

CSC 252: Computer Organization

Spring 2021: Lecture 7

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Announcement

- Programming Assignment 2 is out
 - Details: <https://www.cs.rochester.edu/courses/252/spring2021/labs/assignment2.html>
 - Due on **March 3**, 11:59 PM
 - You (may still) have 3 slip days

14	15	16	17	18	19	20
21	22	23	24	25	26	27
		Today				
28	Mar 1	2	3	4	5	6

Due

Announcement

- Programming assignment 2 is in x86 assembly language.
- Read the instructions before getting started!!!
 - You get 1/4 point off for every wrong answer
 - Maxed out at 10
- TAs are best positioned to answer your questions about programming assignments!!!
- Programming assignments do NOT repeat the lecture materials. They ask you to synthesize what you have learned from the lectures and work out something new.
- Problem set on arithmetics: <https://www.cs.rochester.edu/courses/252/spring2021/handouts.html>.
 - Not to be turned in.

Data Movement

`movq (%rdi), %rdx`

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`addq 8(%rdi), %rdx`

Accessing memory and doing computation in one instruction. Allowed in x86, but not all ISAs allow that (e.g., MIPS).

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Illegal in x86 (and almost all other ISAs). Could make microarchitecture implementation inefficient/inelegant.

Today: Control Instructions

- Control: Conditional branches (`if... else...`)
- Control: Loops (`for, while`)
- Control: Switch Statements (`case... switch...`)

Conditional Branch Example

```
long absdiff
  (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

absdiff:

cmpq	%rsi,%rdi # x:y
jle	.L4
movq	%rdi,%rax
subq	%rsi,%rax
ret	
.L4: # x <= y	
movq	%rsi,%rax
subq	%rdi,%rax
ret	

Register	Use(s)
%rdi	x
%rsi	y
%rax	Return value

0 0 0
ZF SF OF

Conditional Branch Example

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ret	

cmpq sets ZF, SF, OF
jle checks ZF | (SF ^ OF)

0	0	0
ZF	SF	OF

Conditional Branch Example

```
unsigned long absdiff
(unsigned long x, unsigned
long y)
{
    unsigned long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	x
%rsi	y
%rax	Return value

absdiff:

```
    cmpq    %rsi,%rdi # x:y
    jle     .L4
    movq    %rdi,%rax
    subq    %rsi,%rax
    ret

.L4:      # x <= y
    movq    %rsi,%rax
    subq    %rdi,%rax
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```

ZF SF OF

Conditional Branch Example

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Register	Use(s)
%rdi	x
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absdiff:

```
    cmpq    %rsi,%rdi # x:y
    jbe     .L4
    movq    %rdi,%rax
    subq    %rsi,%rax
    ret
.L4:      # x <= y
    movq    %rsi,%rax
    subq    %rdi,%rax
    ret
```

0 0 0
ZF SF OF

Conditional Jump Instruction

```
cmpq    %rsi, %rdi  
jbe    .L4
```

Conditional Jump Instruction

cmpq

%rsi , %rdi

jbe

.L4



Jump to label if
below or equal to

Conditional Jump Instruction

cmpq **%rsi**, **%rdi**
jbe .L4 ← Jump to label if
 below or equal to

- Semantics of **jbe**:

- Treat the data in **%rdi** and **%rsi** as **unsigned values**.
- If **%rdi** is less than or equal to **%rsi**, jump to the part of the code with a label **.L4**

Conditional Jump Instruction

`cmpq`

`%rsi, %rdi`

`jbe`

`.L4`



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- Under the hood:

Conditional Jump Instruction

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- **cmpq** instruction sets the condition codes

Conditional Jump Instruction

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- Under the hood:
 - **cmpq** instruction sets the condition codes
 - **jbe** reads and checks the condition codes
 - If condition met, modify the Program Counter to point to the address of the instruction with a label **.L4**

How Should cmpq Set Condition Codes?

cmpq **%rsi, %rdi**

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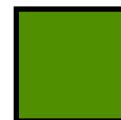
- How do we know $\%rdi \leq \%rsi$? This time for unsigned values
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ZF Zero Flag (result is zero)



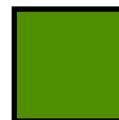
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$$\begin{array}{r} 001 \\ -) 111 \\ \hline \textcolor{red}{C}010 \end{array}$$

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ZF

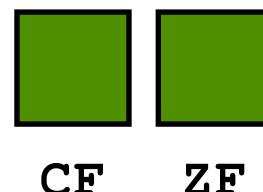
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How Should `cmpq` Set Condition Codes?

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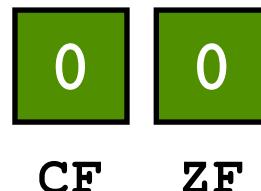
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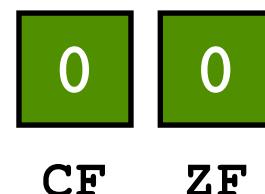
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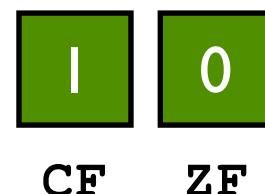
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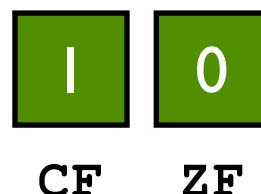
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- $\%rdi \leq \%rsi$ (as unsigned) if and only if:
 - ZF is set, or
 - CF is set
- or simply: **ZF | CF**
- This is what `jbe` checks

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Putting It All Together

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- `cmpq` sets all 4 condition codes simultaneously

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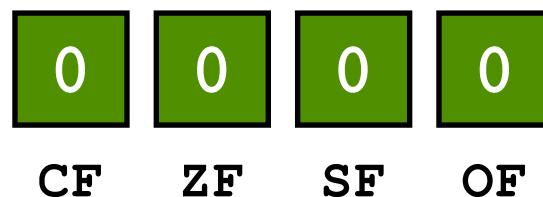
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SF Sign Flag

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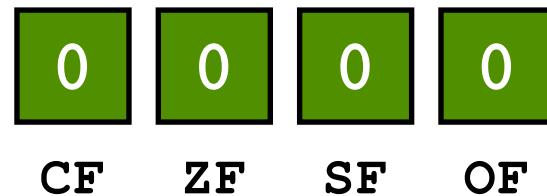
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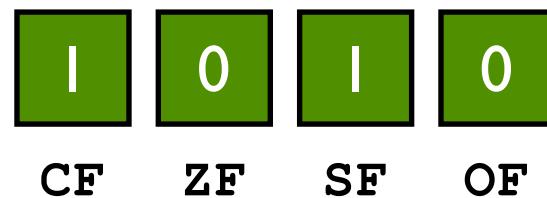
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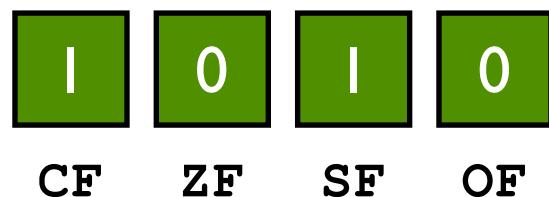
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cmpq    %rsi,%rdi  
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- cmpq sets all 4 condition codes simultaneously
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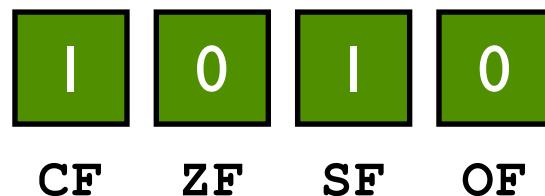
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ZF Zero Flag

CF Carry Flag

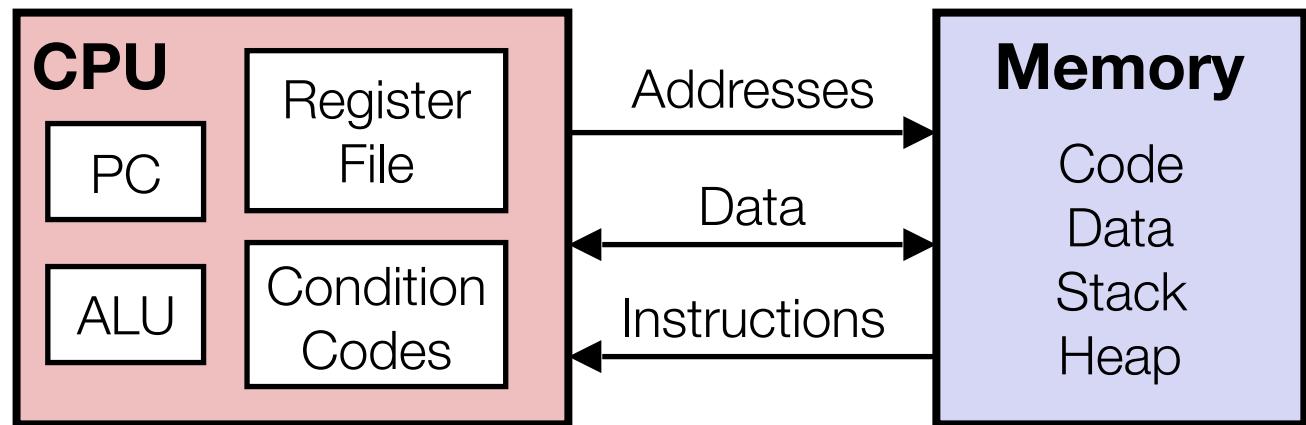
SF Sign Flag

OF Overflow Flag (for signed)



Condition Codes Hold Test Results

Assembly
Programmer's
Perspective
of a Computer



- Condition Codes
 - Hold the status of most recent test
 - 4 common condition codes in x86-64
 - A set of special registers (more often: bits in one single register)
 - Sometimes also called: Status Register, Flag Register



CF Carry Flag

SF Sign Flag

ZF Zero Flag

OF Overflow Flag (for signed)

Jump Instructions

- Jump to different part of code (designated by a label) depending on condition codes

jle	$(SF \wedge OF) \mid ZF$	Less or Equal (Signed)
------------	--------------------------	------------------------

jbe	$CF \mid ZF$	Below or Equal (unsigned)
------------	--------------	---------------------------

Jump Instructions

Instruction	Jump Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	$\sim ZF$	Not Equal / Not Zero
js	SF	Negative
jns	$\sim SF$	Nonnegative
jg	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
jge	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
jl	$(SF \wedge OF)$	Less (Signed)
jle	$(SF \wedge OF) \mid ZF$	Less or Equal (Signed)
ja	$\sim CF \ \& \ \sim ZF$	Above (unsigned)
jae	$\sim CF$	Above or Equal (unsigned)
jb	CF	Below (unsigned)
jbe	$CF \mid ZF$	Below or Equal (unsigned)

Implicit Set Condition Codes

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$$\begin{array}{r} 10000000 \\ +) 11111111 \\ \hline \text{c}01111111 \end{array}$$

0	0	0	0
CF	ZF	SF	OF

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jle .L4

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Implicit Set Condition Codes

addq %rax, %rbx

- Arithmetic instructions implicitly set condition codes (think of it as side effect)

- **CF** set if $\%rax + \%rbx$ generates a carry (i.e., unsigned overflow)
- **ZF** set if $\%rax + \%rbx == 0$
- **SF** set if $\%rax + \%rbx < 0$
- **OF** set if $\%rax + \%rbx$ overflows when $\%rax$ and $\%rbx$ are treated as signed numbers
 - $\%rax > 0$, $\%rbx > 0$, and $(\%rax + \%rbx) < 0$, or
 - $\%rax < 0$, $\%rbx < 0$, and $(\%rax + \%rbx) \geq 0$

```
if((x+y)<0) { addq 0xFF, 0x80  
...  
}
```

jle .L4

I	0	0	I
CF	ZF	SF	OF

Today: Control Instructions

- Control: Conditional branches (`if... else...`)
- Control: Loops (`for, while`)
- Control: Switch Statements (`case... switch...`)

“Do-While” Loop Example

- Popcount: Count number of 1's in argument x

do-while version

```
long pcount_do
(unsigned long x) {
    long result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

“Do-While” Loop Example

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(unsigned long x) {
    long result = 0;
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        x >>= 1;
    } while (x);
    return result;
}
```

goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

“Do-While” Loop Assembly

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

“Do-While” Loop Assembly

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long pcount_goto  
    (unsigned long x) {  
    long result = 0;  
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        x >>= 1;  
        if(x) goto loop;  
        return result;  
}
```

Register	Use(s)
%rdi	Argument x
%rax	result

```
    movl    $0, %rax      # result = 0  
.L2:                      # loop:  
    movq    %rdi, %rdx  
    andl    $1, %rdx      # t = x & 0x1  
    addq    %rdx, %rax    # result += t  
    shrq    $1, %rdi      # x >>= 1  
    jne     .L2           # if (x) goto loop  
    ret
```

“Do-While” Loop Assembly

```
long pcount_goto  
    (unsigned long x) {  
        long result = 0;  
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“Do-While” Loop Assembly

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“Do-While” Loop Assembly

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“Do-While” Loop Assembly

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        result += x & 0x1;  
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“Do-While” Loop Assembly

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    (unsigned long x) {  
    long result = 0;  
    loop:  
        result += x & 0x1;  
        x >>= 1;  
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    addq    %rdx, %rax    # result += t  
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    jne     .L2           # if (x) goto loop  
    ret
```

“Do-While” Loop Assembly

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long pcount_goto  
    (unsigned long x) {  
    long result = 0;  
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        result += x & 0x1;  
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“Do-While” Loop Assembly

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long pcount_goto  
    (unsigned long x) {  
    long result = 0;  
    loop:  
        result += x & 0x1;  
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        return result;  
}
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    addq    %rdx, %rax    # result += t  
    shrq    $1, %rdi      # x >>= 1  
    jne     .L2           # if (x) goto loop  
    ret
```

General “Do-While” Translation

do-while version

```
<before>;  
do {  
    body;  
} while (A < B);  
<after>;
```

goto Version

```
<before>  
.L1: <body>  
      if (A < B)  
          goto .L1  
<after>
```



General “Do-While” Translation

do-while version

```
<before>;  
do {  
    body;  
} while (A < B);  
<after>;
```

goto Version

```
.L1: <body>  
if (A < B)  
    goto .L1  
<after>
```



Replace with a
conditional jump
instruction



General “Do-While” Translation

do-while version

```
<before>;  
do {  
    body;  
} while (A < B);  
<after>;
```

goto Version

```
.L1: <body>  
    if (A < B)  
        goto .L1  
<after>
```

Assembly
Version



```
<before>  
.L1: <body>  
    cmpq B, A  
    jl .L1  
<after>
```

General “While” Translation

while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```

General “While” Translation

while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```



goto Version

```
<before>  
    goto .L2  
.L1: <body>  
.L2: if (A < B)  
        goto .L1  
<after>
```

General “While” Translation

while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```



goto Version

```
<before>  
goto .L2  
.L1: <body>  
.L2: if (A < B)  
      goto .L1  
<after>
```



Assembly
Version

```
<before>  
jmp .L2  
.L1: <body>  
.L2: cmpq A, B  
    jg .L1  
<after>
```

General “While” Translation

while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```



goto Version

```
<before>  
goto .L2  
.L1: <body>  
.L2: if (A < B)  
      goto .L1  
<after>
```



Assembly
Version

```
<before>  
jmp .L2  
.L1: <body>  
.L2: cmpq A, B  
    jg .L1  
<after>
```

General “While” Translation

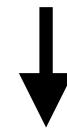
while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```



goto Version

```
<before>  
goto .L2  
.L1: <body>  
.L2: if (A < B)  
      goto .L1  
<after>
```



Assembly
Version

```
<before>  
jmp .L2  
.L1: <body>  
.L2: cmpq A, B  
     jg .L1  
<after>
```

“While” Loop Example

while version

```
long pcount_while
(unsigned long x) {

    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;

}
```

“While” Loop Example

while version

```
long pcount_while
(unsigned long x) {

    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

goto Version

```
long pcount_goto_jtm
(unsigned long x) {
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x) goto loop;
    return result;
}
```

“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

```
//assume unsigned int is 4 bytes  
long pcount_for (unsigned int x)  
{  
  
    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

init
i = 0

```
//assume unsigned int is 4 bytes  
long pcount_for (unsigned int x)  
{  
  
    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

“For” Loop Example

```
for (init; test; update) {  
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}
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//assume unsigned int is 4 bytes  
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    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

init

i = 0

test

i < 32

“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

```
//assume unsigned int is 4 bytes  
long pcount_for (unsigned int x)  
{  
  
    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

init

i = 0

test

i < 32

update

i++

“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

```
//assume unsigned int is 4 bytes  
long pcount_for (unsigned int x)  
{  
  
    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

init

i = 0

test

i < 32

update

i++

body

{

result += (x >> i)
 & 0x1;
}

Convert “For” Loop to “While” Loop

For Version

```
before;  
for (init; test; update) {  
    body;  
}  
after
```

Convert “For” Loop to “While” Loop

For Version

```
before;  
for (init; test; update) {  
    body;  
}  
after
```



While Version

```
before;  
init;  
while (test) {  
    body;  
    update;  
}  
after;
```

Convert “For” Loop to “While” Loop

For Version

```
before;  
for (init; test; update) {  
    body;  
}  
after
```

While Version

```
before;  
init;  
while (test) {  
    body;  
    update;  
}  
after;
```

Assembly Version

```
before  
init  
jmp .L2  
.L1: body  
      update  
.L2: cmpq A, B  
      jg .L1  
after
```



Today: Control Instructions

- Control: Conditional branches (`if... else...`)
- Control: Loops (`for, while`)
- Control: Switch Statements (`case... switch...`)

Switch Statement Example

```
long switch_eg (long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

Switch Statement Example

```
long switch_eg (long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:      Fall-through case
            w = y/z;
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

Switch Statement Example

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long switch_eg (long x, long y, long z)
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    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;      Fall-through case
        case 3:
            w += z;
            break;
        case 5:
        case 6:      Multiple case
            w -= z;      labels
            break;
        default:
            w = 2;
    }
    return w;
}
```

Switch Statement Example

```
long switch_eg (long x, long y, long z)
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    long w = 1;
    switch(x) {
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        case 3:
            w += z;
            break;
        case 5:
        case 6:          Multiple case
            w -= z;      labels
            break;
        default:         For missing
            w = 2;      cases, fall back
    }                  to default
    return w;
}
```

Switch Statement Example

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            w += z;
            break;
        case 5:
        case 6:      Multiple case
            w -= z;      labels
            break;
        default:
            w = 2;      ← For missing
    }                  cases, fall back
    return w;          to default
}
```

Converting to a cascade of if-else statements is simple, but cumbersome with too many cases.

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ...  
    case val_{n-1}:  
        Block n-1  
}
```

Jump Targets

Targ₀: Code Block 0

Targ₁: Code Block 1

Targ₂: Code Block 2

•
•
•

Targ_{n-1}: Code Block n-1

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Jump Table

JTab:

Targ0
Targ1
Targ2
•
•
•
Targn-1

Jump Targets

Targ0:

Code Block 0

Targ1:

Code Block 1

Targ2:

Code Block 2

•
•
•

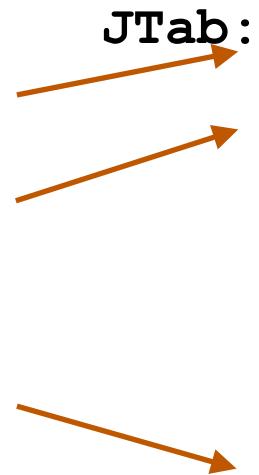
Targn-1:

Code Block n-1

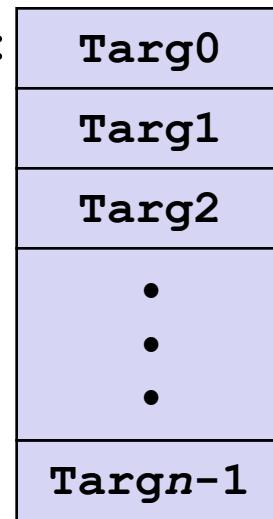
Implementing Switch Using Jump Table

Switch Form

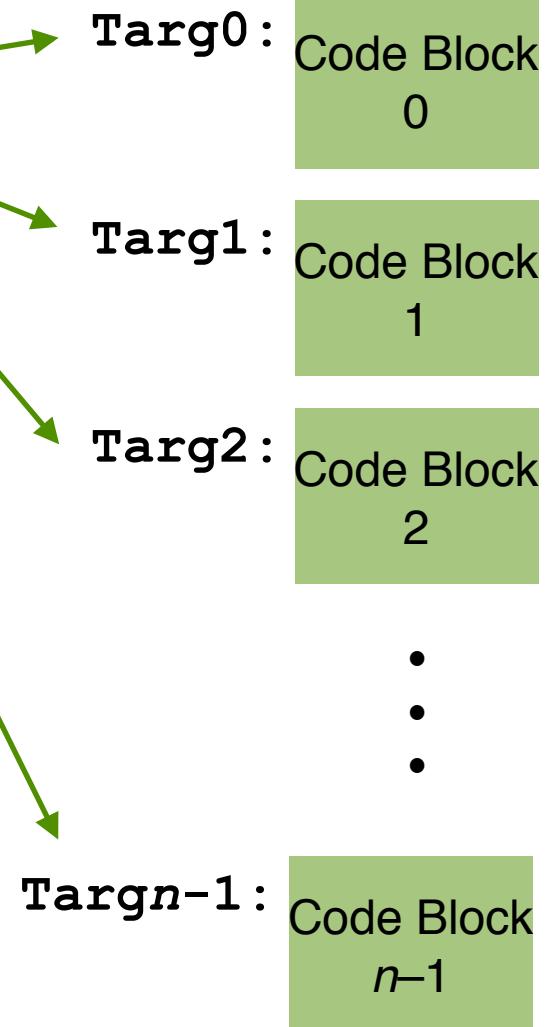
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switch(x) {  
    case val_0:  
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    case val_1:  
        Block 1  
    ....  
    case val_{n-1}:  
        Block n-1  
}
```



Jump Table



Jump Targets

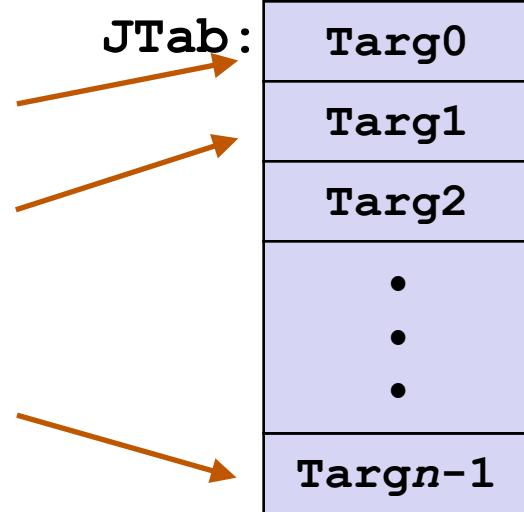


Implementing Switch Using Jump Table

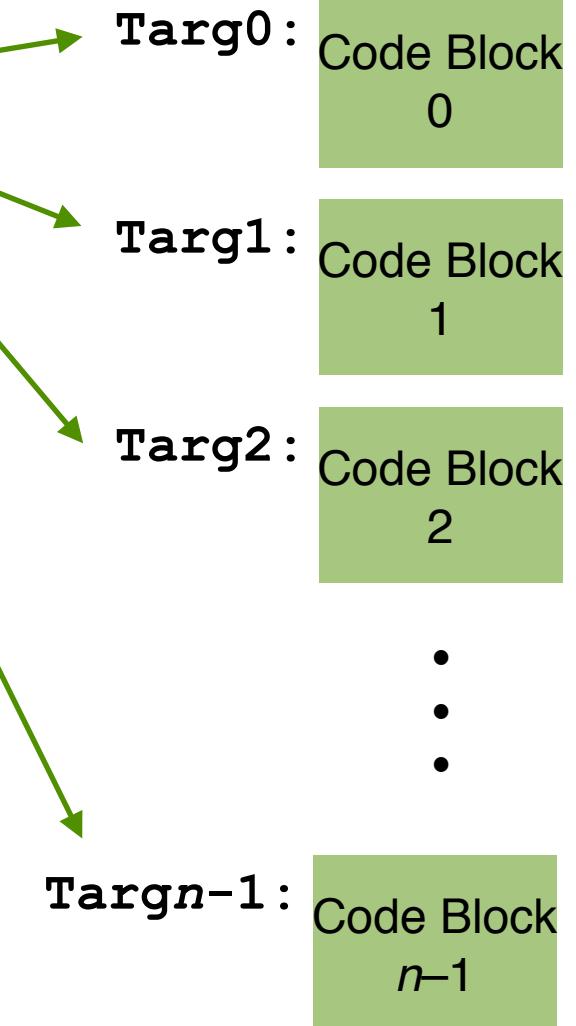
Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Jump Table



Jump Targets



- Each code block starts from a unique address (Targ0, Targ1, ...)
- Jump table stores all the target address
- Use the case value to index into the jump table to find where to jump to

Assembly Directives (Pseudo-Ops)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

- Directives:
 - Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

Assembly Directives (Pseudo-Ops)

```
.section .rodata
.align 8
.L4:
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.quad .L5 # x = 5
.quad .L5 # x = 6
```

- **.quad**: tells the assembler to set aside the next 8 bytes in memory and initialize with the value of the operand (a label here, which itself is an address)

- Directives:
 - Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

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.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

- **.quad**: tells the assembler to set aside the next 8 bytes in memory and initialize with the value of the operand (a label here, which itself is an address)
- **.align**: tells the assembler that addresses of the the following data will be aligned to 8 bytes

- Directives:
 - Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

Assembly Directives (Pseudo-Ops)

```
.section .rodata
.align 8
.L4:
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.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

- Directives:

- Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

- **.quad**: tells the assembler to set aside the next 8 bytes in memory and initialize with the value of the operand (a label here, which itself is an address)
- **.align**: tells the assembler that addresses of the the following data will be aligned to 8 bytes
- **.section**: denotes different parts of the object file

Assembly Directives (Pseudo-Ops)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

- Directives:
 - Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

- **.quad**: tells the assembler to set aside the next 8 bytes in memory and initialize with the value of the operand (a label here, which itself is an address)
- **.align**: tells the assembler that addresses of the the following data will be aligned to 8 bytes
- **.section**: denotes different parts of the object file
- **.rodata**: read-only data section

Jump Table and Jump Targets

Jump Table

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

jmp .L3 will go
to .L3 and start
executing from there

Jump Targets

```
.L1:          # Case 1
    movq    %rsi, %rax
    imulq   %rdx, %rax
    jmp     .done
.L2:          # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx
.L3:          # Case 3
    addq    %rcx, %rax
    jmp     .done
.L5:          # Case 5, 6
    subq    %rdx, %rax
    jmp     .done
.LD:          # Default
    movl    $2, %eax
    jmp     .done
```

Implementing Switch Using Jump Table

Switch Form

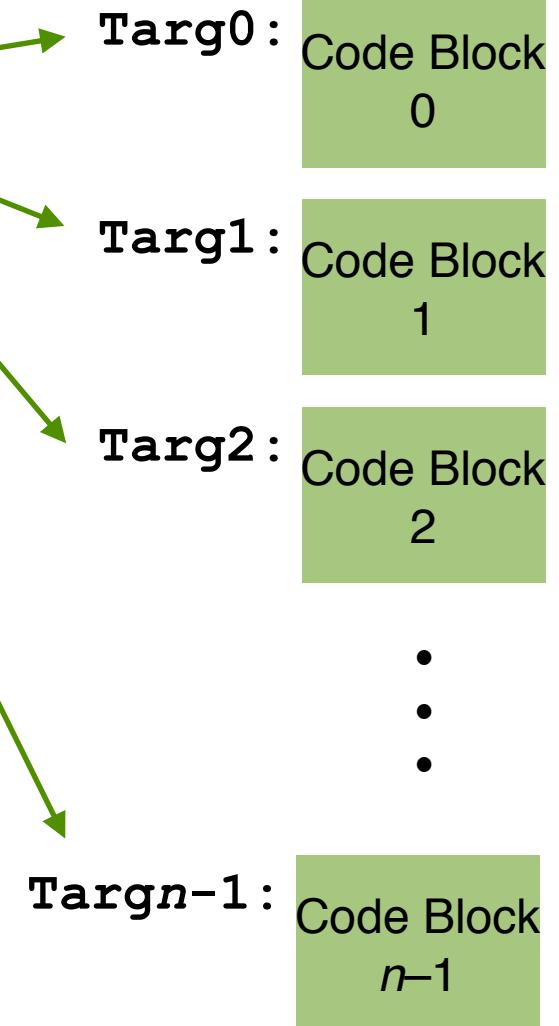
```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_{n-1}:  
        Block n-1  
}
```

JTab:

Jump Table



Jump Targets

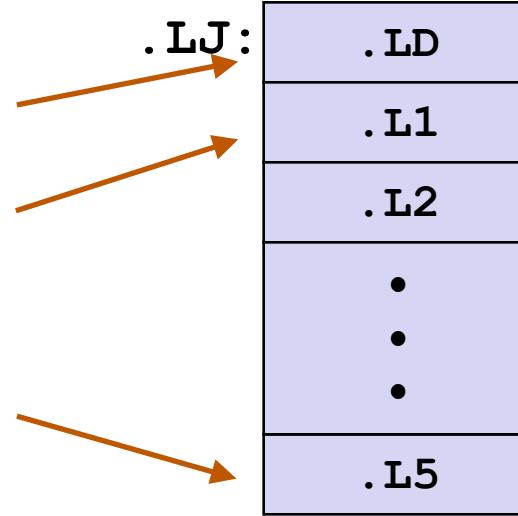


Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Jump Table



Jump Targets

- .LD : Code Block 0
- .L1 : Code Block 1
- .L2 : Code Block 2
-
-
-
- .L5 : Code Block n-1
- .Done :

Code Blocks ($x == 1$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

Code Blocks ($x == 1$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
case 1:           // .L1
    w = y*z;
    break;
...
}
```

Code Blocks ($x == 1$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
case 1:          // .L1
    w = y*z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks ($x == 1$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
case 1:           // .L1
    w = y*z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L1:
    movq    %rsi, %rax    # y
    imulq   %rdx, %rax    # y*z
    jmp     .done
```

Code Blocks ($x == 1$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
case 1:           // .L1
    w = y*z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L1:
    movq    %rsi, %rax    # y
    imulq   %rdx, %rax    # y*z
    jmp     .done
```

Code Blocks ($x == 2$, $x == 3$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks ($x == 2$, $x == 3$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 2:           // .L2
    w = y/z;
    /* Fall Through */
case 3:           // .L3
    w += z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks ($x == 2$, $x == 3$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 2:           // .L2
    w = y/z;
    /* Fall Through */
case 3:           // .L3
    w += z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L2:          # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx          # y/z

.L3:          # Case 3
    addq    %rcx, %rax # w += z
    jmp     .done
```

Code Blocks ($x == 2$, $x == 3$)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 2:           // .L2
    w = y/z;
    /* Fall Through */
case 3:           // .L3
    w += z;
    break;
...
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L2:          # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx          # y/z
.L3:          # Case 3
    addq    %rcx, %rax # w += z
    jmp     .done
```

Code Blocks ($x == 5$, $x == 6$, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks ($x == 5$, $x == 6$, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 5: // .L5
case 6: // .L5
    w -= z;
    break;
default: // .LD
    w = 2;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Code Blocks ($x == 5$, $x == 6$, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 5: // .L5
case 6: // .L5
    w -= z;
    break;
default: // .LD
    w = 2;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L5:                      # Case 5,6
    subq %rdx, %rax # w -= z
    jmp .done
.LD:                      # Default:
    movl $2, %eax   # 2
    jmp .done
```

Code Blocks ($x == 5$, $x == 6$, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 5: // .L5
case 6: // .L5
    w -= z;
    break;
default: // .LD
    w = 2;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

```
.L5:                      # Case 5,6
    subq %rdx, %rax # w -= z
    jmp .done
.LD:                      # Default:
    movl $2, %eax   # 2
    jmp .done
```

Code Blocks ($x == 5$, $x == 6$, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
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.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 5: // .L5
case 6: // .L5
    w -= z;
    break;
default: // .LD
    w = 2;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

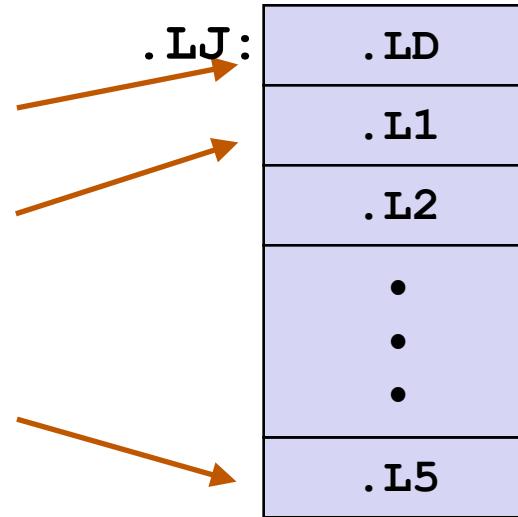
```
.L5:                                # Case 5,6
    subq %rdx, %rax # w -= z
    jmp .done
.LD:                                # Default:
    movl $2, %eax   # 2
    jmp .done
```

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Jump Table



Jump Targets

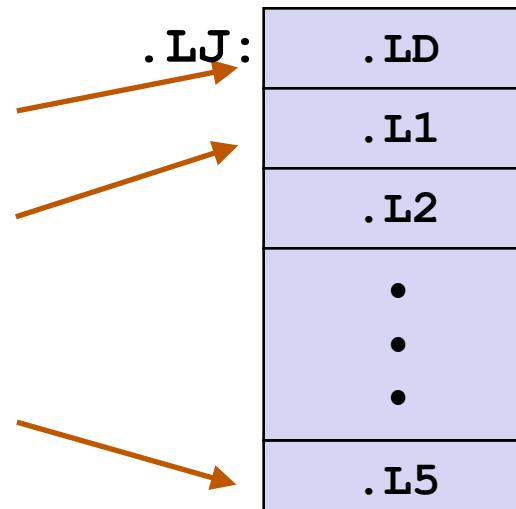
- .LD : Code Block 0
- .L1 : Code Block 1
- .L2 : Code Block 2
-
-
-
- .L5 : Code Block n-1
- .Done :

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
    case val_0:  
        Block 0  
    case val_1:  
        Block 1  
    ....  
    case val_n-1:  
        Block n-1  
}
```

Jump Table



Jump Targets

.LD:	Code Block 0
.L1:	Code Block 1
.L2:	Code Block 2
.L3:	Code Block 3
.L5:	Code Block n-1
.Done:	

- The only thing left...
 - How do we jump to different locations in the jump table depending on the case value?

Indirect Jump Instruction

```
.section  .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

Indirect Jump Instruction

The address we want to jump to is stored at .LJ + 8 * x

```
.section  .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

Indirect Jump Instruction

The address we want to jump to is stored at `.LJ + 8 * x`

```
.section .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
# assume x in %rdi
movq .LJ(,%rdi,8), %rax
jmp *%rax
```

Indirect Jump Instruction

The address we want to jump to is stored at `.LJ + 8 * x`

```
.section .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
# assume x in %rdi
movq .LJ(,%rdi,8), %rax
jmp *%rax
```

- Indirect Jump: `jmp *%rax`
 - `%rax` specifies the address to jump to (`PC = %rax`)

Indirect Jump Instruction

The address we want to jump to is stored at `.LJ + 8 * x`

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.section .rodata
.align 8
.LJ:
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.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
# assume x in %rdi
movq .LJ(,%rdi,8), %rax
jmp *%rax
```

- Indirect Jump: `jmp *%rax`
 - `%rax` specifies the address to jump to (`PC = %rax`)
- Direct Jump (`jmp .LJ`), directly specifies the jump address

Indirect Jump Instruction

The address we want to jump to is stored at $.LJ + 8 * x$

```
.section .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
# assume x in %rdi
movq .LJ(,%rdi,8), %rax
jmp *%rax
```

- Indirect Jump: **jmp *%rax**
 - %rax specifies the address to jump to ($PC = \%rax$)
- Direct Jump (**jmp .LJ**), directly specifies the jump address
- Indirect Jump specifies where the jump address is located

Indirect Jump Instruction

The address we want to jump to is stored at `.LJ + 8 * x`

```
.section .rodata
.align 8
.LJ:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
# assume x in %rdi
movq .LJ(,%rdi,8), %rax
jmp *%rax
```

- Indirect Jump: `jmp *%rax`
 - `%rax` specifies the address to jump to (`PC = %rax`)
- Direct Jump (`jmp .LJ`), directly specifies the jump address
- Indirect Jump specifies where the jump address is located

An equivalent syntax in x86: `jmp * .LJ(,%rdi,8)`