Final Exam CSC 252 4 May 2022 Computer Science Department University of Rochester

Instructor: Yuhao Zhu

TAs: Nisarg Ujjainkar, Abhishek Tyag, Kalen Frieberg, Gunnar Hammonds, Mandar Juvekar, Zihao Lin, Vladimir Maksimovski, Yiyao (Jack) Yu

Name:	
Problem 0 (3 points):	
Problem 1 (12 points):	
Problem 2 (18 points):	
Problem 3 (25 points):	
Problem 4 (20 points):	
Problem 5 (26 points)	
Problem 6 (26 points):	
Total (130 points):	

Remember "**I don't know**" is given 15% partial credit, but you must erase everything else. This does not apply to extra credit questions.

Your answers to all questions must be contained in the given boxes. Use spare space to show all supporting work to earn partial credit.

You have 180 minutes to work.

Please sign the following. I have not given nor received any unauthorized help on this exam.

Signature:_____

GOOD LUCK!!!

Problem o: Warm-up (3 Points)

Do you think 252 could be taught in high school?

Problem 1: Miscellaneous (12 points)

Part a) (3 points) Write down the sum of 32_8 and 0x32 in base 2.

Part b) (3 points) Generally, the access time of a direct-mapped cache is _____ than that of a fully associative cache that is the same total size. Answer with <, >, =, <=, or >=.

Part c) (3 points) (True or False) Virtual memory has no size limitation.

Part d) (3 points) Consider the following C struct Person:

```
struct Person{
    long id;
    float age;
    float weight;
    float height;
    char name[20];
    char sex[7];
    struct Person * nextPerson;
};
```

What's the size of the struct Person?

Problem 2: ISA (18 points)

A stack-based ISA is designed. The ISA uses a hardware stack, and instructions in this ISA manipulate this stack. Each entry on the stack is one byte long, and the memory is byte-addressable. All instructions are 8-bit long, and are classified into two categories.

R Category

Binary encoding:

OpCode<7-5 bits>	00000<4-0 bits>
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Instruction list:

Instruction	OpCode	Role	
рор	001	Remove the item at the top of the stack.	
halt	001	Halt the processor.	

I Category

Binary encoding:

OpCode<7-5 bits>	Immediate value in 2's component<4-0 bits>
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Instruction list:

Instruction	OpCode	Role
pushi	000	Push sign-extended immediate value on the stack.
load	101	Let the top entry of the stack be A. Compute Address = A + sign-extended immediate value; Pop A from the stack; Use Address to load a byte from the main memory and push the byte on the stack.
store	100	Let the top entry of the stack be A, and the second entry of the stack be B. Compute Address = A + sign-extended immediate value; Pop A from the stack; Store B at the memory location Address; Pop B from the stack.

loadd		Address = sign-extended immediate value. Load a byte from the memory using Address and push the byte on the stack.
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Part a) (12 points) Encode the following instructions in binary.

(4 points) store -8

(4 points) pushi 7

(4 points) halt

Part b) (6 points) You want to implement a function that pops the top value from the stack and then pushes it back on the stack twice. Memory location at address 0 is reserved for this instruction to use temporarily. Write an assembly program using the existing instructions to implement this function.

Problem 3: Floating-Point Arithmetic (25 points)

Suppose that the IEEE decided to add a new *n*-bit floating-point standard, with its main characteristics consistent with the other IEEE standards. This *n*-bit standard can precisely represent the value $6\frac{3}{16}$, but cannot precisely represent $11\frac{7}{16}$. The smallest **positive normalized value** that can be represented in this standard is 2^{-30} .

Part a) (3 points) Convert $6\frac{3}{16}$ to Binary Normalized Form

Part b) (3 points) How many of the *n* bits are fraction bits?

Part c) (3 points) What's the bias of this standard?

Part d) (3 points) How many of the *n* bits are exponent bits?

Part e) (3 points) What is n?

Part f) (10 points) Suppose that using this new IEEE standard you perform two separate calculations, assume nearest-even rounding is used:

- 1. $(256 + 2\frac{1}{4}) 256$
- 2. $(256 256) + 2\frac{1}{4}$

What would be the result of these calculations? Do they both result in equivalent mathematically precise answers? Show your math to earn partial credit.

Problem 4: Cache (20 points)

For all the questions in this problem, assume that we are using a 12-bit machine with a byte-addressable memory and a direct-mapped cache. The cache can hold up to 16 cache lines.

Part a) (3 points) How many bits do you need for the set index?



Part b) (17 points) The following sequence of 9 memory accesses generates the hits/misses shown. Some miss/hit entries are intentionally left blank. The cache is initially empty. Note that the addresses are written in binary with spaces added between each 4 bits for readability. These are not necessarily the tag/index/offset boundaries.

#	Address	Hit/Miss
1	1101 1111 0000	Miss
2	0000 1101 1111	Miss
3	1101 0111 0101	Miss
4	0000 1101 1100	Hit
5	1101 1111 0011	Miss
6	1111 0111 0010	
7	1101 1111 0000	
8	0000 1101 1101	
9	1111 0111 0100	

(3 points) What is the number of tag bits?

(3 points) What is the number of offset bits?

(3 points) What is the size of each cache line (ignore the valid bit, dirty bit, and tag bits etc.)? Show the formula you used to calculate this, and the value you get from it.

(8 points) Fill the miss/hit for each of the blank entries.

Problem 5: Assembly Programming (26 points)

Conventions:

- 1. For this section, the assembly shown uses the AT&T/GAS syntax opcode src, dst for instructions with two arguments where src is the source argument and dst is the destination argument. For example, this means that mov a, b moves the value a into b.
- 2. All C code is compiled on a **64-bit machine**, where arrays grow toward higher addresses.
- 3. For functions that take an argument, the argument is stored in **%rdi** at the time the function is called. The return value of this function is stored in **%rax** at the time the function returns.

Consider the assembly of a C function void foobar (int *x) which takes a single int pointer parameter x.

```
000000000001159 <foobar>:
```

1159:	mov	(%rdi),%eax
115b:	inc	%eax
115d:	mov	<pre>%eax,(%rdi)</pre>
115f:	lea	0x2ecf,%rdi
1166:	cmp	\$0x1,%eax
1169:	je	1177 <foobar+0x1e></foobar+0x1e>
116b:	cmp	\$0x2,%eax
116e:	jne	117c <foobar+0x23></foobar+0x23>
1170:	lea	0x2eb4,%rdi
1177:	call	1030 <puts@plt></puts@plt>
117c:	ret	

The addresses **0x2ecf** and **0x2eb4** contain the beginning of character strings "**foo**" and "**bar**" respectively. **puts** () is a standard C function that prints a character string given the string start address as a parameter.

Part a) (9 points)

(3 points) When the initial value of ***x** is 0, what is printed by the program?

(3 points) When the initial value of ***x** is 1, what is printed by the program?

(3 points) The call on line 1177 is replaced with a jmp. Would this program still run correctly? Explain your answer.

Part b) (17 points)

Assuming that there are two threads, each of which executes the same **foobar()** function with the same exact parameter **x**. ***x** is initially set to **0**. For the sake of this problem, assume that the instructions inside the **puts()** function are always executed as an atomic unit (they are either executed together without interruption or not executed at all).

(4 points) What are all the possible values of *x after both threads finish execution of foobar()?

(4 points) What are all the possible strings printed by this multi-threaded program?

(3 points) Suppose the first three instructions (1159 to 115d) could be replaced with a single atomic instruction called csc252. Would this guarantee that the program always ends with x containing the value 2? Show you work to earn partial credit.

(3 points) Would this guarantee that the program always prints the same string no matter how many times you run it? Show you work to earn partial credit.

(3 points) Assuming that the two threads execute on two different processors, each with a separate cache. What is something that the processor designers have to pay attention to in order to correctly implement the csc252 instruction?

Problem 6: Virtual Memory (26 points)

Assume a virtual memory system that has the following characteristics:

- 1. The virtual address space is 32 KB and is byte addressable
- 2. Physical memory size is 8 KB and is byte addressable
- 3. Page size is 128 Bytes
- 4. One level page table, where each page table entry contains a valid bit, a dirty bit, and the physical page number
- 5. PTBR is ox3ADC
- 6. There is a data TLB that stores only the last page table entry

The format of a PTE is as shown below. MSB is the valid bit followed by the dirty bit. Last few bits are the physical page number (PPN).

valid<1 bit>	dirty<1 bit>	PPN <n bits=""></n>
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Part a) (3 points) What are the number of physical and virtual pages?

Part b) (3 points) What is the total size of the page table?

Consider the following C program:

```
void fibbo(int a[64]) {
    a[0] = 0
    a[1] = 1
    for (int i=2; i < 64; i++) {
        a[i] = a[i-2] + a[i-1];
    }
}</pre>
```

Suppose that the virtual address of the array 'a' is 0x0400. Assume the data TLB is empty when the code starts execution. The table below shows a part of the main memory before the code executes.

Address	Data
3ADC	B8
3ADD	3A
3ADE	CD
3ADF	78
3AE4	F9
3CDC	В6
3DDC	4F
3EDC	Fo

Part c) (4 points) How many pages does array 'a' occupy?

Part d) (4 points) To read a [1], what virtual page number(s) is(are) accessed?

Part e) (8 points) What physical memory addresses are accessed when reading a [1]?

Part f) (4 points) How many data TLB misses will occur in the execution of the program?