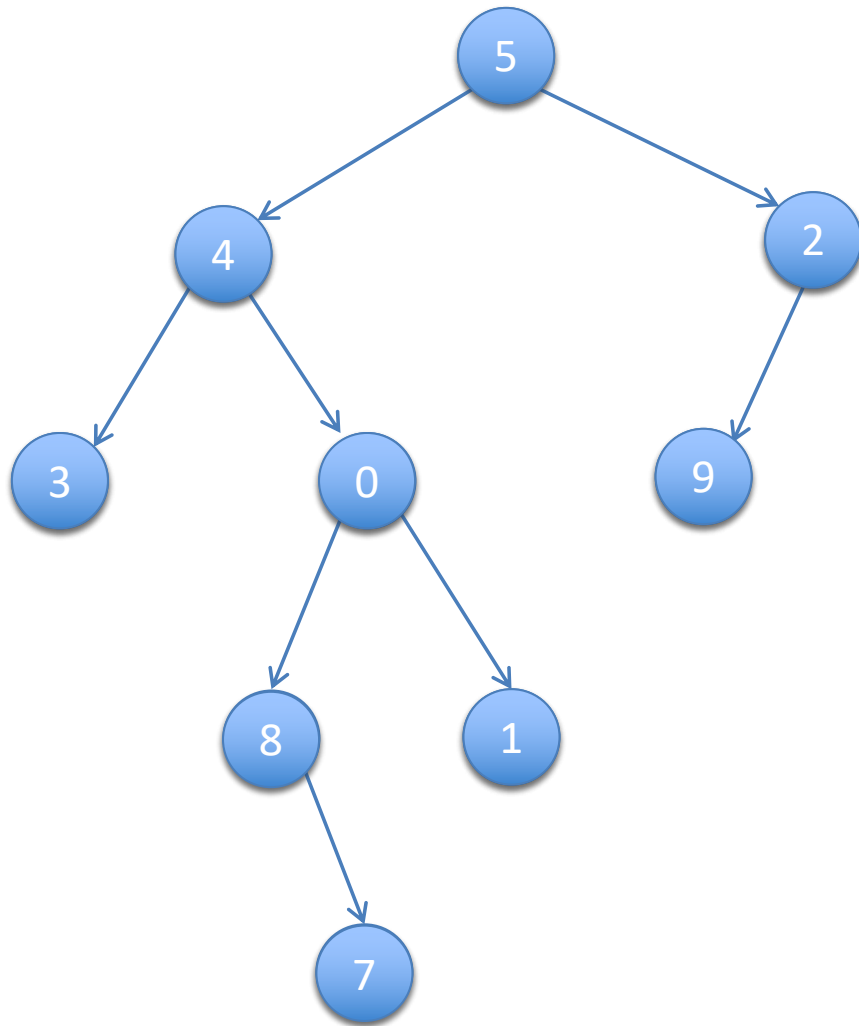
# What is a traversal order good for?

- Many things
- E.g., **Evaluate**(root) of an **expression tree**
  - If root is an operand, return the operand
  - Else
    - A = **Evaluate**(root.left)
    - B = **Evaluate**(root.right)
    - Return **A** **root.key** **B**
      - **root.key** is one of the operators
- What traversal order is the above?

# Level-Order Traversal



# How to do level-order traversal?



A (FIFO) Queue

# Level-Order Print in Java

```
void levelorder_print(BTNode root) {
```

```
// Implement
```

```
}
```

# Level-Order Print in Java

```
void printLevelOrder()
{
    Queue<BTNode> queue = new LinkedList<BTNode>();
    queue.offer(root);
    while (!queue.isEmpty())
    {
        BTNode currNode = queue.poll();
        System.out.print(currNode.getPayLoad() + " ");

        if (currNode.left != null) {
            queue.offer(currNode.left);
        }

        if (currNode.right != null) {
            queue.offer(currNode.right);
        }
    }
}
```

# What if we change the Queue into a Stack

```

void print_____Order()
{
    Stack<BTNode> stack = new Stack<BTNode>();
    stack.push(root);
    while (!stack.isEmpty())
    {
        BTNode tempNode = stack.pop();
        System.out.print(tempNode.getPayload() + " ");

        /*Enqueue left child */
        if (tempNode.left != null) {
            stack.push(tempNode.left);
        }

        /*Enqueue right child */
        if (tempNode.right != null) {
            stack.push(tempNode.right);
        }
    }
}

```

What if we change the Queue into a Stack

What traversal is this?

What traversal is this?

Reverse Preorder

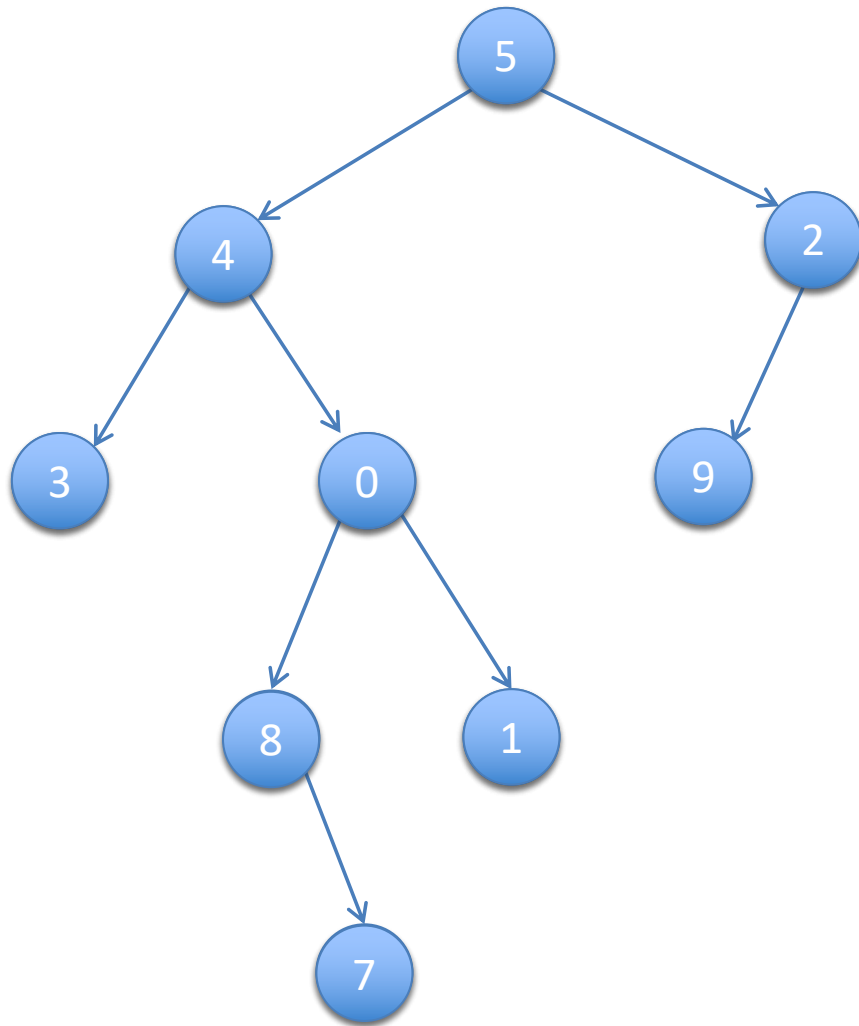
Reverse Preorder

CSC 172, Fall 2017

# **ANOTHER EXAMPLE WITHOUT RECURSION**



# How to do in-order traversal? (without recursion)

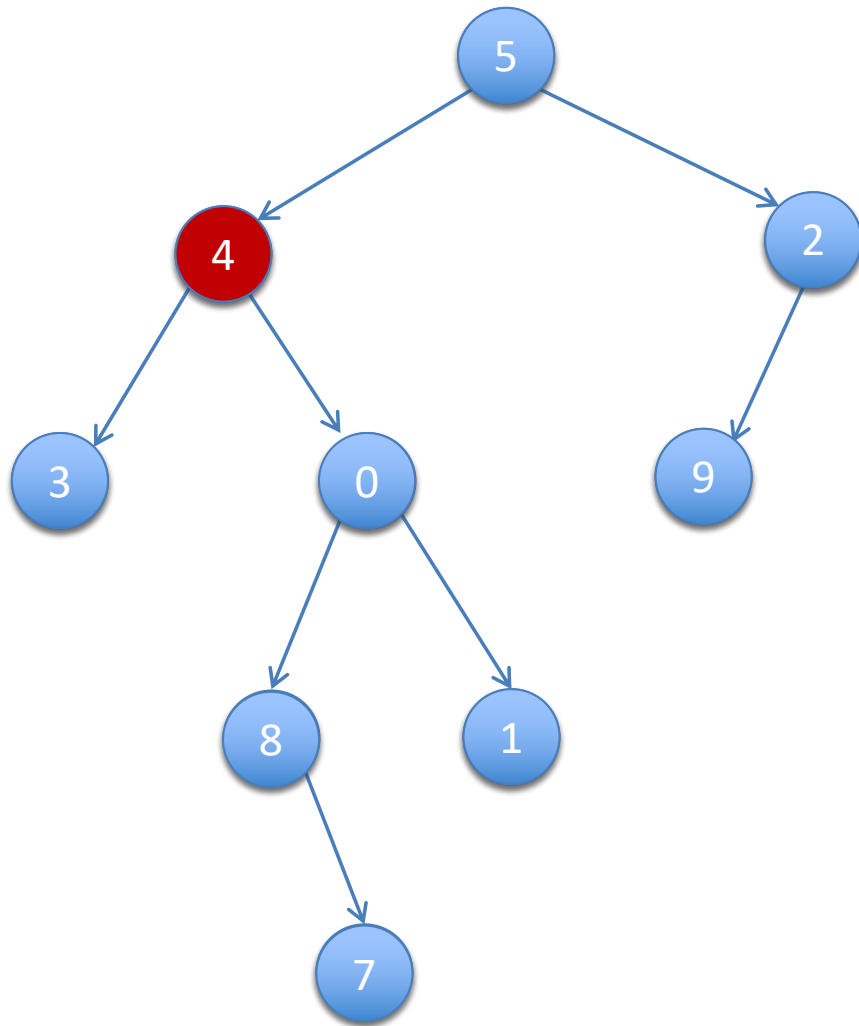


Using a Stack!  
Just use a Stack instead of a  
Queue

A Stack

Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a  
Queue

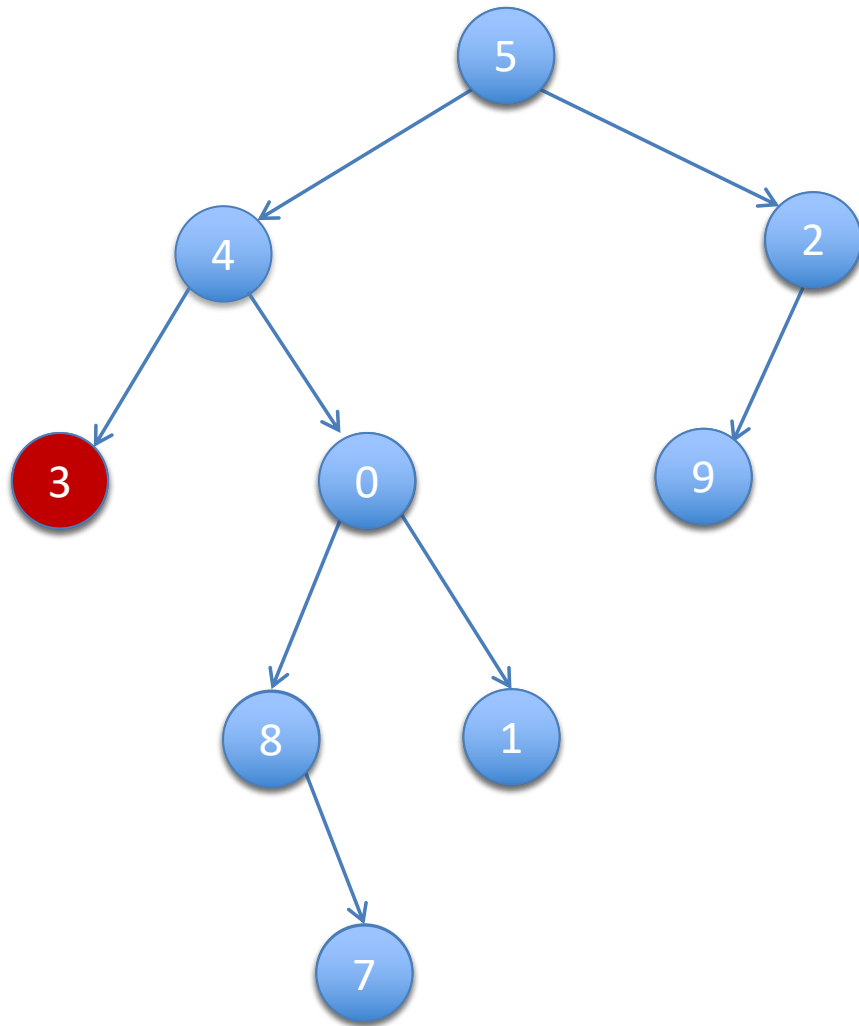


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

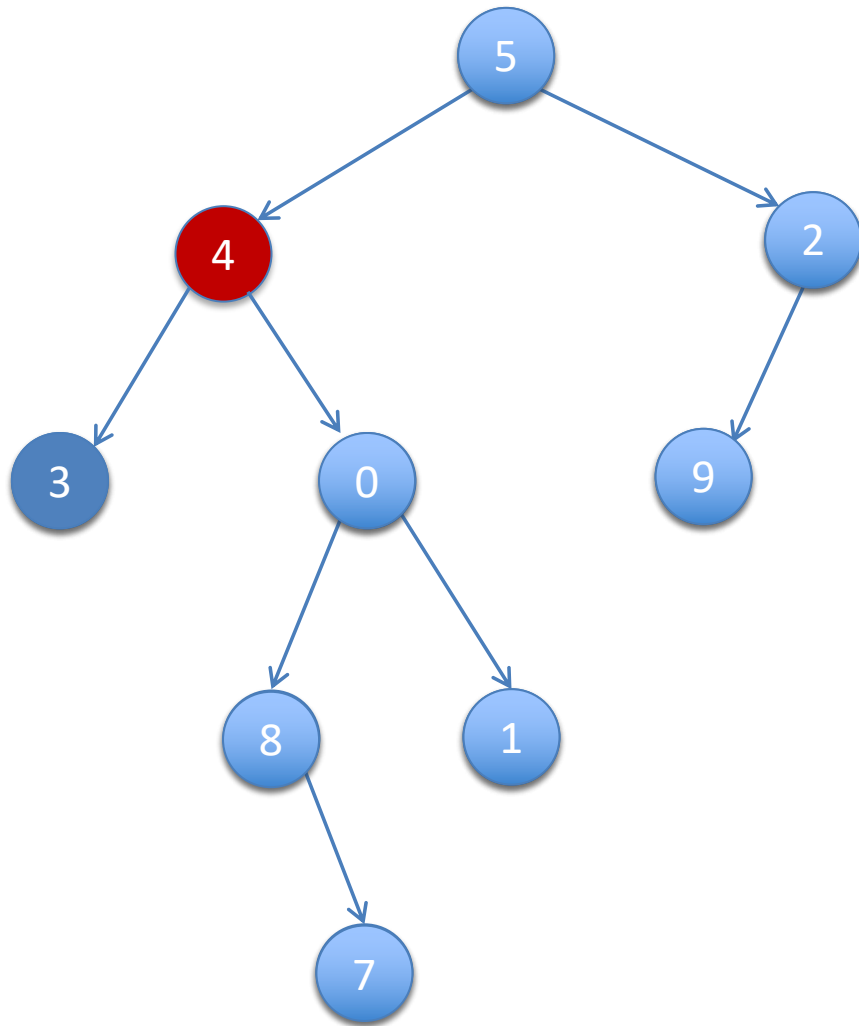


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a  
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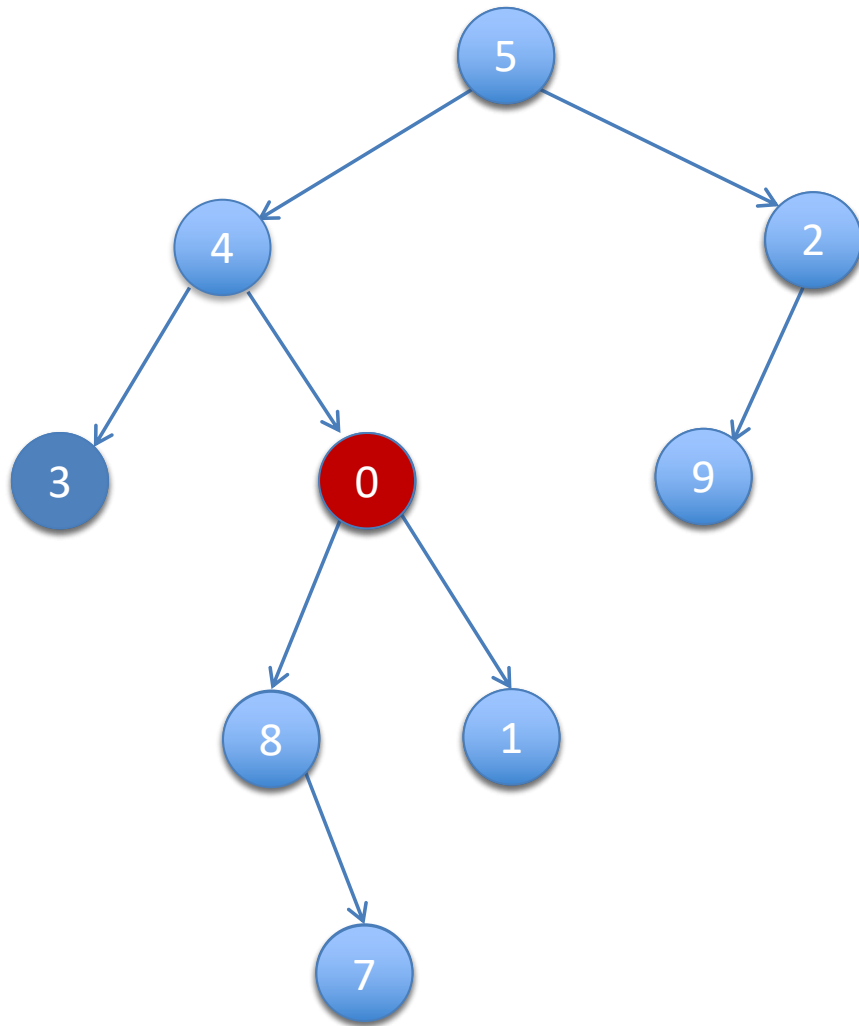


A Stack



Output

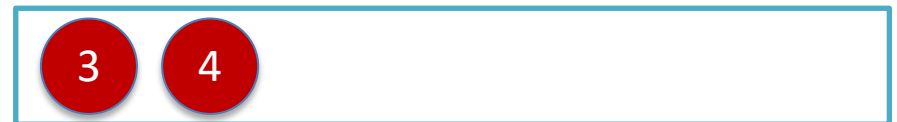
# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

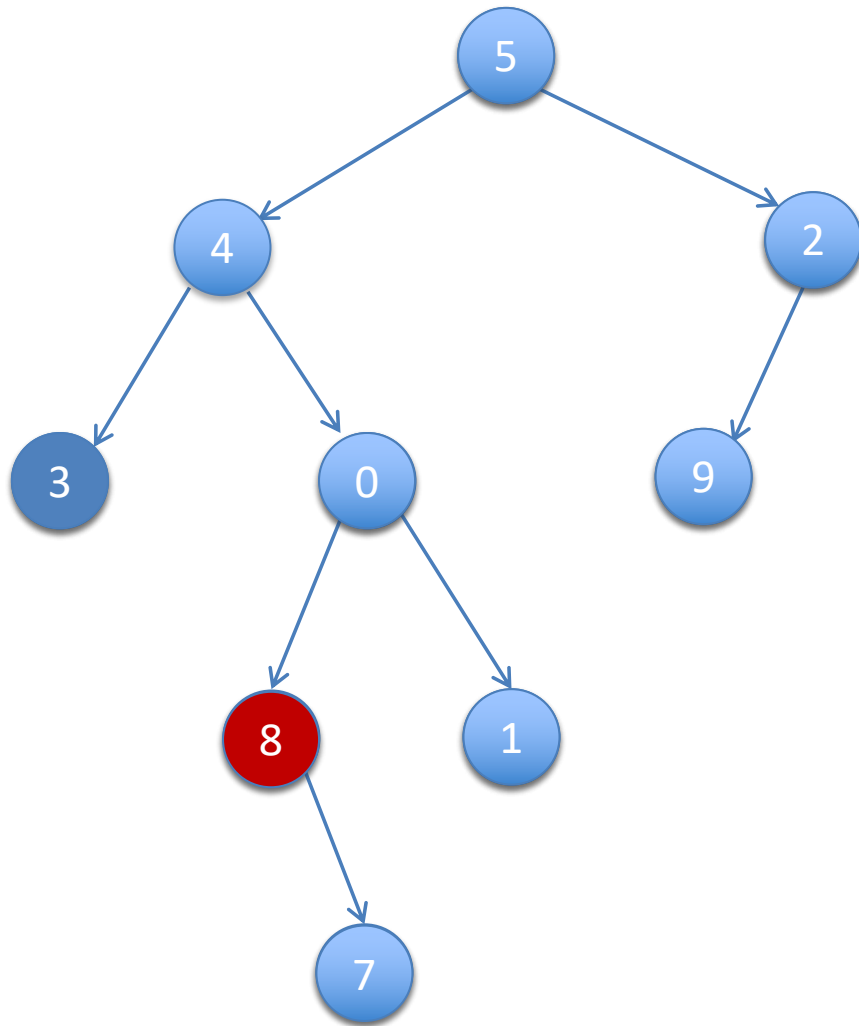


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a  
Queue

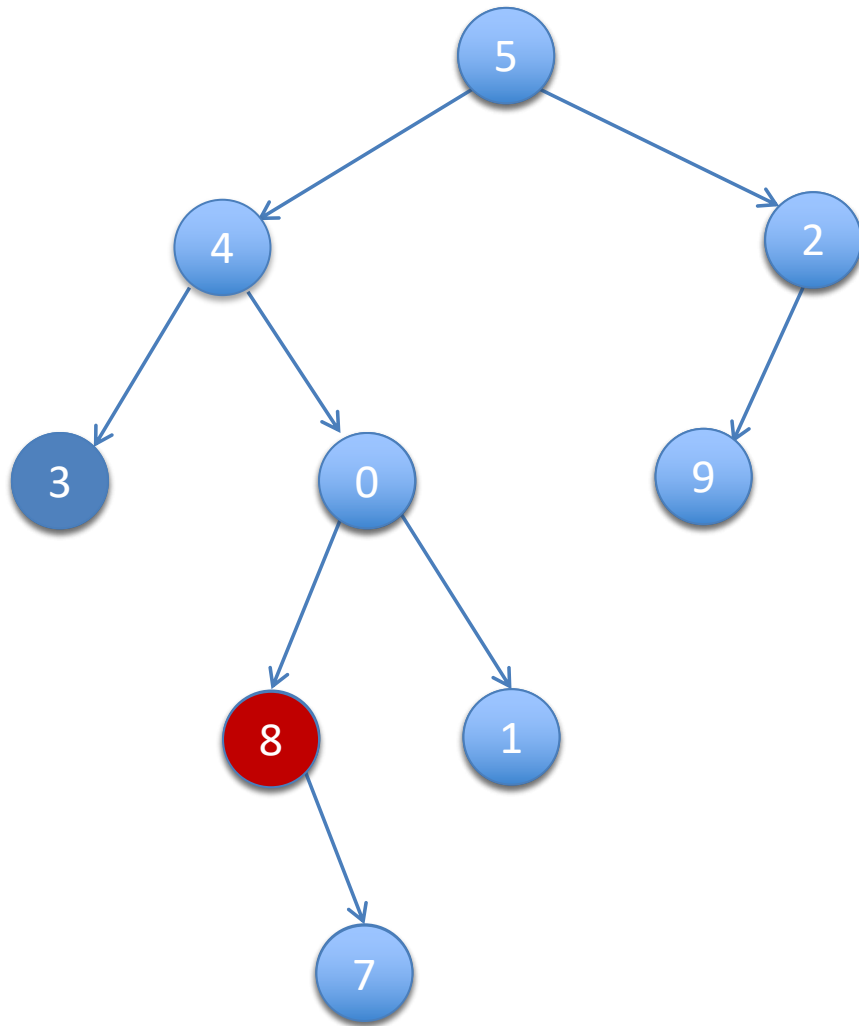


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

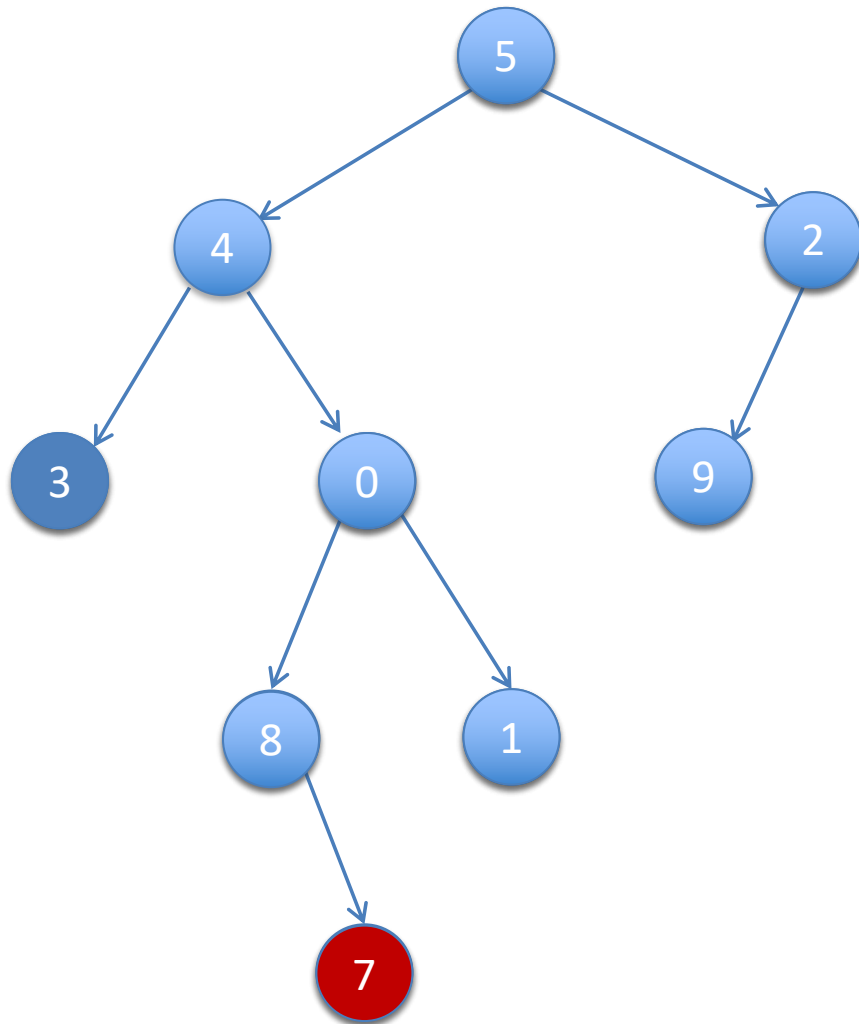


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue



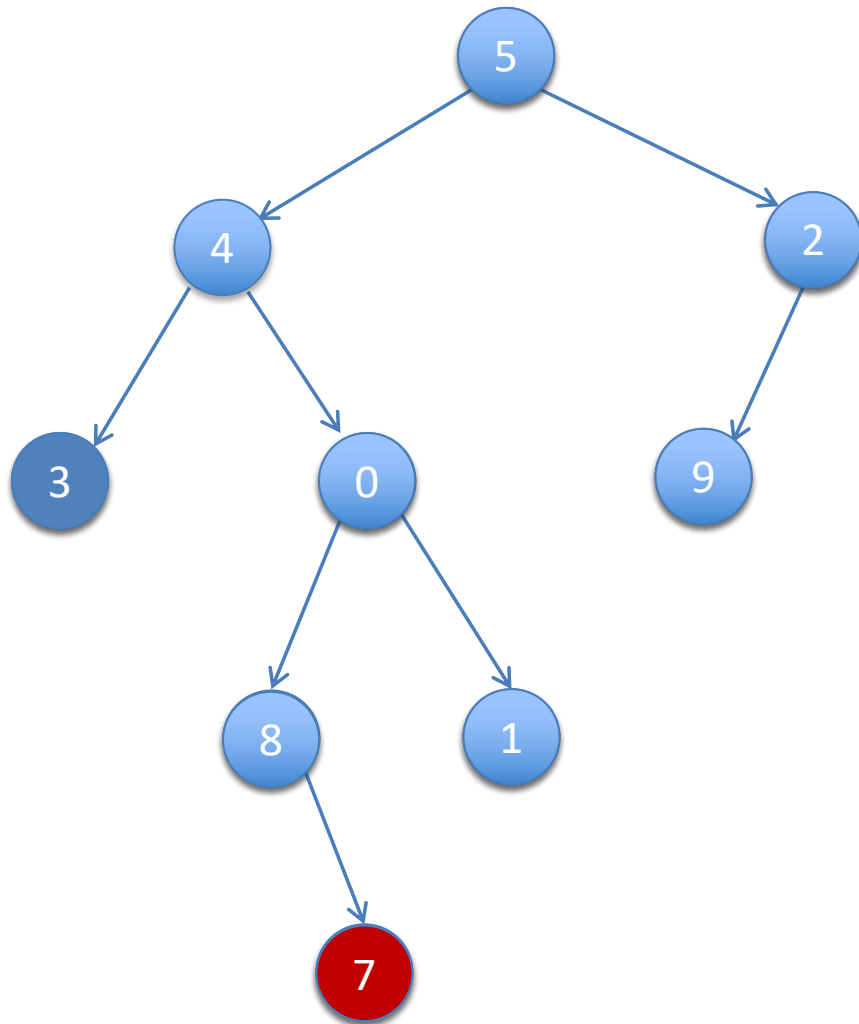
A Stack



Output



# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

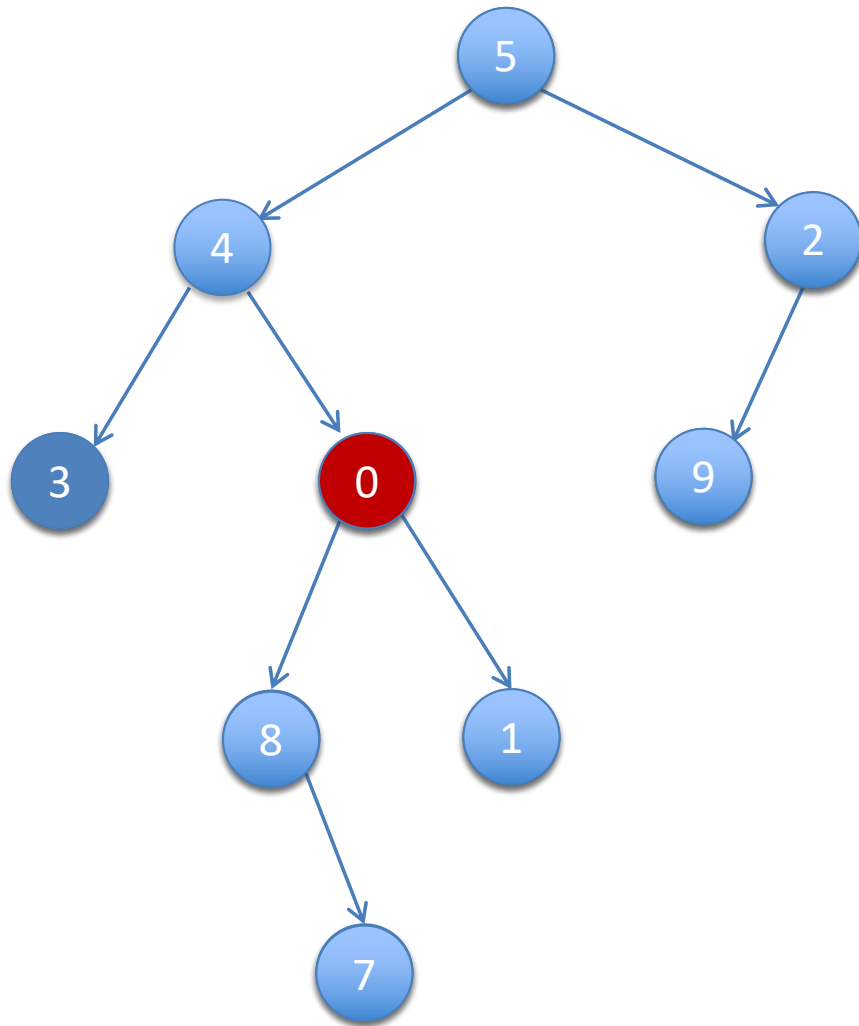


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

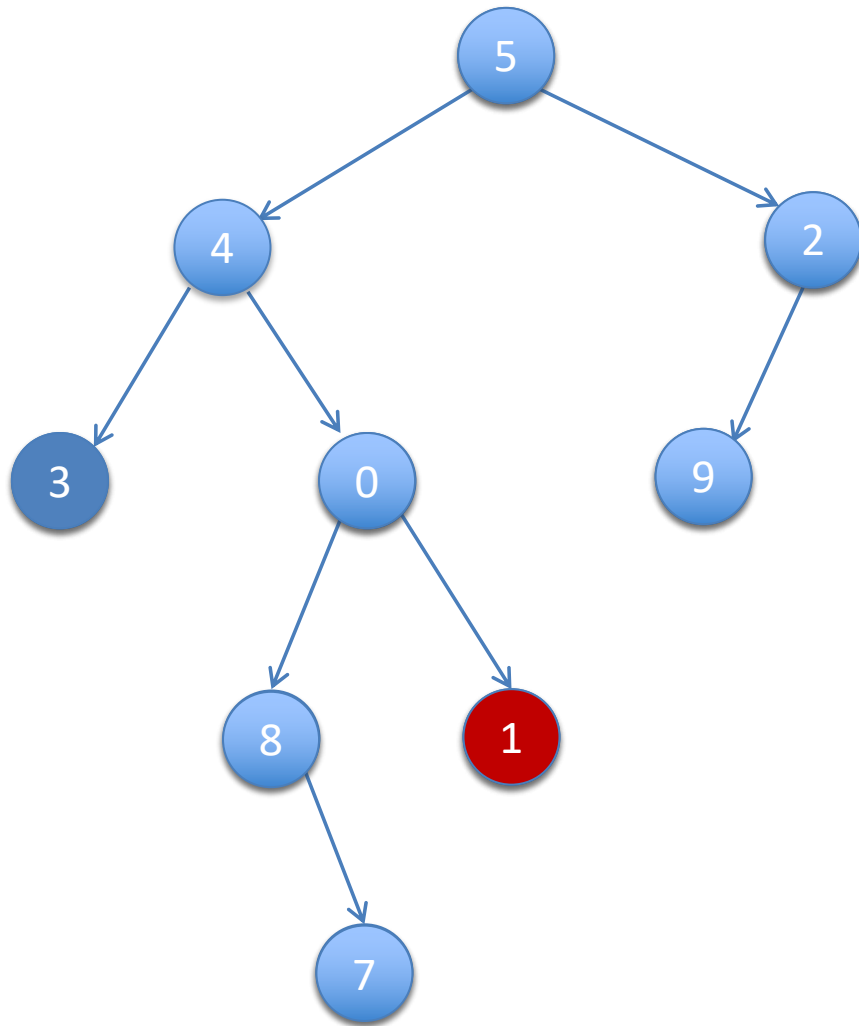


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

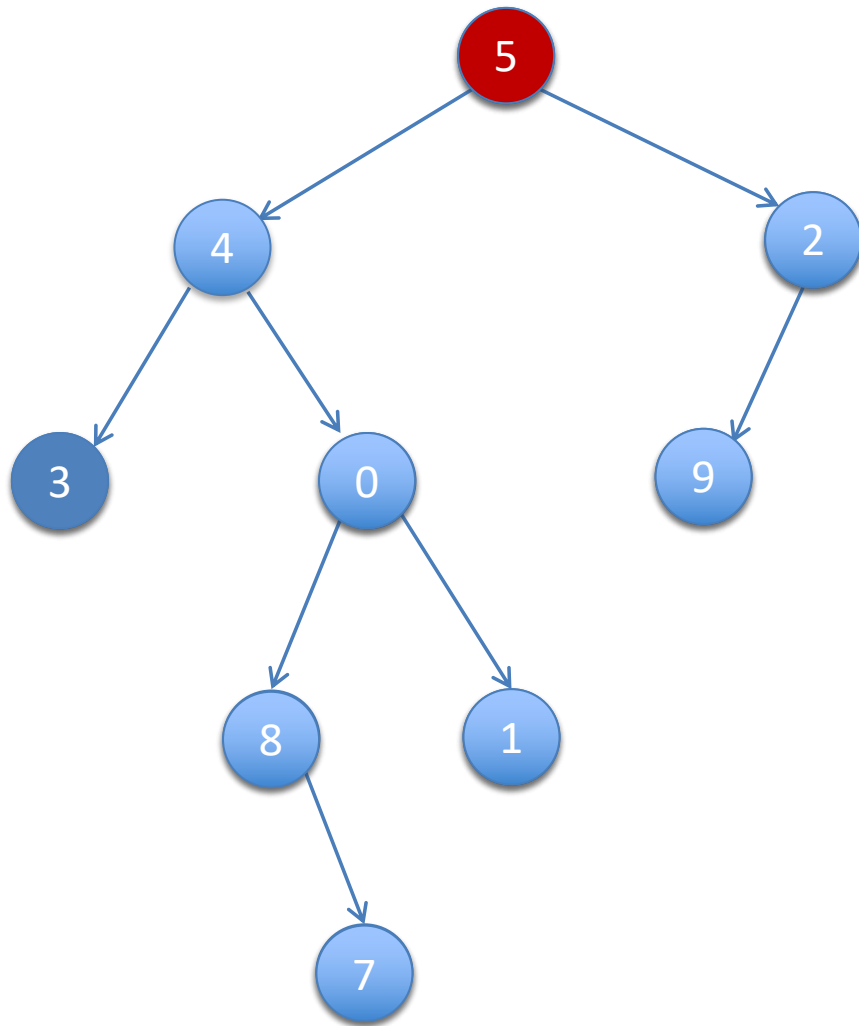


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

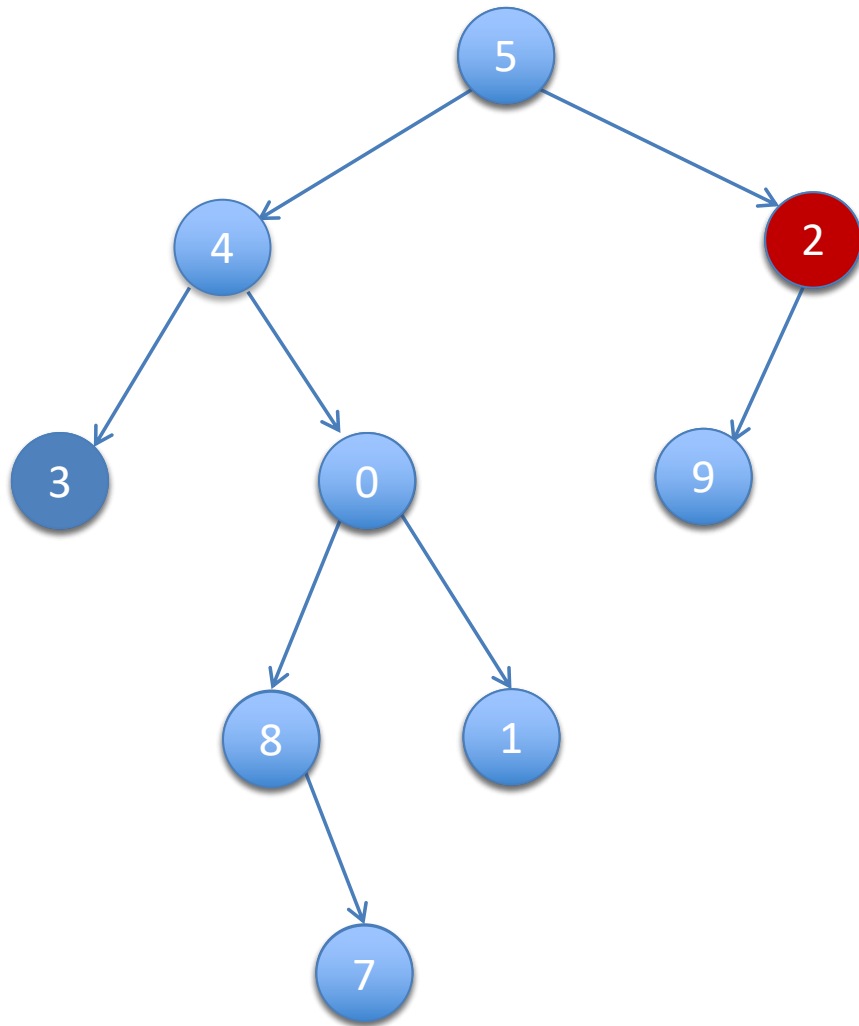


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

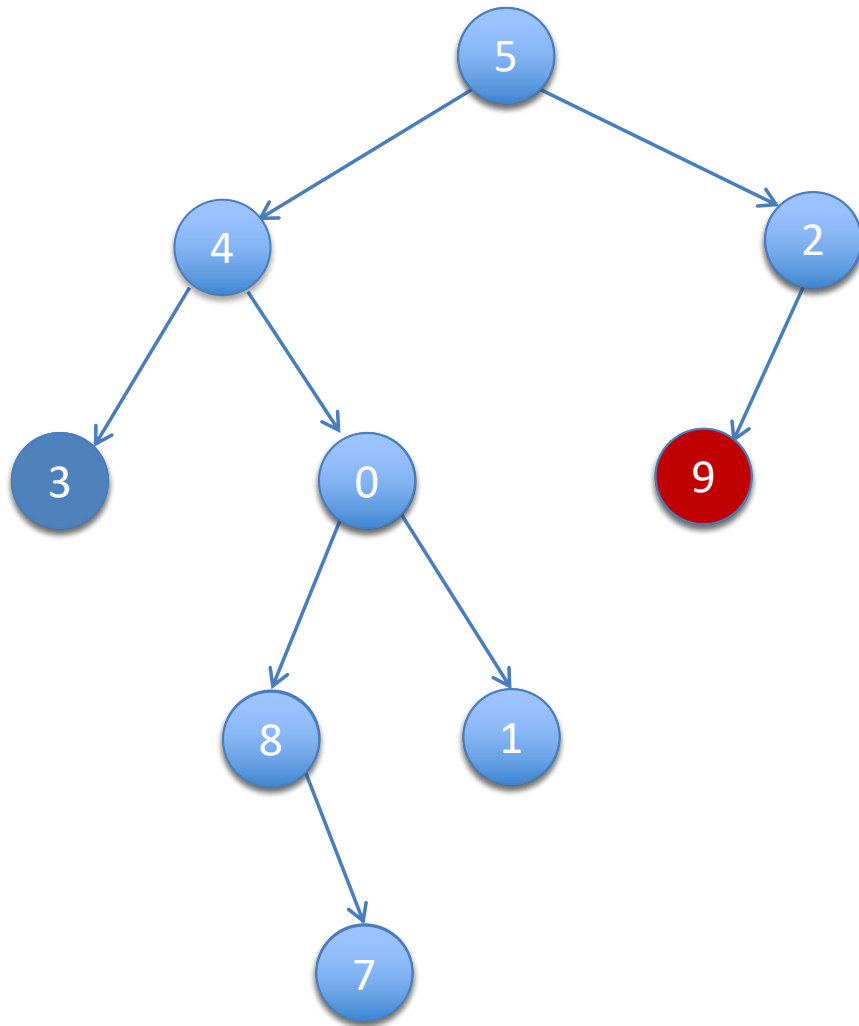


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

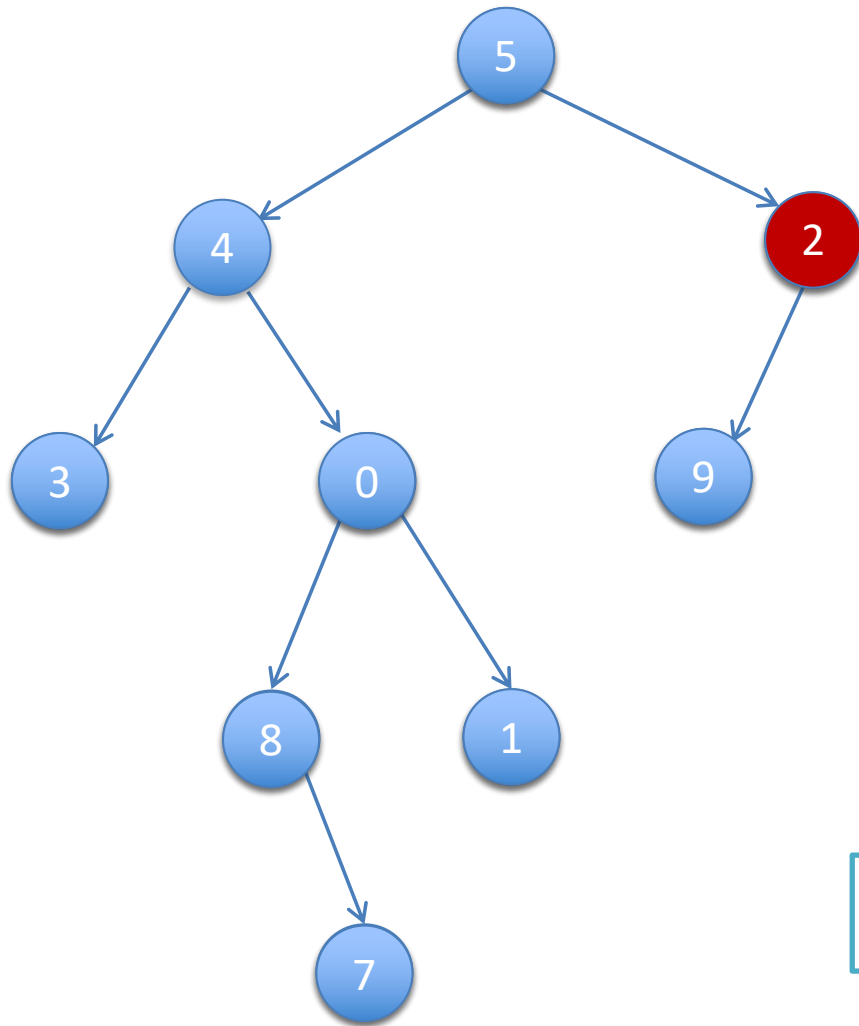


A Stack



Output

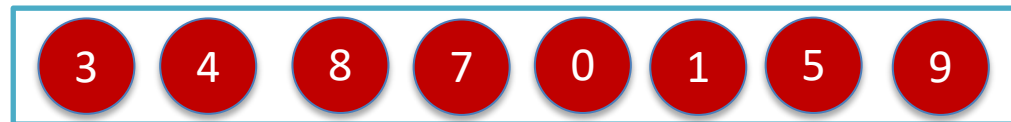
# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue

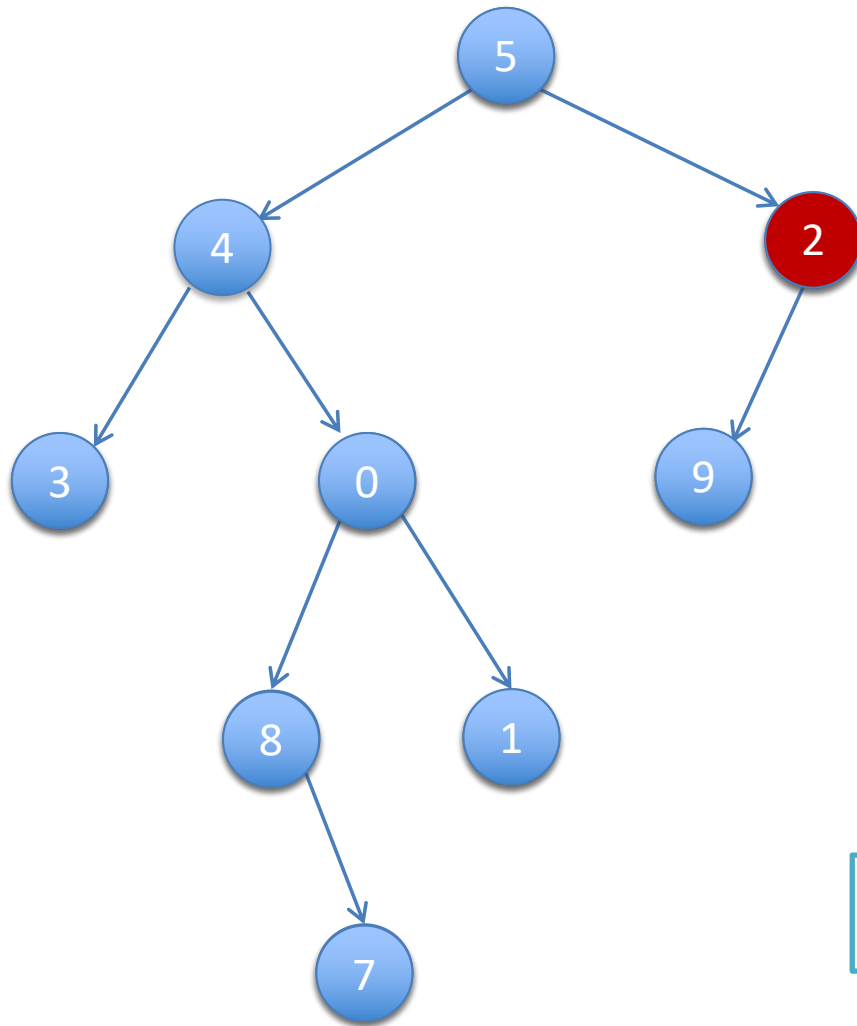


A Stack



Output

# How to do in-order traversal? (without recursion)



Using a Stack!  
Just use a Stack instead of a Queue



A Stack



Output



# Summary

We have covered:

- Definition of a tree data structure and its components
- LMC-RS vs 2-Children Tree
- Concepts of:
  - Root, internal, and leaf nodes
  - Parents, children, and siblings
  - Ancestors and descendants
  - Full vs Complete Tree
  - Tree Traversal

# Acknowledgement

- Douglas Wilhelm Harder.
  - Thanks for making an excellent set of slides for ECE 250 *Algorithms and Data Structures* course
- Prof. Hung Q. Ngo:
  - Thanks for those beautiful slides created for CSC 250 (Data Structures) course at UB.
- Many of these slides are taken from these two sources.