

# Semantics for Transactional Languages

**Michael L. Scott**



ASPLOS PC “mini symposium”

21 October 2011

thoughts inspired by a series of wonderful students

# A Simple Idea [Herlihy & Moss '93]

- User labels atomic sections

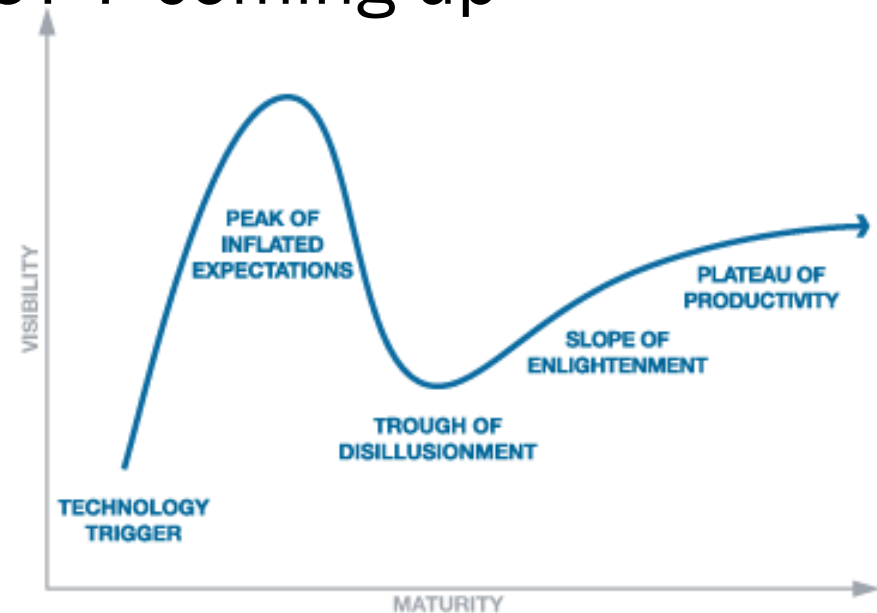
```
atomic {  
    ...  
}
```



- Underlying system ensures atomicity; executes in parallel when possible (speculation)
- Motivations
  - » performance (via lock elision)
  - » **simplicity: as easy as coarse-grain locks**

# Since Then

- Surge of interest ~2003, with the multicore revolution and with breakthroughs in HW (UWisc, UIUC) and SW (Sun, Cambridge)
- Scores of papers & systems; at least a dozen active groups; TRANSACT 7 coming up
- Through the trough of disillusionment?



# A Fleeting Opportunity

- HTM is coming
  - » Azul, Sun Rock, AMD ASF
  - » IBM has announced for Blue Gene/Q
  - » ... ?
- STM for backward compatibility, fallback on HW overflow
- Language support essential
- Narrow window in time to “get the semantics right”

# Outline

- Assertions
  - » **atomicity** is central
  - » speculation is an implementation issue (only)
  - » small transactions are what matter
  - » **privatization** is essential
    - necessary for correctness
    - solves the problem of legacy synchronization
- Open Questions
  - » **non-transactional reads and writes**
  - » big transactions, integration with system transactions
  - » relationship to “**deterministic parallel programming**”

# Memory Model

- Transactional sequential consistency (TSC)
  - » ideal but expensive: global total order on accesses
    - consistent w/ program order  $<_p$
    - w/ each transaction contiguous — atomic
- Strict serializability (SS)
  - » txns globally totally ordered wrt one another
  - » also ordered wrt preceding & following accesses of same thread (though those accesses aren't necessarily globally ordered wrt one another)
- Transactional Data-Race Freedom (TDRF)
  - » if all conflicting accesses are ordered by SS
  - » then the program appears to be TSC

# Strong Isolation Is a Non-Issue

- Hard to explain to the programmer
  - » what is a memory access?
- Heavy performance penalty in STM
- Only matters in racy programs
  - » constrains the behavior of buggy code
  - » less than you want (TSC); more than you need to build what you want (TSC given TDRF)
  - » may be useful for debugging, but a good race detector is better

# Opacity Is a Semantic Non-Issue

- Aborted transactions do not appear in (language-level) histories
- Opacity is simply one end of the implementation spectrum: validate at every read
- Sandboxing is the other end: validate before every “dangerous” operation (and periodically)
- Some very promising implementations in the middle: delayed/out-of-band validation
  - » ask me later!



# Privatization Is Essential

- Definition: transaction  $T$  with history prefix  $P$  privatizes datum  $D$  if
  - »  $\exists$  extensions of  $P$  in which a first access to  $D$  after  $P$  occurs in different threads
  - »  $\forall$  extensions of  $P+T$ , the first access to  $D$  after  $P+T$  occurs in the same one thread
- Crucial for performance with STM
- Solves the problem of legacy synchronization
  - » locking is privatization —  
`acquire` and `release` are small atomic blocks

# Transactions $\neq$ Critical Sections

~~L.acquire()  
...  
L.release()  $\equiv$  atomic {  
...  
}~~

L.acquire()  
...  
L.release()  $\equiv$  atomic { ... }  
...  
atomic { ... }

# Open Questions

# Non-transactional Accesses

- Want reads for, e.g., ordered speculation, high-performance hybrid TM
  - » clearly important at the HTM ISA level
  - » not clear whether needed/wanted at language API level
- Want writes out of aborted txns for debugging
  - » again, clearly important at the HTM ISA level
    - and probably more useful if immediate
  - » probably important at the language level too
    - not as clear that these need to be immediate
- Immediate writes, and writes in aborted txns, a challenge for the memory model

# Atomicity and Determinism

- See our paper at DISC'11
  - » languages/idioms that guarantee all abstract executions will be “equivalent” in some well-defined sense
- Independent split-merge an obvious foundation for language-level determinism
- Atomic commutative [+associative] ops an obvious extension
- Is there anything else?
  - » atomic event handlers, perhaps?

# The Bottom Line: Keep It Simple!

- Atomicity is central
- Speculation is an implementation issue (only)
- Small transactions are what matter
- Privatization is essential (and solves the problem of legacy synchronization)



UNIVERSITY *of*  
ROCHESTER

[www.cs.rochester.edu/research/synchronization/](http://www.cs.rochester.edu/research/synchronization/)

Thanks to Bill Scherer, Virendra Marathe,  
Mike Spear, Luke Dalessandro, Li Lu, ...