

## Recursion

equally powerful to iteration, and as efficient in cases where you can use tail recursion.

mechanical transformations back and forth  
often more intuitive (sometimes less)

**naive** implementation less efficient

no special syntax required

fundamental to functional languages like Scheme

tail recursion

```
(* OCaml: *)
let rec gcd b c =
  if b = c then b
  else if b < c then gcd b (c - b)
  else gcd (b - c) c;;
```

implemented as

```
gcd (b c)
start:
  if b = c
    return b
  if b < c
    c := c - b
    goto start
  if b > c
    b := b - c
    goto start
```

changes to create tail recursion (e.g. pass along an accumulator)

```
(* OCaml: *)
let rec summation f low high =
  if low == high then f low
  else f low + summation f (low+1) high;;
```

becomes

```
let rec summation2 f low high st =
  if low == high then st + f low
  else summation2 f (low+1) high (st + f low);;
```

and then

```
let summation3 f low high =
  let rec helper low st =
    let new_st = st + f low in
    if low == high then new_st
    else helper (low+1) new_st in
  helper low 0;;
```

More generally (absent an associative operator), pass along a **continuation**.

This is perfectly natural to someone used to programming in a functional language. Note that the summation example depends for correctness on the associativity of addition. To sum the elements in the **same** order we could have counted down from high instead of up from low, but that makes a more drastic change to the structure of the recursive calls.

There is no perfectly general algorithm to discover tail-recursive versions of functions, but compilers for functional languages recognize all sorts of common cases.

Sisal and pH have "iterative" syntax for tail recursion:

```
function sum (f : function (n : integer returns integer),
              low : integer, high : integer returns integer)
for initial
  st := f (low);
while low <= high
  low := old low + 1
  st := old st + f (low)
returns value of st
end for
end function
```

The Sisal compiler was **really** good at finding tail recursive forms.

## Concurrency

specifies that statements are to occur (at least logically) concurrently  
concurrency is fundamental to probably half the research in computer science today  
subject of chapter 13

## Nondeterminacy

choice "doesn't matter"

periodically popular, promoted by Dijkstra for use with selection (**guarded command** syntax)

can apply to execution order as well

useful for certain kinds of concurrency

process server

```
do
  receive read request ->
  reply with data
[]
  receive write request ->
  update data and reply
od
```

also nice for certain axiomatic proof schemes

raises issues of "randomness", "fairness", "liveness", etc.