

CSC 252: Computer Organization

Spring 2021: Lecture 19

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Announcements

- Cache problem set: <https://www.cs.rochester.edu/courses/252/spring2021/handouts.html>
- Not to be turned in. Won't be graded.
- Assignment 4 out.

4	5	6	7	8	9	10
11	12	13	14	15	16	17
18	19	20	21	22	23	24

Today

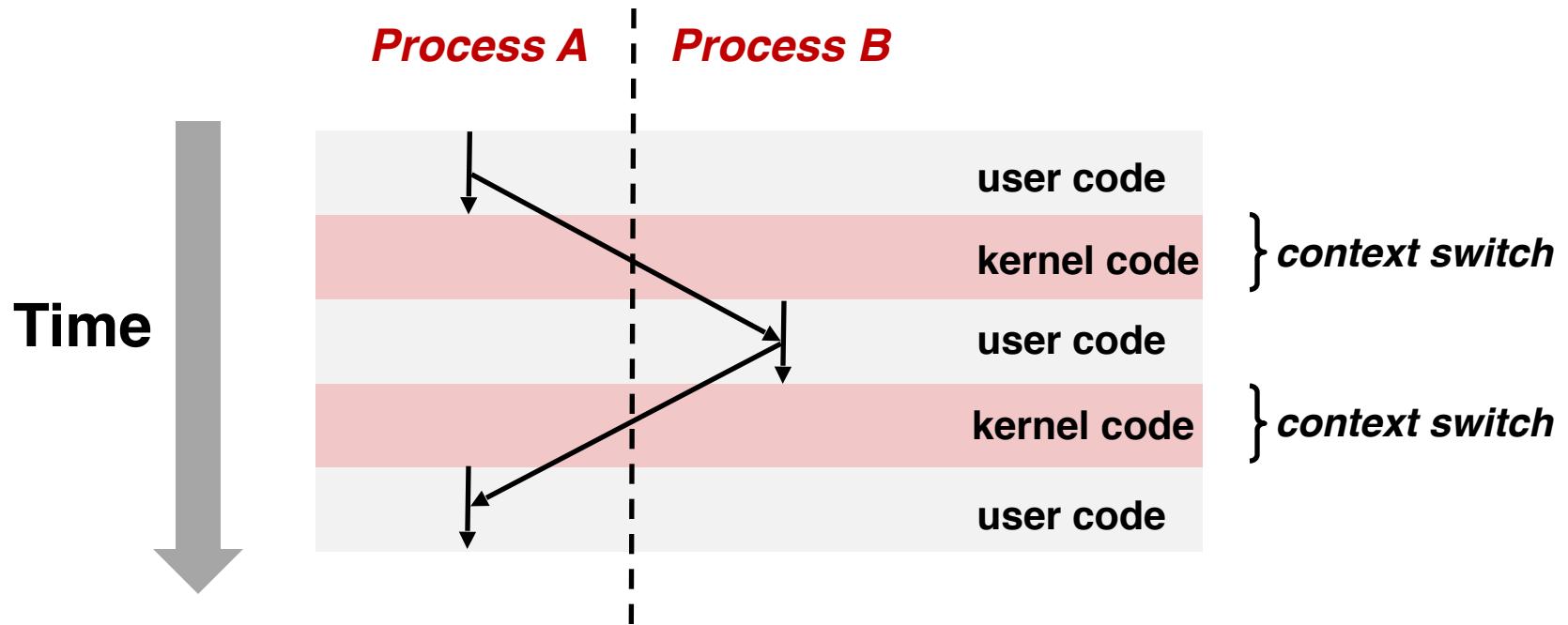
Due

Today

- Process Control
- Signals: The Way to Communicate with Processes

Context Switching

- Processes are managed by a shared chunk of memory-resident OS code called the *kernel*
 - Important: the kernel is not a separate process, but rather runs as part of some existing process.
- Control flow passes from one process to another via a *context switch*



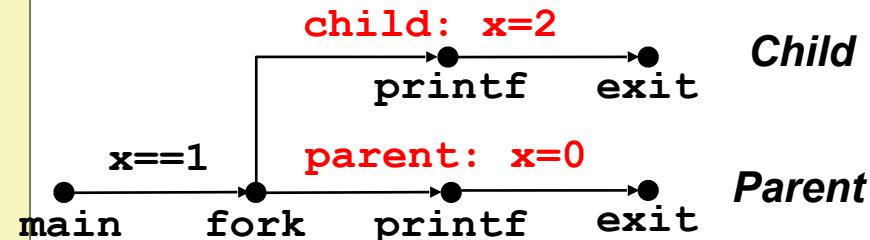
Process Graph Example

```
int main()
{
    pid_t pid;
    int x = 1;

    pid = Fork();
    if (pid == 0) { /* Child */
        printf("child : x=%d\n", ++x);
        exit(0);
    }

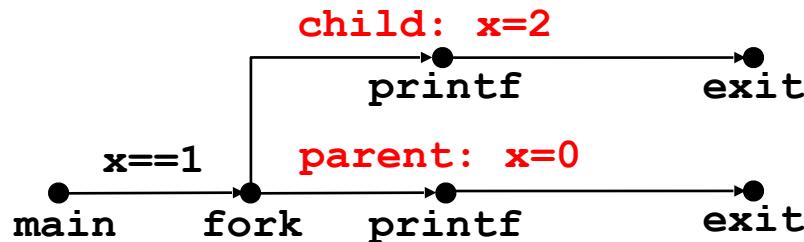
    /* Parent */
    printf("parent: x=%d\n", --x);
    exit(0);
}
```

fork.c

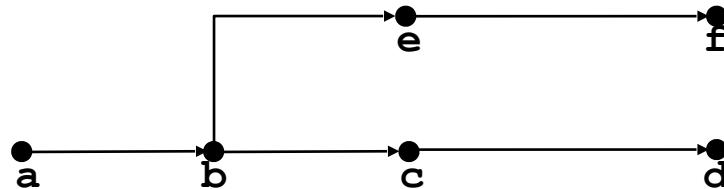


Interpreting Process Graphs

- Original graph:

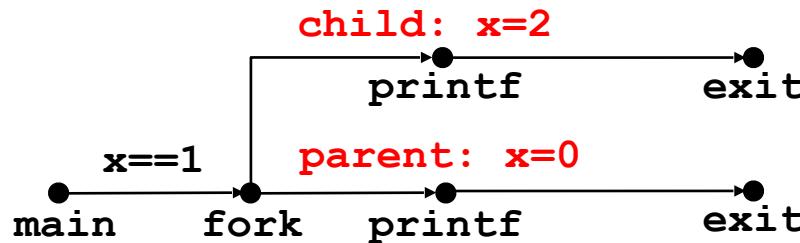


- Abstracted graph:

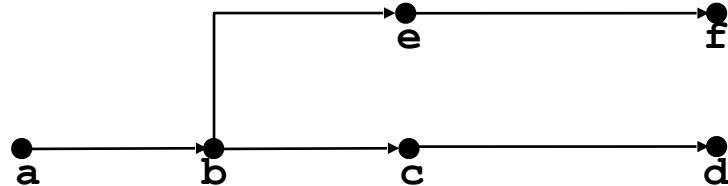


Interpreting Process Graphs

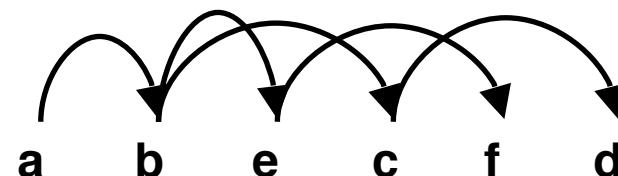
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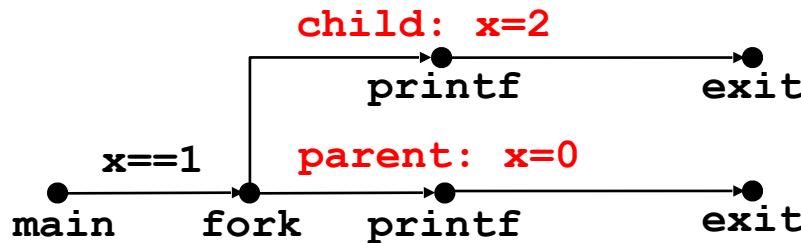


Feasible execution ordering:

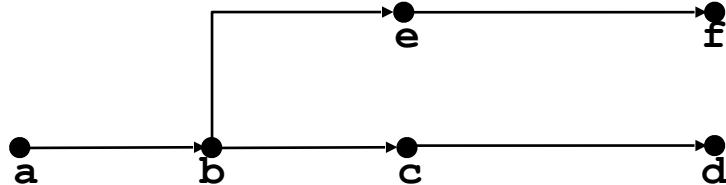


Interpreting Process Graphs

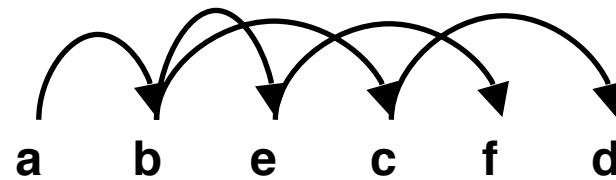
- Original graph:



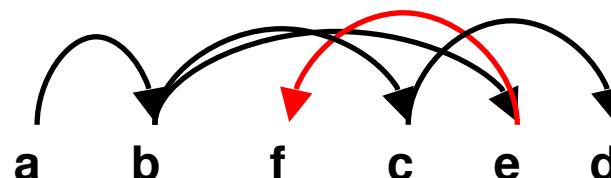
- Abstracted graph:



Feasible execution ordering:



Infeasible execution ordering:



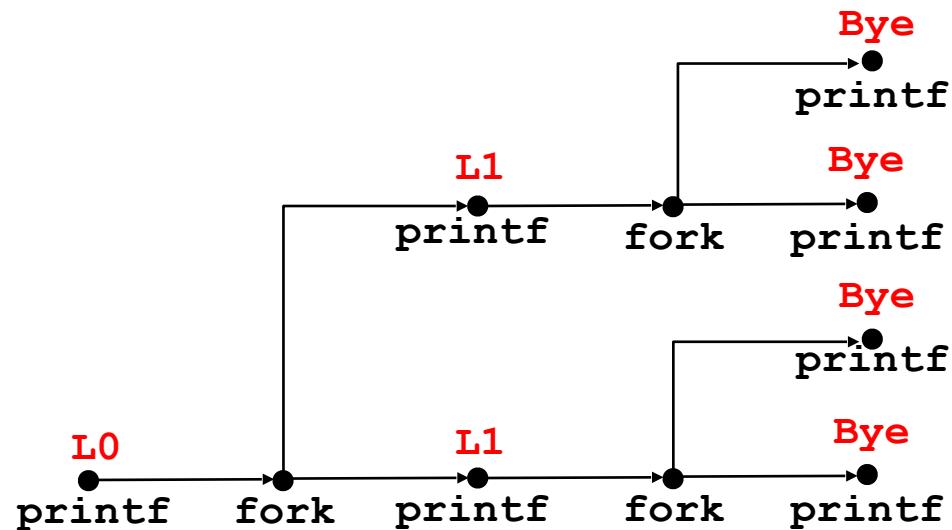
fork Example: Two consecutive forks

```
void fork2()
{
    printf("L0\n");
    fork();
    printf("L1\n");
    fork();
    printf("Bye\n");
}
```

forks.c

fork Example: Two consecutive forks

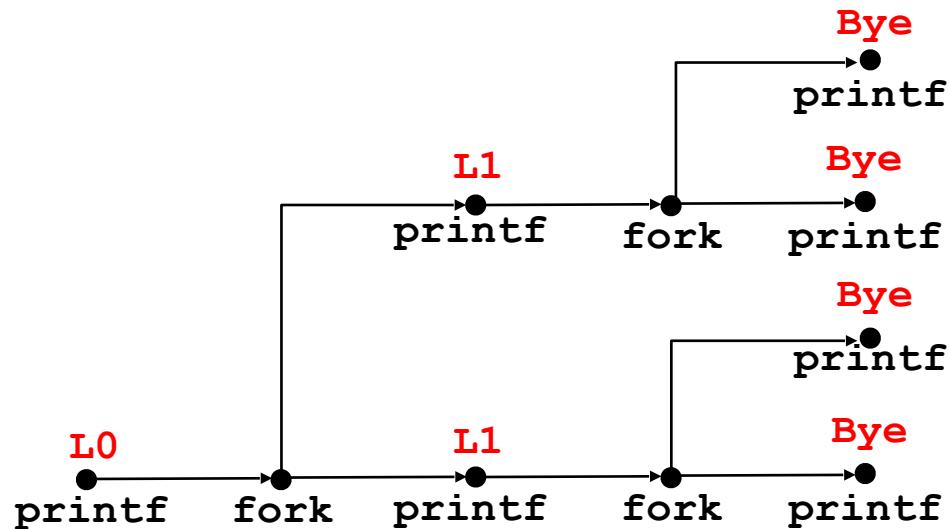
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forks.c
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forks.c



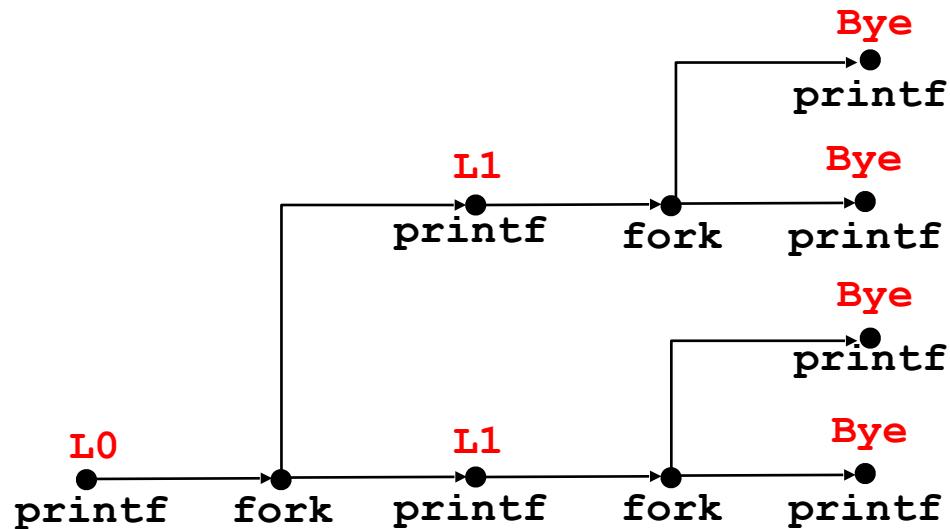
Feasible output:

L0
L1
Bye
Bye
L1
Bye
Bye

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forks.c



Feasible output:

L0
L1
Bye
Bye
L1
Bye
Bye

Infeasible output:

L0
Bye
L1
Bye
L1
Bye
Bye

fork Example: Nested forks in parent

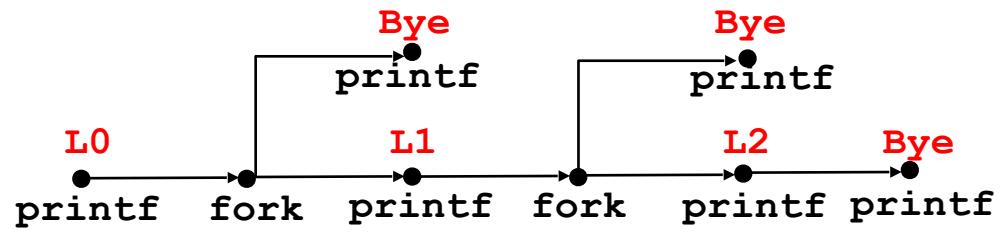
```
void fork4()
{
    printf("L0\n");
    if (fork() != 0) {
        printf("L1\n");
        if (fork() != 0) {
            printf("L2\n");
        }
    }
    printf("Bye\n");
}
```

forks.c

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        }
    }
    printf("Bye\n");
}
```

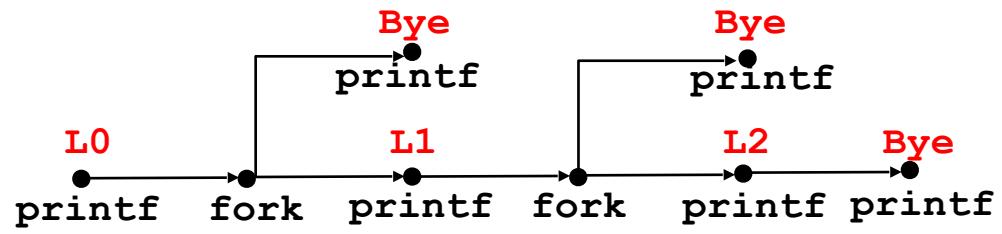
forks.c



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forks.c



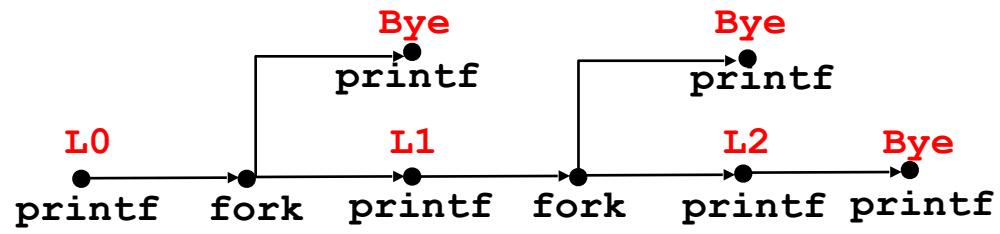
Feasible output:

L0
L1
Bye
Bye
L2
Bye

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forks.c



Feasible output:

L0
L1
Bye
Bye
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Infeasible output:

L0
Bye
L1
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L2

fork Example: Nested forks in children

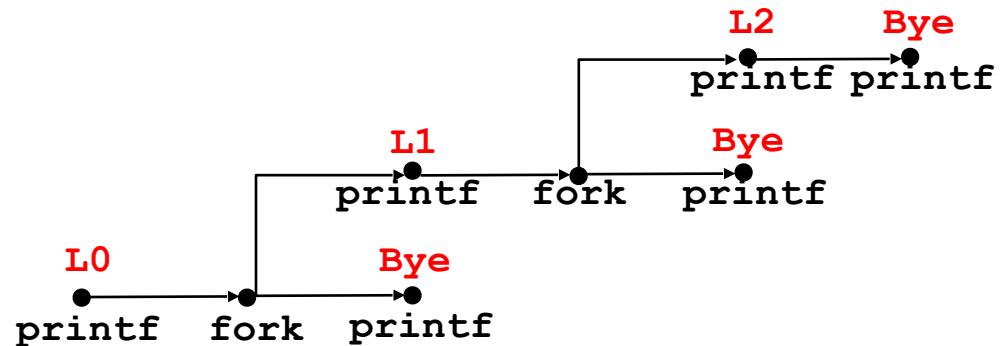
```
void fork5()
{
    printf("L0\n");
    if (fork() == 0) {
        printf("L1\n");
        if (fork() == 0) {
            printf("L2\n");
        }
    }
    printf("Bye\n");
}
```

forks.c

fork Example: Nested forks in children

```
void fork5()
{
    printf("L0\n");
    if (fork() == 0) {
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        }
    }
    printf("Bye\n");
}
```

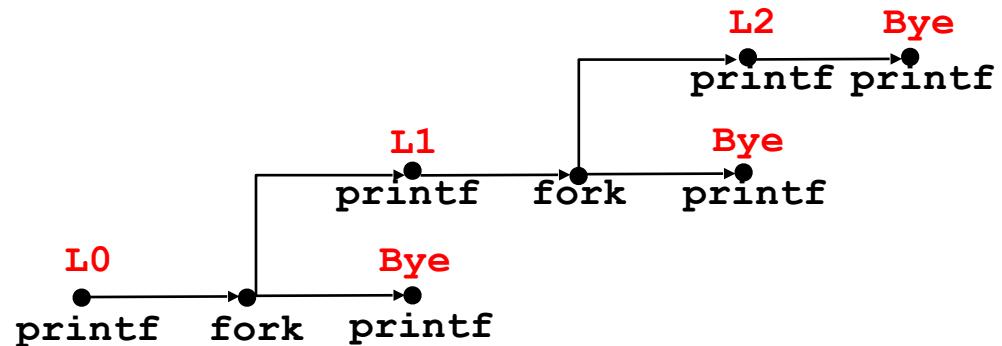
forks.c



fork Example: Nested forks in children

```
void fork5()
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    printf("L0\n");
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        }
    }
    printf("Bye\n");
}
```

forks.c



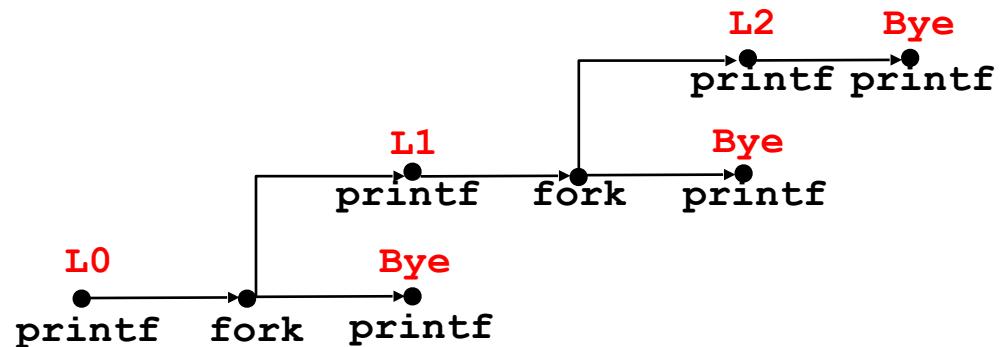
Feasible output:

L0
Bye
L1
L2
Bye
Bye

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```

forks.c



Feasible output:

L0
Bye
L1
L2
Bye
Bye

Infeasible output:

L0
Bye
L1
Bye
Bye
L2

Reaping Child Processes

- When process terminates, it still consumes system resources
 - Examples: Exit status, various OS tables
 - Called a “zombie”: Living corpse, half alive and half dead
- Reaping
 - Performed by parent on terminated child (using `wait` or `waitpid`)

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 - Parent is given exit status information
 - Kernel then deletes zombie child process
- What if parent doesn’t reap?
 - If any parent terminates without reaping a child, then the orphaned child will be reaped by `init` process (`pid == 1`)

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 - Parent is given exit status information
 - Kernel then deletes zombie child process
- What if parent doesn’t reap?
 - If any parent terminates without reaping a child, then the orphaned child will be reaped by `init` process (`pid == 1`)
 - So, only need explicit reaping in long-running processes
 - e.g., shells and servers

```
void fork7() {
    if (fork() == 0) {
        /* Child */
        printf("Terminating Child, PID = %d\n", getpid());
        exit(0);
    } else {
        printf("Running Parent, PID = %d\n", getpid());
        while (1)
            ; /* Infinite loop */
    }
}
```

forks.c

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        while (1)
            ; /* Infinite loop */
    }
}

```

forks.c

```

linux> ./forks 7 &
[1] 6639
Running Parent, PID = 6639
Terminating Child, PID = 6640
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9      00:00:00 tcsh
 6639 ttyp9      00:00:03 forks
 6640 ttyp9      00:00:00 forks <defunct>
 6641 ttyp9      00:00:00 ps
linux> kill 6639
[1]  Terminated
linux> ps
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 6585 ttyp9      00:00:00 tcsh
 6642 ttyp9      00:00:00 ps

```

```

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forks.c

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linux> kill 6639
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```

- ps shows child process as “defunct” (i.e., a zombie)



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forks.c

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```

```

linux> kill 6639
[1]  Terminated
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  PID TTY          TIME CMD
 6585 ttys000    00:00:00 tcsh
 6642 ttys000    00:00:00 ps

```

- ps shows child process as “defunct” (i.e., a zombie)
- Killing parent allows child to be reaped by init

```
void fork8()
{
    if (fork() == 0) {
        /* Child */
        printf("Running Child, PID = %d\n",
               getpid());
        while (1)
            ; /* Infinite loop */
    } else {
        printf("Terminating Parent, PID = %d\n",
               getpid());
        exit(0);
    }
}
```

forks.c

```

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    if (fork() == 0) {
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               getpid());
        while (1)
            ; /* Infinite loop */
    } else {
        printf("Terminating Parent, PID = %d\n",
               getpid());
        exit(0);
    }
}

```

forks.c

```

linux> ./forks 8
Terminating Parent, PID = 6675
Running Child, PID = 6676
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9      00:00:00 tcsh
 6676 ttyp9      00:00:06 forks
 6677 ttyp9      00:00:00 ps
linux> kill 6676
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9      00:00:00 tcsh
 6678 ttyp9      00:00:00 ps

```

```

void fork8()
{
    if (fork() == 0) {
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        printf("Running Child, PID = %d\n",
               getpid());
        while (1)
            ; /* Infinite loop */
    } else {
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        exit(0);
    }
}

```

forks.c

```

linux> ./forks 8
Terminating Parent, PID = 6675
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 6676 ttyp9      00:00:06 forks
 6677 ttyp9      00:00:00 ps
linux> kill 6676
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9      00:00:00 tcsh
 6678 ttyp9      00:00:00 ps

```

- Child process still active even though parent has terminated. Can't be reaped since it's still running!
- Must kill child explicitly, or else will keep running indefinitely

wait: Synchronizing with Children

```
void fork9() {
    int child_status;

    if (fork() == 0) {
        printf("HC: hello from child\n");
        exit(0);
    } else {
        printf("HP: hello from parent\n");
        wait(&child_status);
        printf("CT: child has terminated\n");
    }
    printf("Bye\n");
}
```

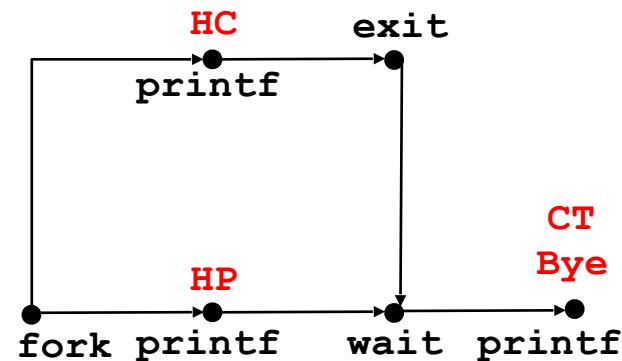
forks.c

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    if (fork() == 0) {
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forks.c

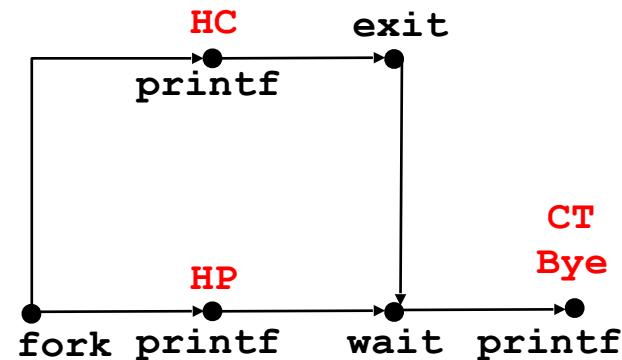


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    printf("Bye\n");
}
```

forks.c



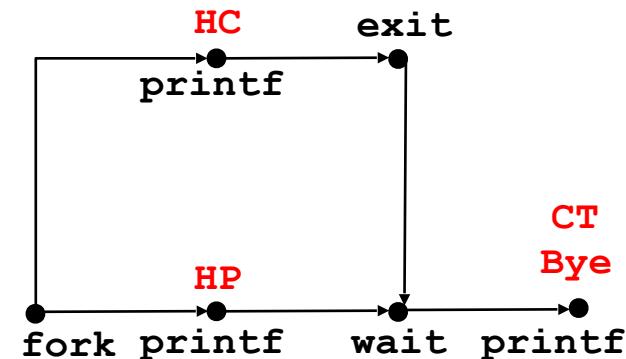
Feasible output:
HC
HP
CT
Bye

wait: Synchronizing with Children

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    if (fork() == 0) {
        printf("HC: hello from child\n");
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        printf("HP: hello from parent\n");
        wait(&child_status);
        printf("CT: child has terminated\n");
    }
    printf("Bye\n");
}
```

forks.c



Feasible output:
HC
HP
CT
Bye

Infeasible output:
HP
CT
Bye
HC

wait: Synchronizing with Children

- Parent reaps a child by calling the `wait` function
- `int wait(int *child_status)`
 - Suspends current process until one of its children terminates
 - Return value is the `pid` of the child process that terminated
 - If `child_status != NULL`, then the integer it points to will be set to a value that indicates reason the child terminated and the exit status:
 - Checked using macros defined in `wait.h`
 - `WIFEXITED`, `WEXITSTATUS`, `WIFSIGNALED`,
`WTERMSIG`, `WIFSTOPPED`, `WSTOPSIG`,
`WIFCONTINUED`
 - See textbook for details

Another wait Example

- If multiple children completed, will take in arbitrary order
- Can use macros WIFEXITED and WEXITSTATUS to get information about exit status

```
void fork10() {
    int i, child_status;

    for (i = 0; i < N; i++)
        if (fork() == 0) {
            exit(100+i); /* Child */
        }
    for (i = 0; i < N; i++) { /* Parent */
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child_status))
            printf("Child %d terminated with exit status %d\n",
                   wpid, WEXITSTATUS(child_status));
        else
            printf("Child %d terminate abnormally\n", wpid);
    }
}
```

forks.c

waitpid: Waiting for a Specific Process

- pid_t waitpid(pid_t pid, int &status, int options)
 - Suspends current process until specific process terminates
 - Various options (see textbook)

```
void fork11() {
    pid_t pid[N];
    int i;
    int child_status;

    for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0)
            exit(100+i); /* Child */
    for (i = N-1; i >= 0; i--) {
        pid_t wpid = waitpid(pid[i], &child_status, 0);
        if (WIFEXITED(child_status))
            printf("Child %d terminated with exit status %d\n",
                   wpid, WEXITSTATUS(child_status));
        else
            printf("Child %d terminate abnormally\n", wpid);
    }
}
```

forks.c

execve: Loading and Running Programs

Executes “/bin/ls -lt /usr/include” in child process using current environment:

```
char *myargv[] = {"/bin/ls", "-lt", "/usr/include"};
char *environ[] = {"USER=droh", "PWD=/usr/droh"};
```



```
if ((pid = Fork()) == 0) { /* Child runs program */
    if (execve(myargv[0], myargv, environ) < 0) {
        printf("%s: Command not found.\n", myargv[0]);
        exit(1);
    }
}
```

execve: Loading and Running Programs

- `int execve(char *filename, char *argv[], char *envp[])`

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- Loads and runs in the current process:
 - Executable file `filename`
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 - By convention `argv[0]==filename`
 - Environment variable list `envp`
 - “name=value” strings (e.g., `USER=droh`)

execve: Loading and Running Programs

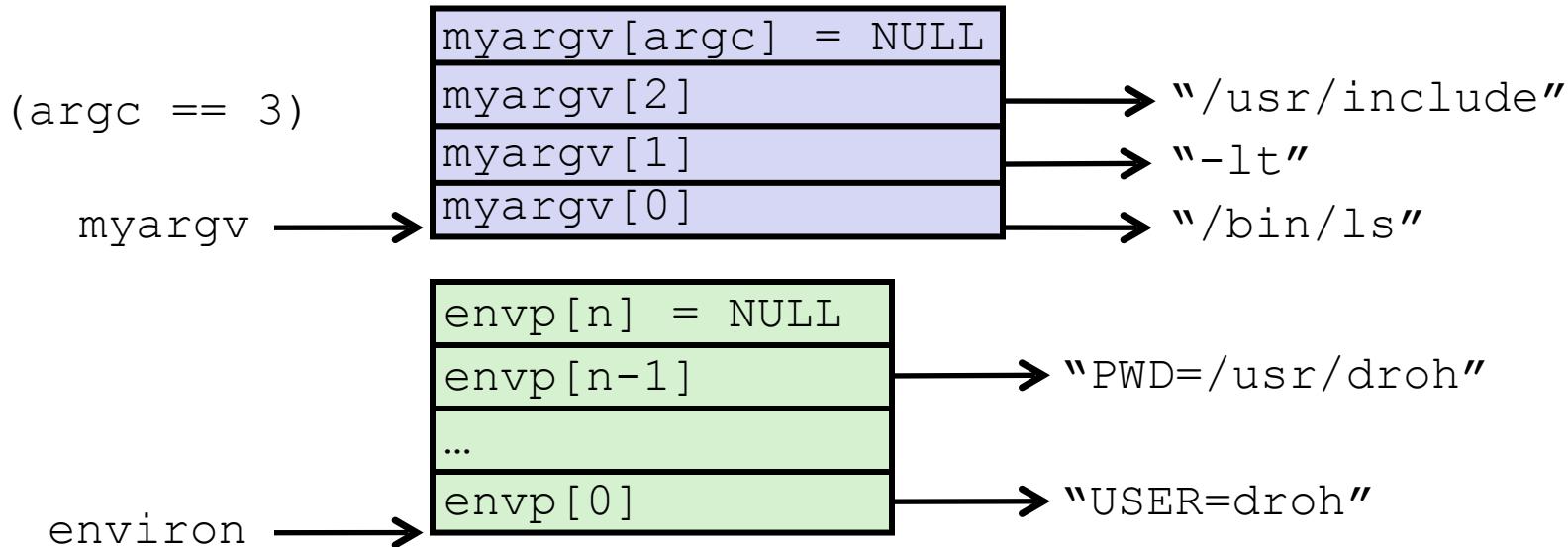
- `int execve(char *filename, char *argv[], char *envp[])`
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- Overwrites code, data, and stack
 - Retains PID, open files and signal context

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 - “name=value” strings (e.g., `USER=droh`)
- Overwrites code, data, and stack
 - Retains PID, open files and signal context
- Called **once** and **never** returns
 - ...except if there is an error

execve Example

Executes “/bin/ls -lt /usr/include” in child process using current environment:



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if ((pid = Fork()) == 0) { /* Child runs program */
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        exit(1);
    }
}
```

Summary

- Processes
 - At any given time, system has multiple active processes
 - Only one can execute at a time on a single core, though
 - Each process appears to have total control of processor + private memory space
- Spawning processes
 - Call `fork`
 - One call, two returns
- Process completion
 - Call `exit`
 - One call, no return
- Reaping and waiting for processes
 - Call `wait` or `waitpid`
- Loading and running programs
 - Call `execve` (or variant)
 - One call, (normally) no return

Today

- Process Control
- Signals: The Way to Communicate with Processes

Signals

- A **signal** is a small message that notifies a process that an event of some type has occurred in the system
 - Sent from the **OS kernel**
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<i>ID Name</i>	<i>Default Action</i>	<i>Corresponding Event</i>
2 SIGINT	Terminate	User typed ctrl-c
9 SIGKILL	Terminate	Kill program (cannot override or ignore)
11 SIGSEGV	Terminate	Segmentation violation
14 SIGALRM	Terminate	Timer signal
17 SIGCHLD	Ignore	Child stopped or terminated

Signal Concepts: Sending a Signal

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 - Another process has invoked the `kill` system call to explicitly request the kernel to send a signal to the destination process.
 - Note: `kill` doesn't mean you are going to kill the target process. It is just a system call that allows you to send signals. Of course the signal you send could be SIGKILL.

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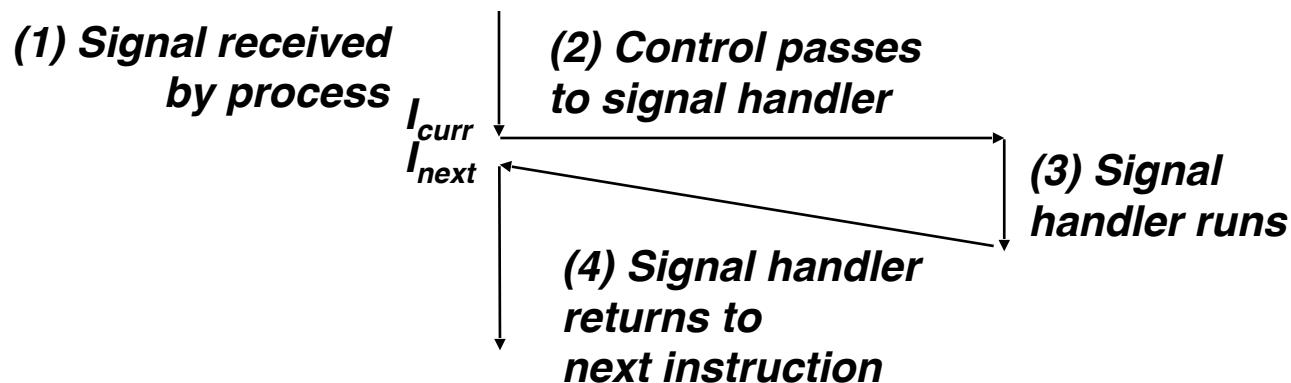
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Sending Signals with /bin/kill Program

- /bin/kill program sends arbitrary signal to a process
- Examples

- **/bin/kill -9 24818**

Send SIGKILL to process 24818

- /bin/kill itself doesn't kill the process. 9 is the ID for the SIGKILL signal, which terminates the process

```
linux> ./forks 16
Child1: pid=24818 pgrp=24817
Child2: pid=24819 pgrp=24817
```

```
linux> ps
 PID TTY          TIME CMD
24788 pts/2        00:00:00 tcsh
24818 pts/2        00:00:02 forks
24819 pts/2        00:00:02 forks
24820 pts/2        00:00:00 ps
```

Sending Signals with /bin/kill Program

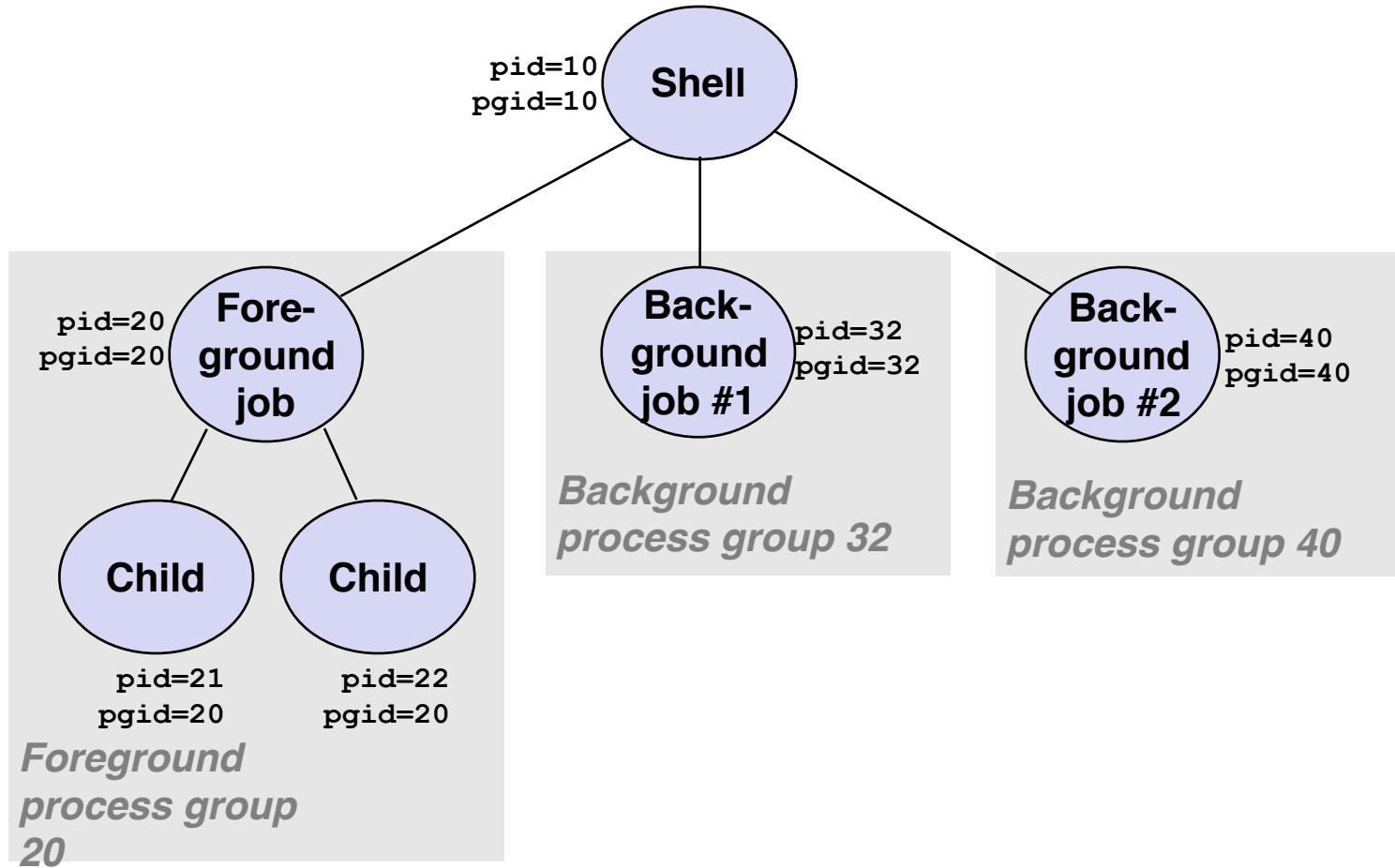
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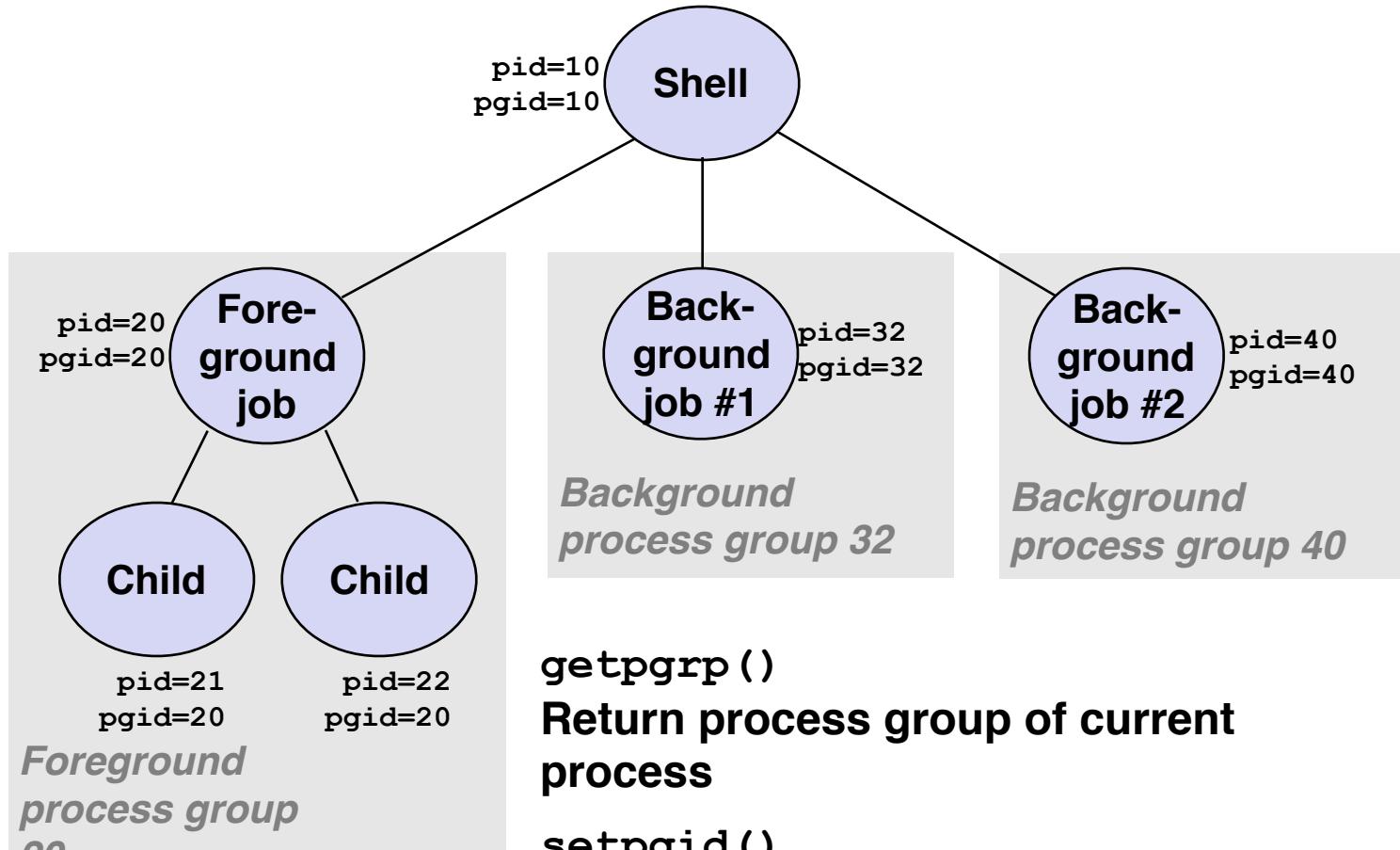
Process Groups

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`getpgrp()`

Return process group of current process

`setpgid()`

Change process group of a process

Sending Signals with /bin/kill Program

- /bin/kill program sends arbitrary signal to a process or process group

- Examples

- **/bin/kill -9 -24817**
Send SIGKILL to every process in process group 24817

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Child1: pid=24818 pgrp=24817
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```

```
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Sending Signals from the Keyboard

- Typing `ctrl-c` causes the kernel to send a `SIGINT` to every process in the foreground process group.
 - `SIGINT` – default action is to terminate each process
- Typing `ctrl-z` causes the kernel to send a `SIGTSTP` to every job in the foreground process group.
 - `SIGTSTP` – default action is to stop (suspend) each process

Example of ctrl-c and ctrl-z

```
bluefish> ./forks 17
Child: pid=28108 pgrp=28107
Parent: pid=28107 pgrp=28107
```

```
<types ctrl-z>
Suspended
bluefish> ps w
  PID TTY      STAT      TIME  COMMAND
27699 pts/8    Ss        0:00  -tcsh
28107 pts/8    T         0:01  ./forks 17
28108 pts/8    T         0:01  ./forks 17
28109 pts/8    R+        0:00  ps w
```

```
bluefish> fg
./forks 17
<types ctrl-c>
bluefish> ps w
  PID TTY      STAT      TIME  COMMAND
27699 pts/8    Ss        0:00  -tcsh
28110 pts/8    R+        0:00  ps w
```

STAT (process state)

Legend:

First letter:

S: sleeping

T: stopped

R: running

Second letter:

s: session leader

+: foreground proc group

See “man ps” for more details

Sending Signals with kill Function

```
void fork12()
{
    pid_t pid[N];
    int i;
    int child_status;

    for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0) {
            /* Child: Infinite Loop */
            while(1)
                ;
        }

    for (i = 0; i < N; i++) {
        printf("Killing process %d\n", pid[i]);
        kill(pid[i], SIGINT);
    }

    for (i = 0; i < N; i++) {
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child_status))
            printf("Child %d terminated with exit status %d\n",
                   wpid, WEXITSTATUS(child_status));
        else
            printf("Child %d terminated abnormally\n", wpid);
    }
}
```

forks.c