

CSC 162  
DATA STRUCTURES

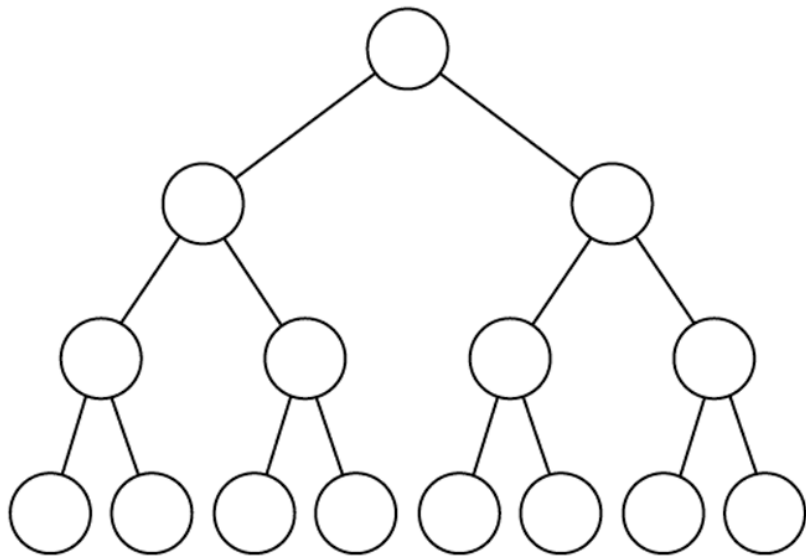
# A PROBLEM WITH BSTs

Common operations on *balanced* BST are  $O(\log(n))$

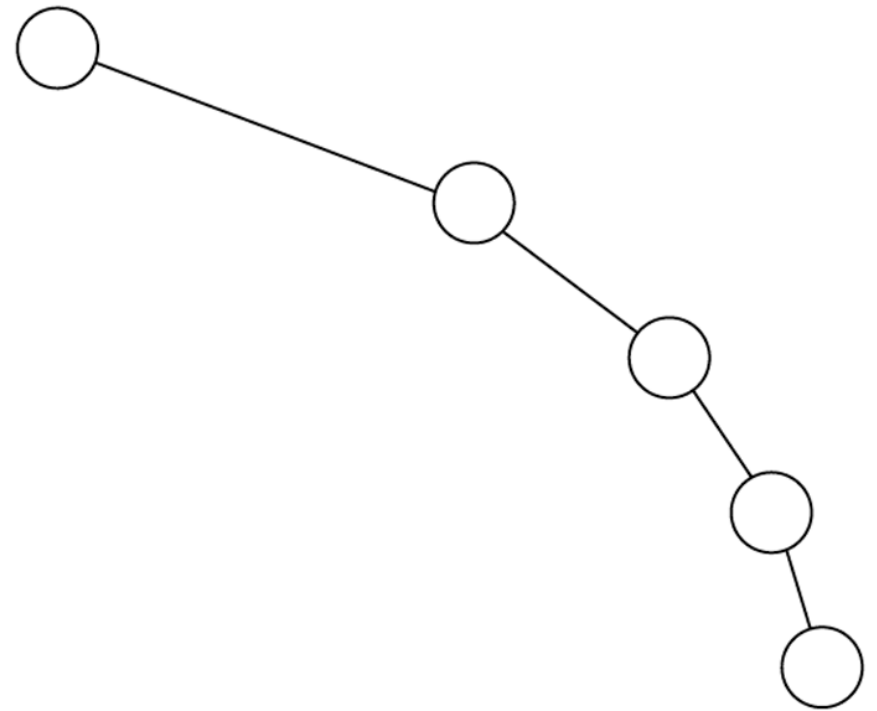
Alas, when the tree goes out of balance, performance degrades (worst case : chain  $O(n)$ )

There are several data structures that modify the BST to maintain balance.

# BST Structure : Log & Linear

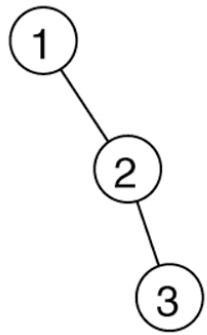


(a)

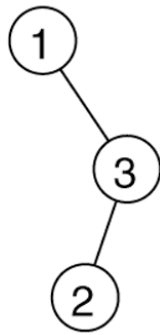


(b)

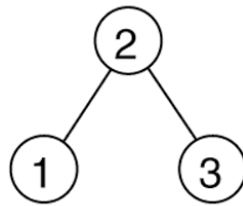
# Permutations of 1,2,3



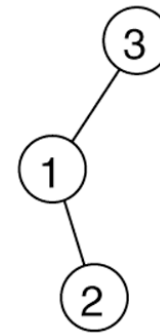
(a)



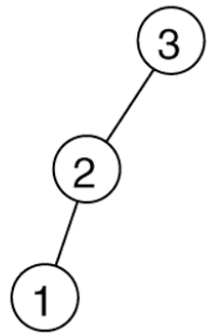
(b)



(c)



(d)



(e)

# AVL Trees

The first balanced binary search tree

Named after discoverers

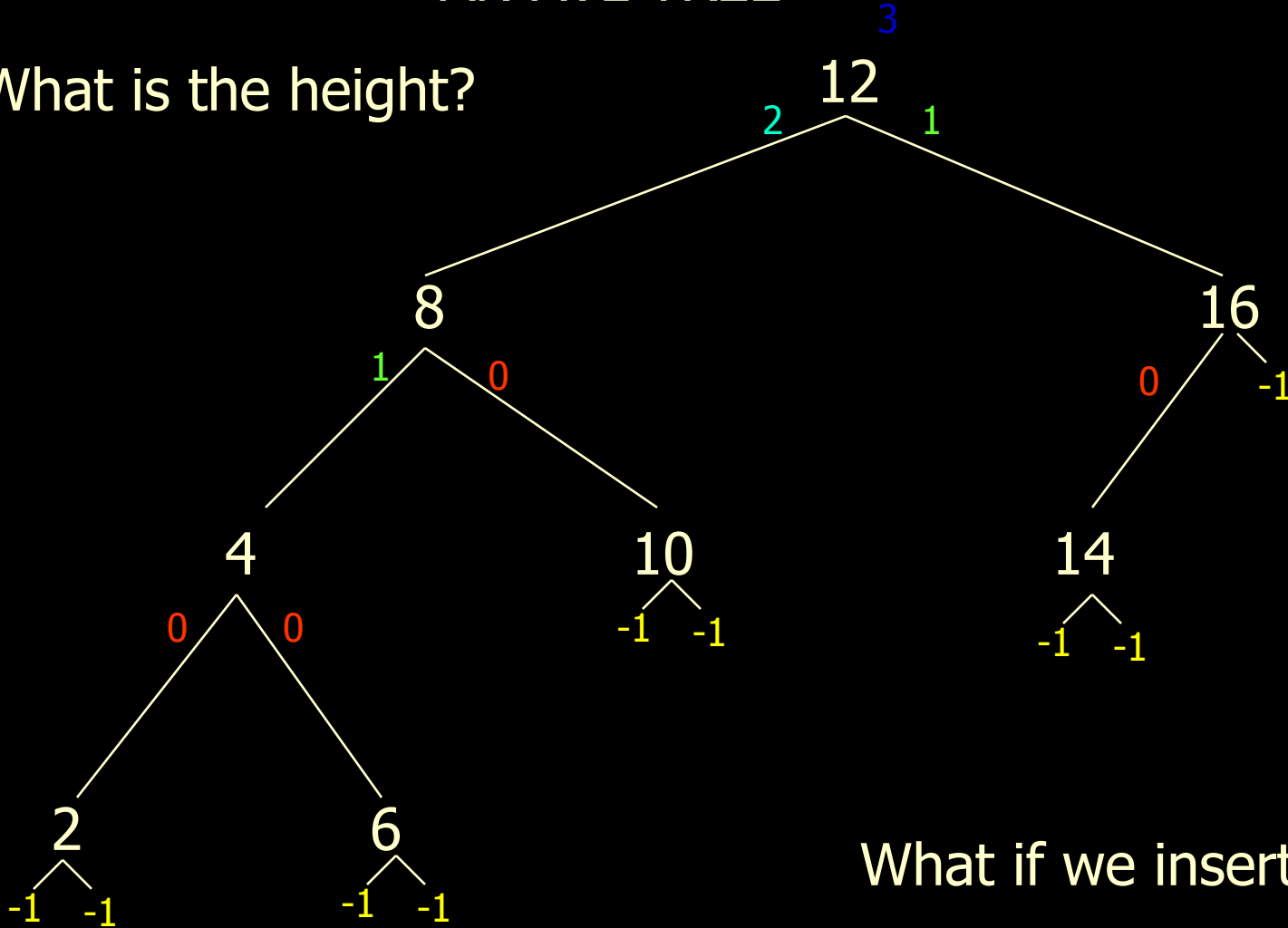
Adelson-Velskii and Landis.

## DEFINITION:

An AVL tree is a binary search tree with the additional balance property that, for any node in the tree, the height of the left and right subtrees can differ by at most 1. (Height of an empty subtree is  $-1$ ).

# AN AVL TREE

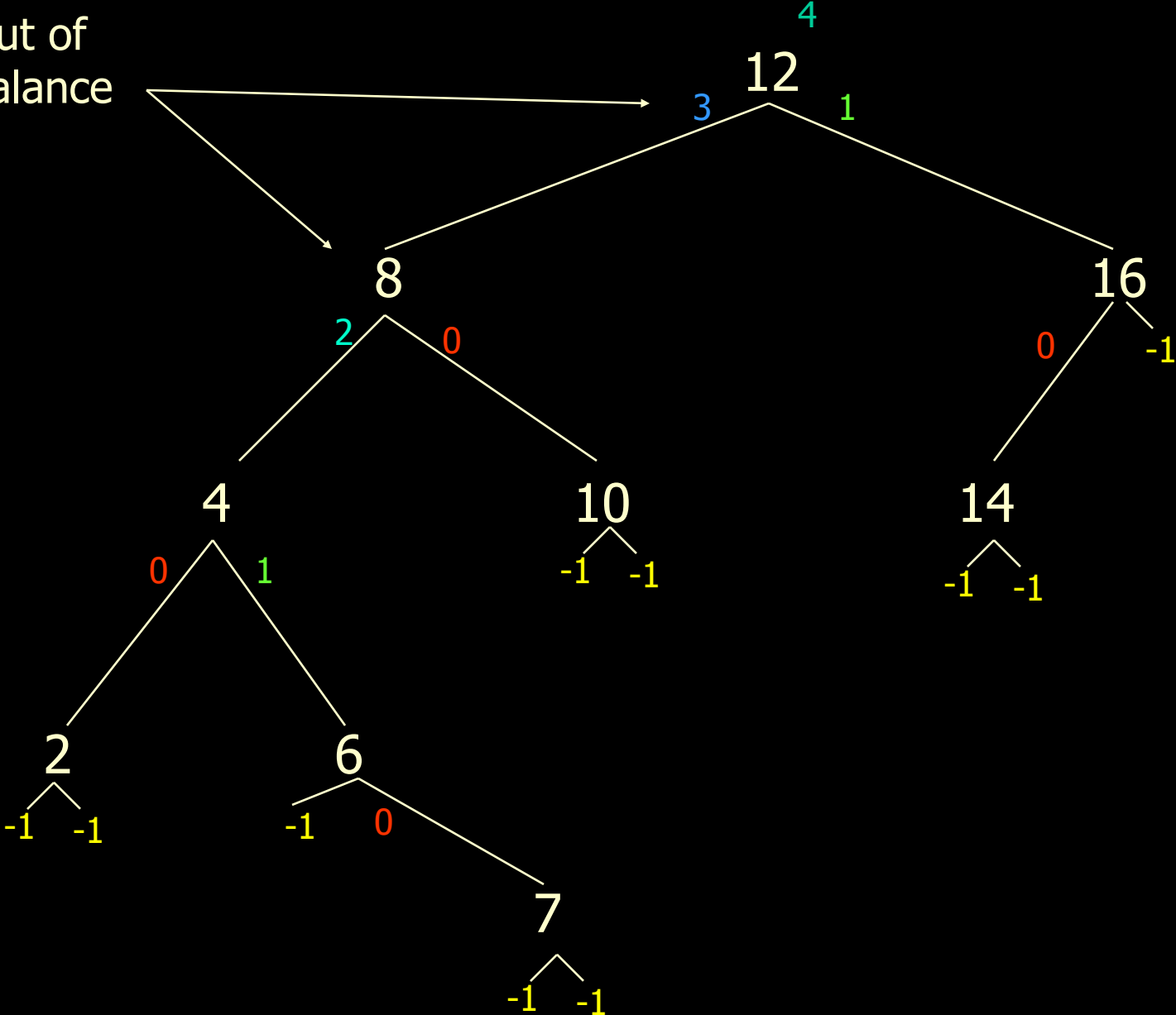
What is the height?



What if we insert "7"?

# NOT AN AVL TREE

Out of balance



# AVL THEOREM

An AVL tree of  $n$  element has height  $O(\log n)$

In fact, An AVL tree of height  $H$  has at least  $F_{H+3} - 1$  nodes, where  $F_i$  is the  $i^{\text{th}}$  Fibonacci number



# AVL THEOREM

Let  $S_H$  be the size of smallest AVL tree of height  $H$ .

Clearly  $S_0=1$  and  $S_1=2$

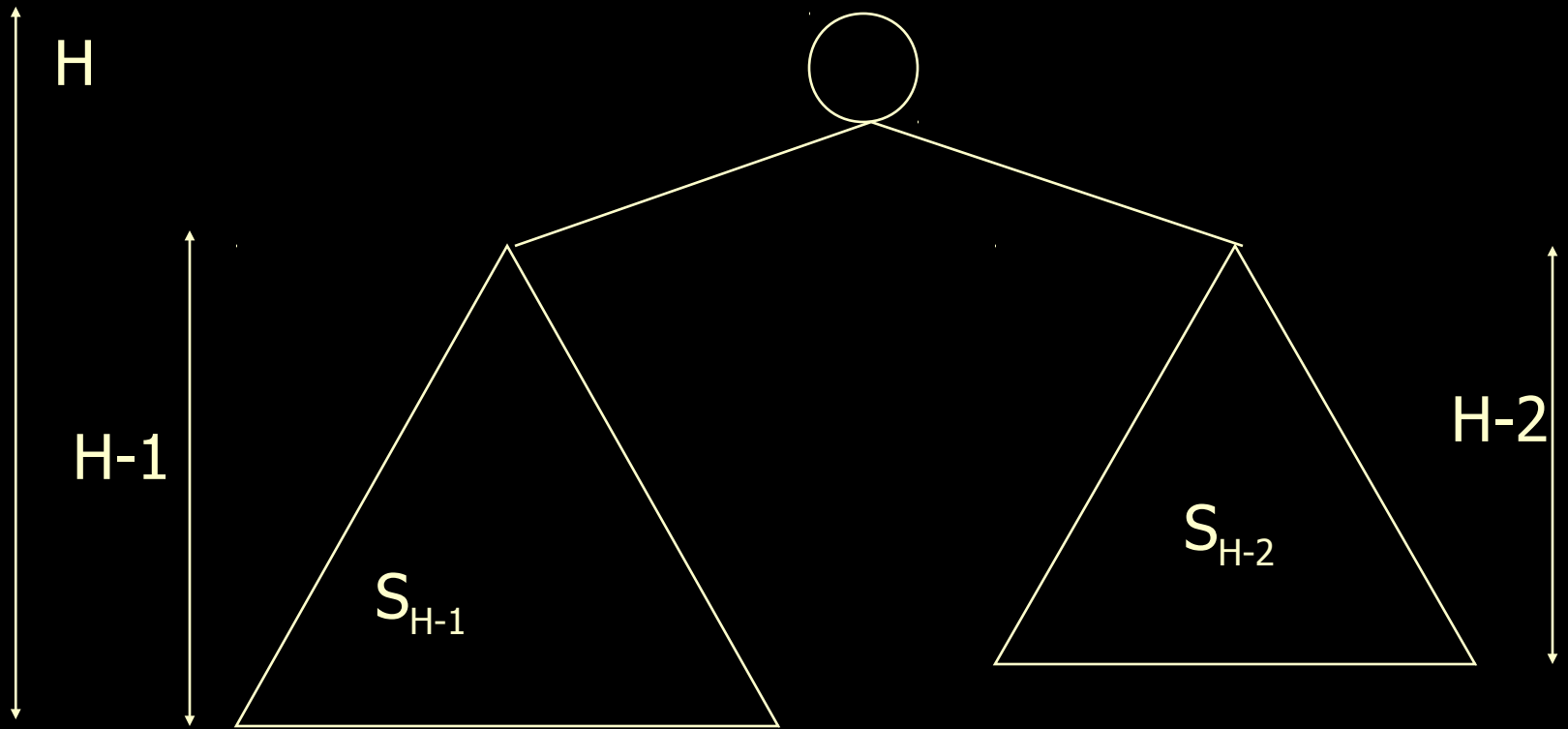
$S_H$  must have subtrees of height  $H-1$  and  $H-2$

These subtrees must have fewest nodes for height

$$\text{So, } S_H = S_{H-1} + S_{H-2} + 1$$

Minimum number AVL trees are *Fibonacci trees*

# Smallest AVL Tree of height H



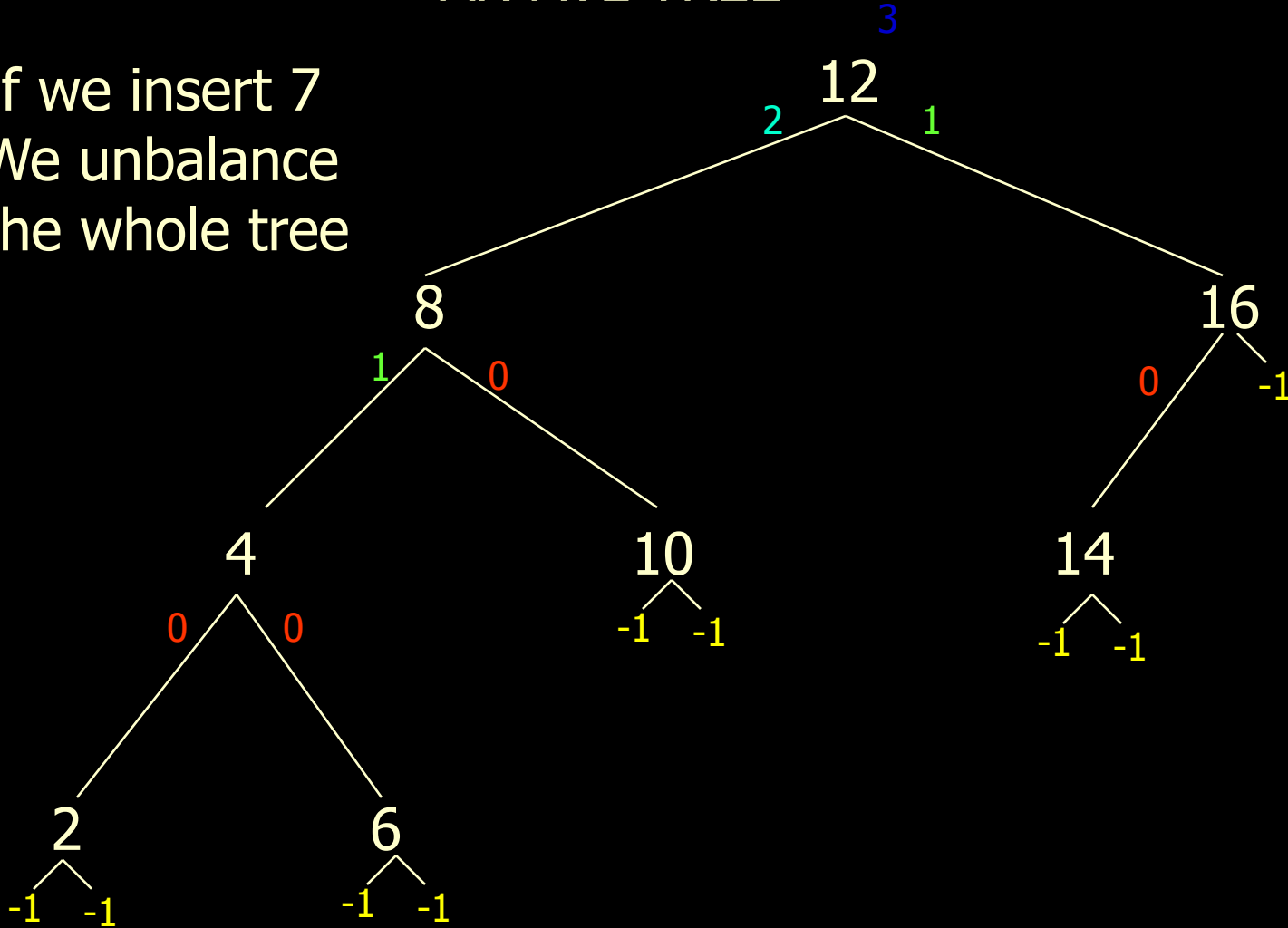
# Adds can unbalance a BST

Only nodes on the path from the root to the insertion point can have their balances altered

If we restore the unbalance node, we balance the tree

# AN AVL TREE

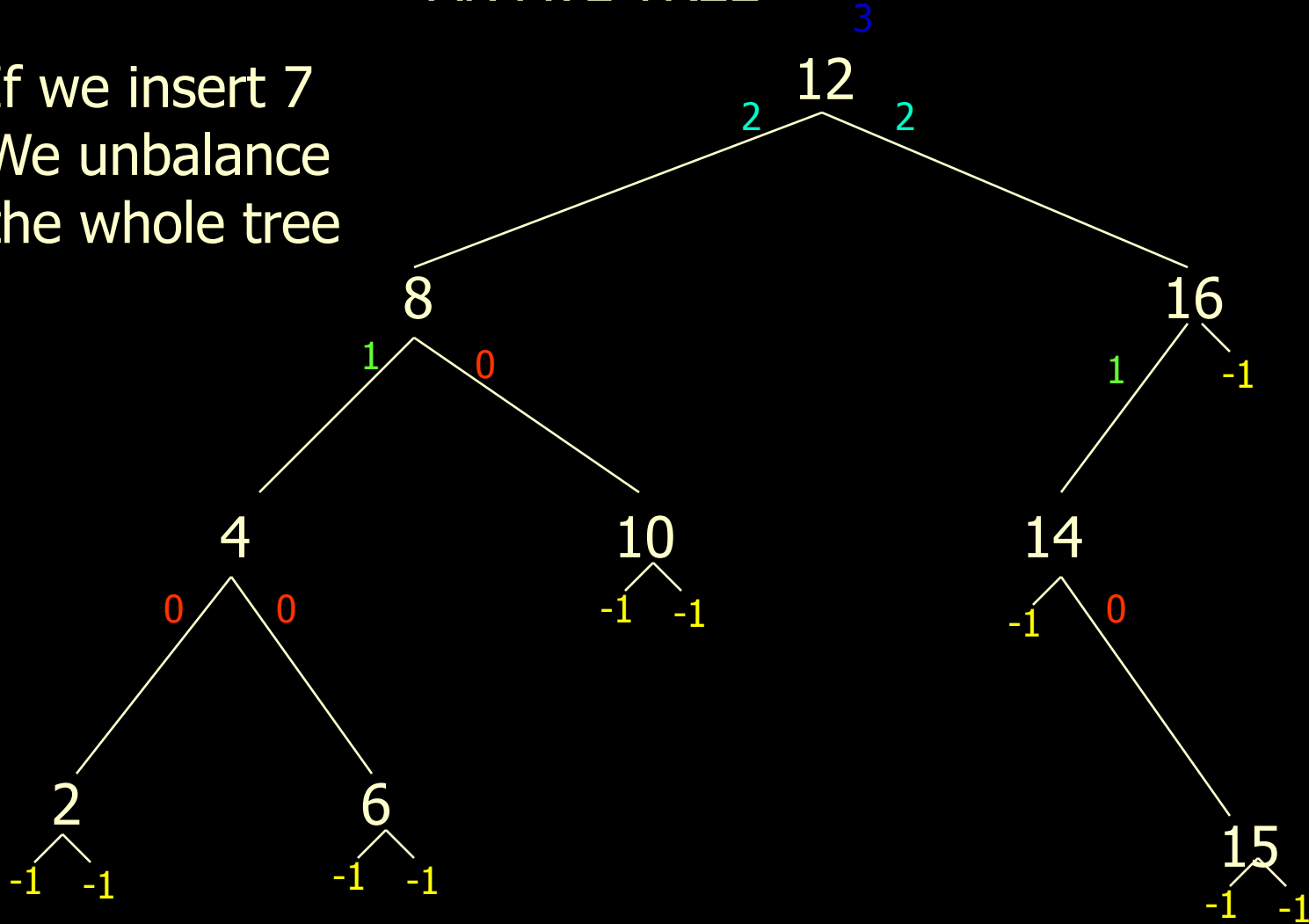
If we insert 7  
We unbalance  
the whole tree



What if we insert "15"?

# AN AVL TREE

If we insert 7  
We unbalance  
the whole tree



What if we insert "15"?

# 4 cases

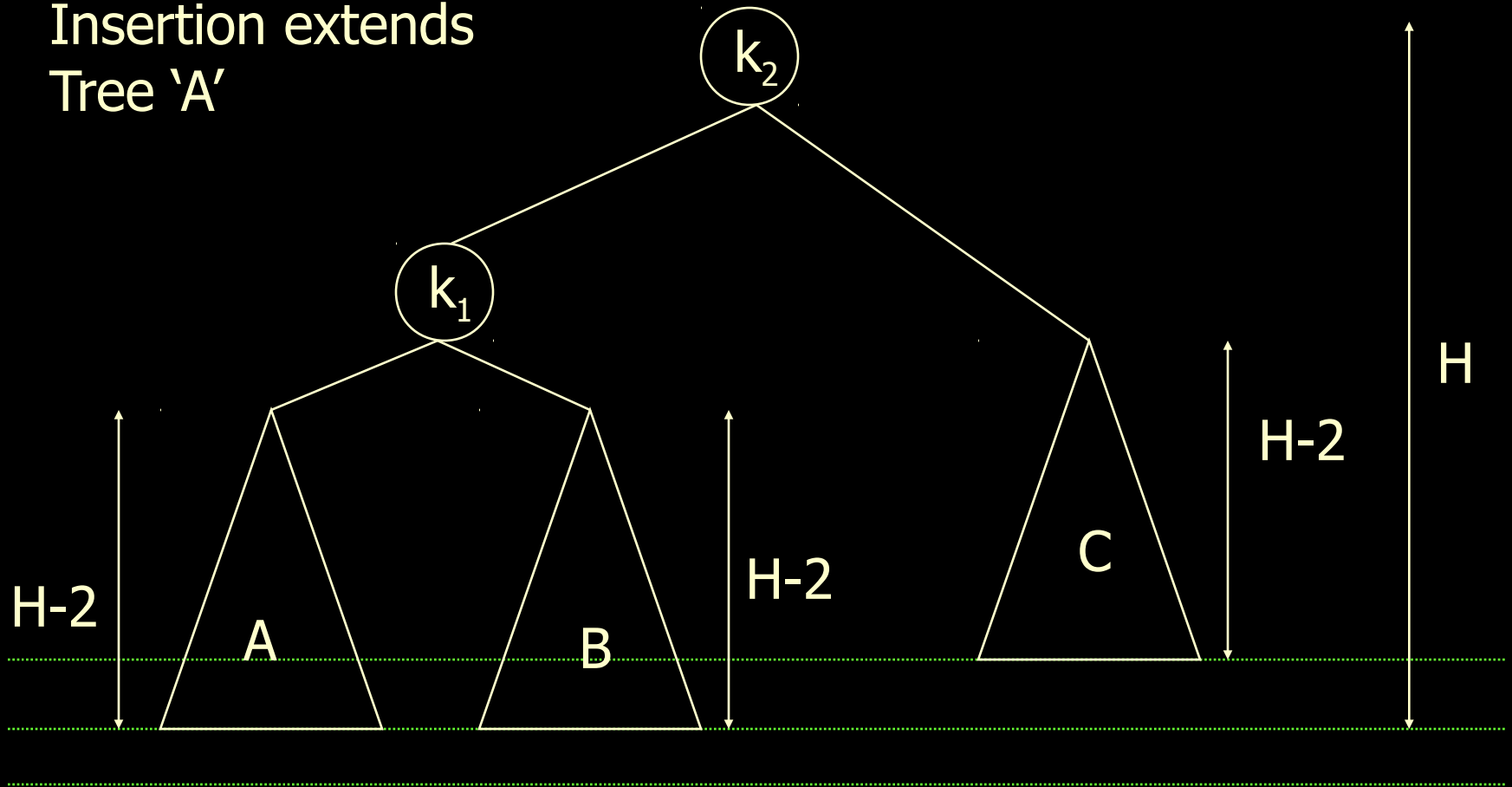
1. Insertion in left sub-tree of left child
2. Insertion in right sub-tree of left child
3. Insertion in left sub-tree of right child
4. Insertion in right sub-tree of right child

1 & 4 are symmetric

2 & 3 are symmetric

# Cases 1 & 4

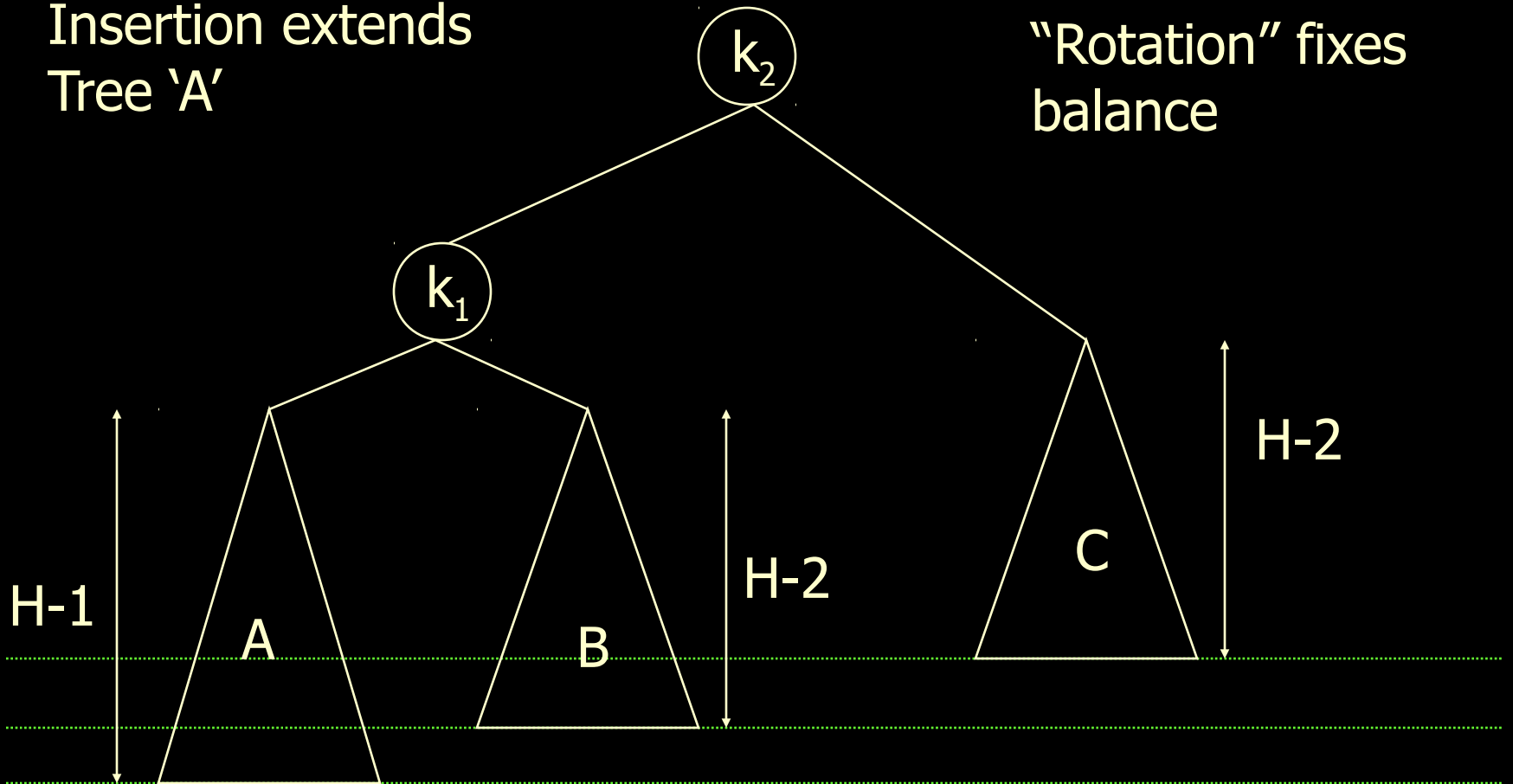
Insertion extends  
Tree 'A'



# Cases 1 & 4

Insertion extends  
Tree 'A'

"Rotation" fixes  
balance

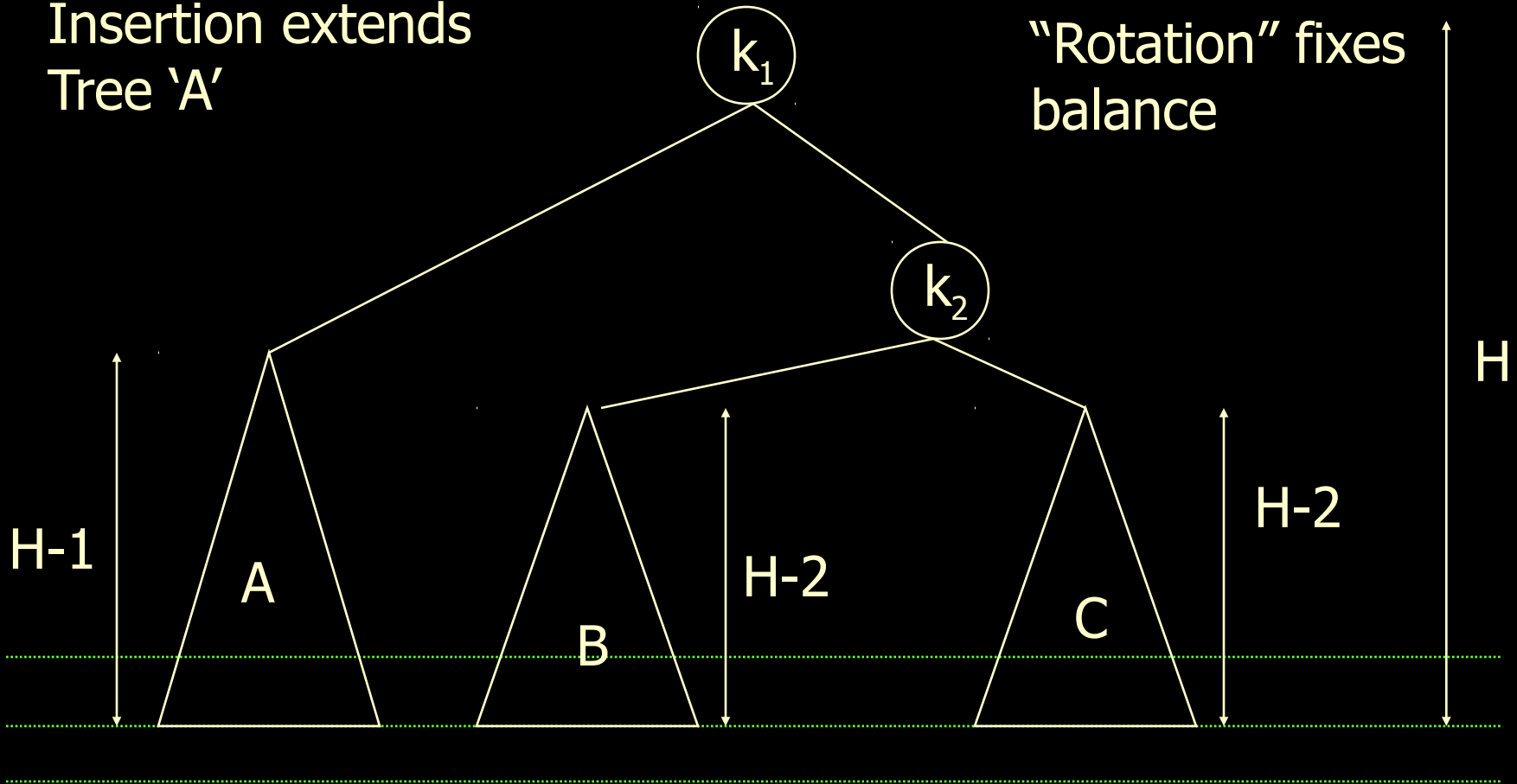




# Cases 1 & 4

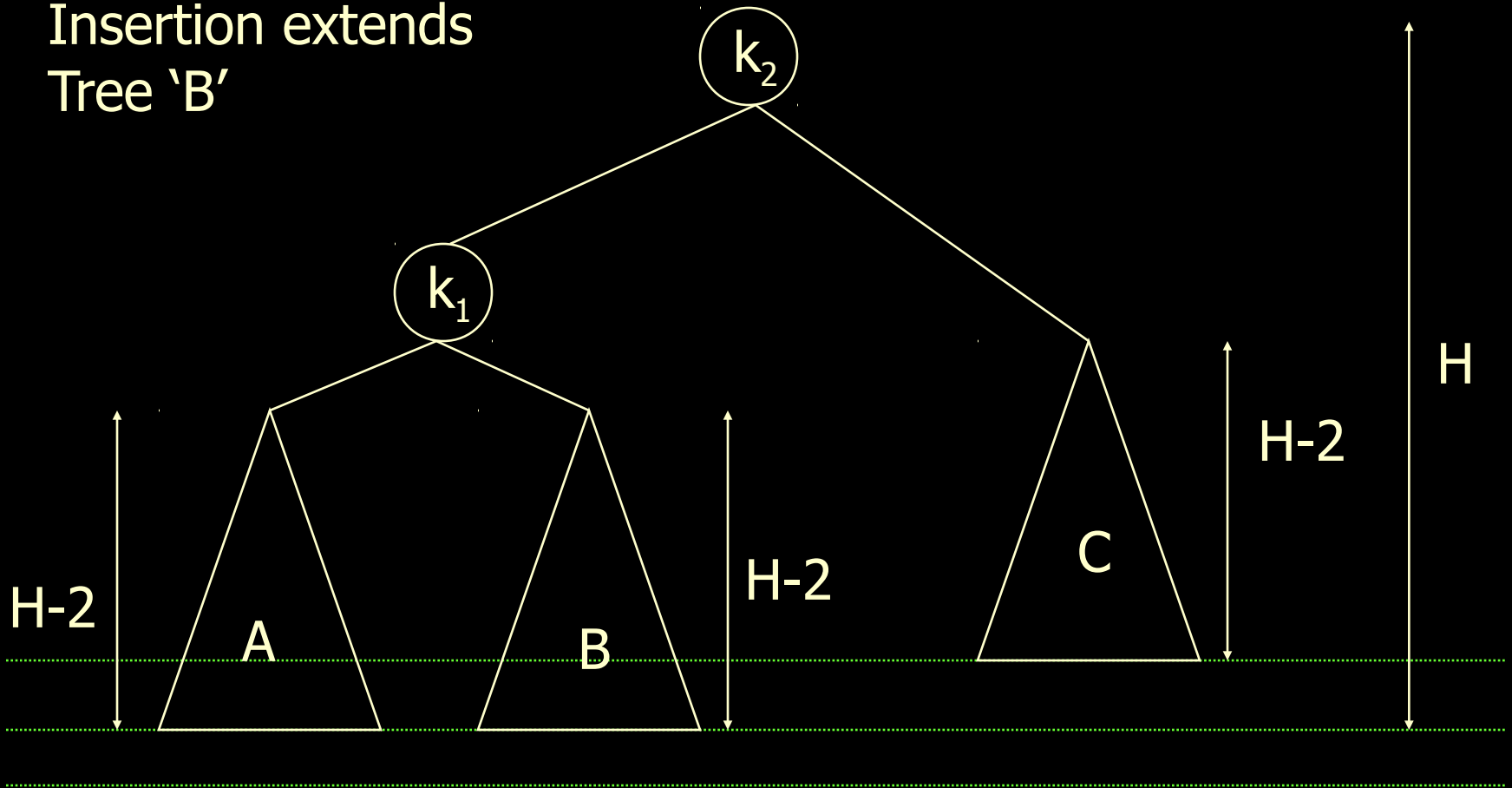
Insertion extends  
Tree 'A'

"Rotation" fixes  
balance



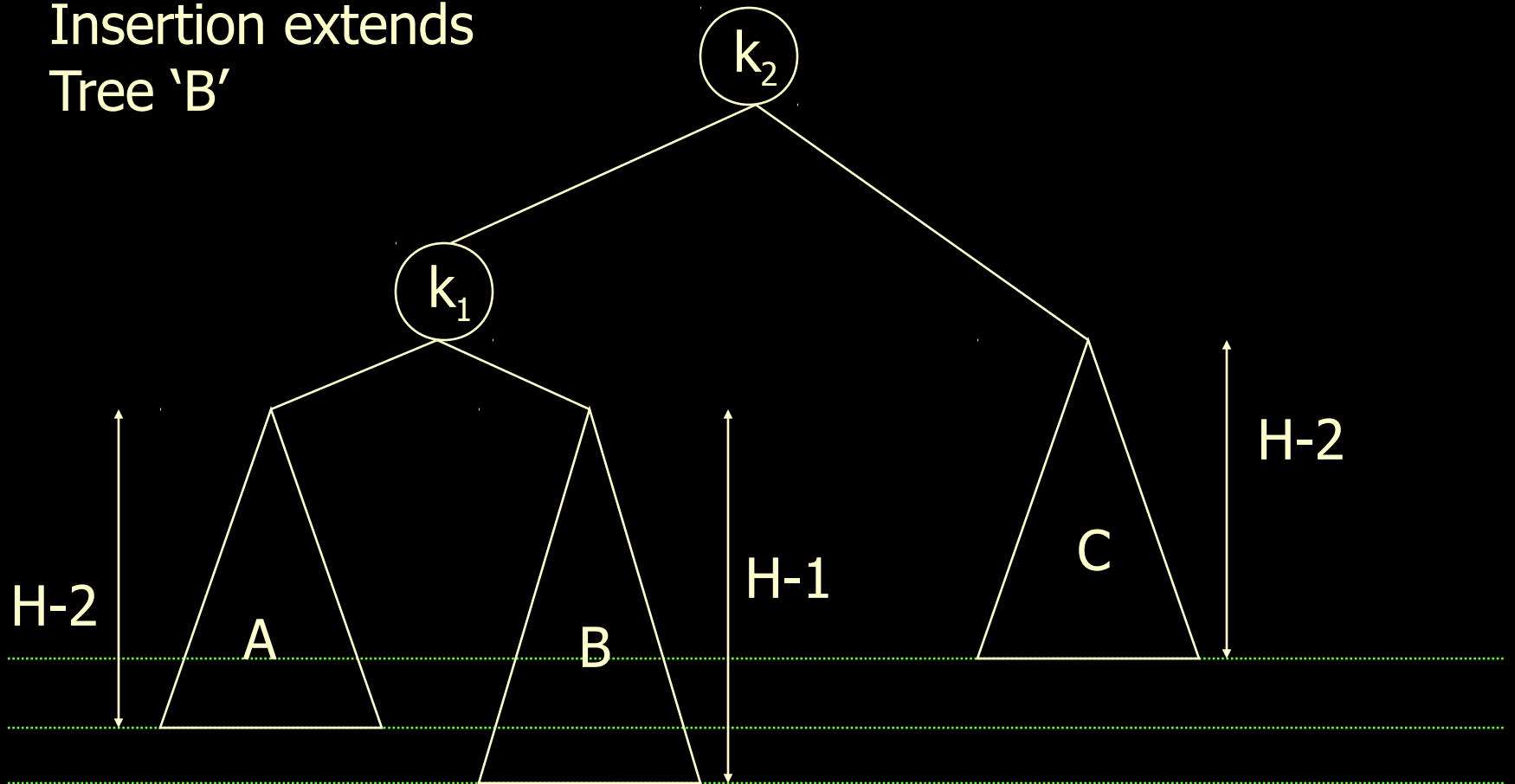
# Cases 2 & 3

Insertion extends  
Tree 'B'



# Cases 2 & 3

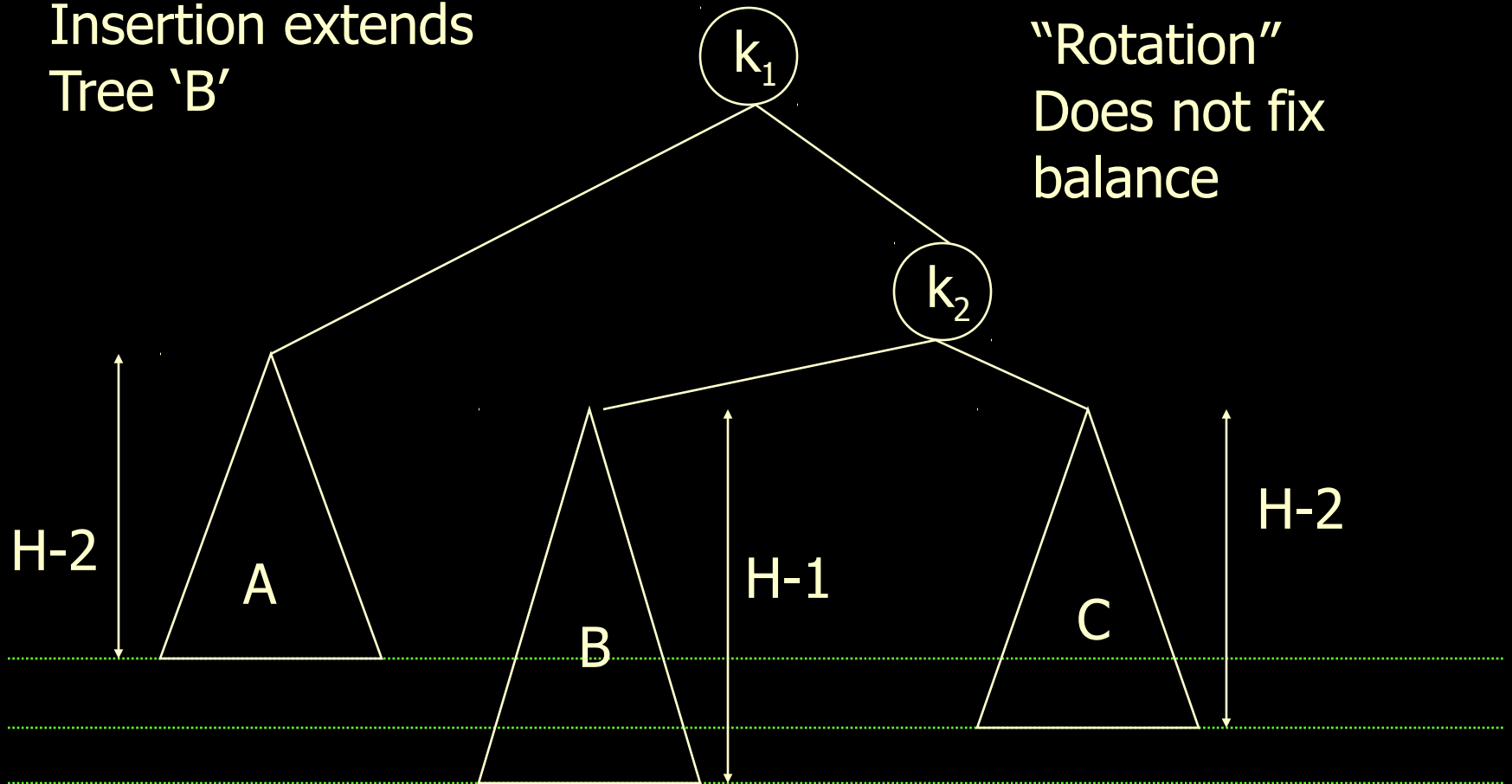
Insertion extends  
Tree 'B'



# Cases 2 & 3

Insertion extends  
Tree 'B'

"Rotation"  
Does not fix  
balance



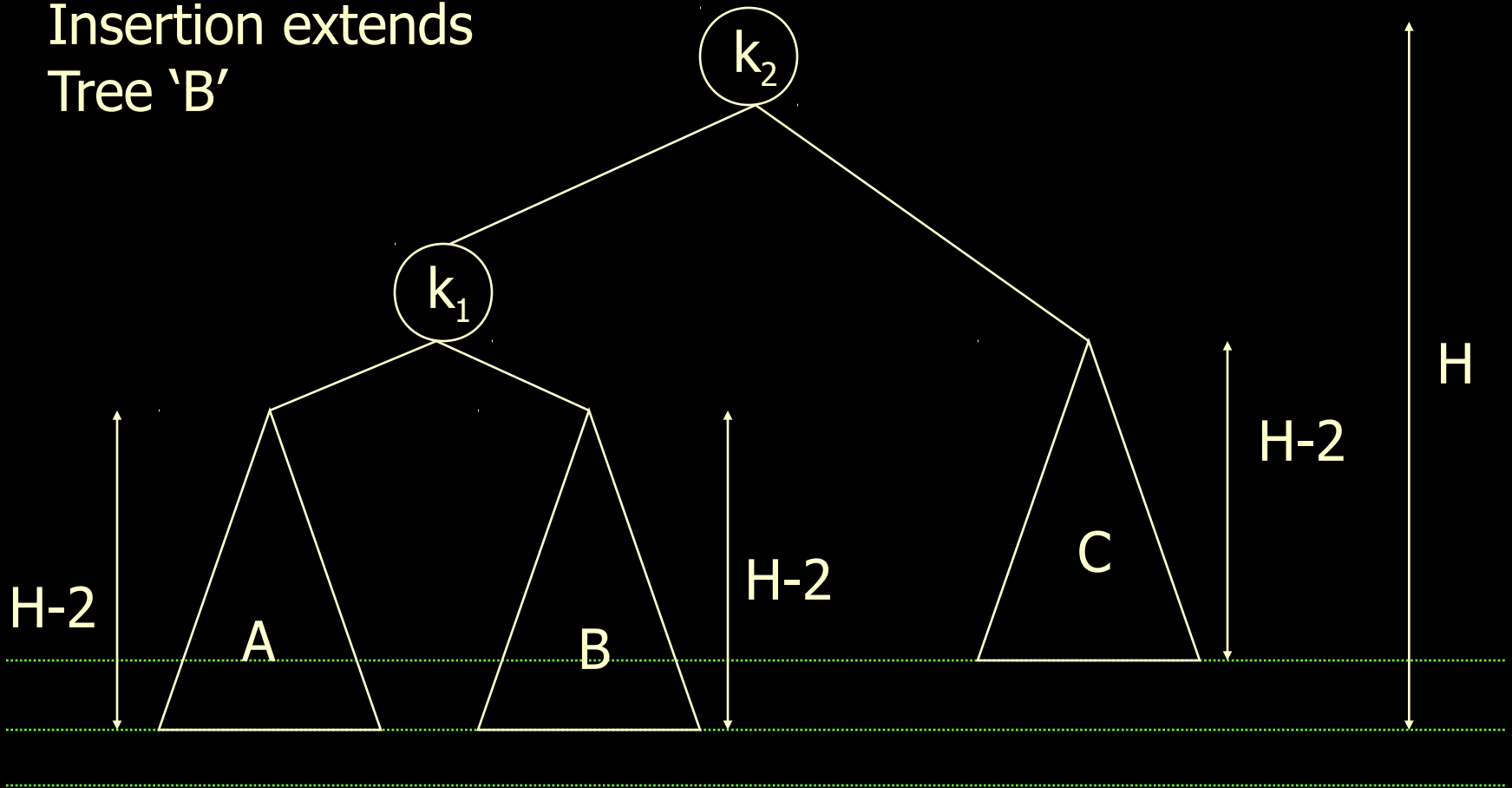
# Double Rotation

If x is out of balance for cases 2&3

1. Rotate between X's child and grandchild
2. Rotate between X and it's new child

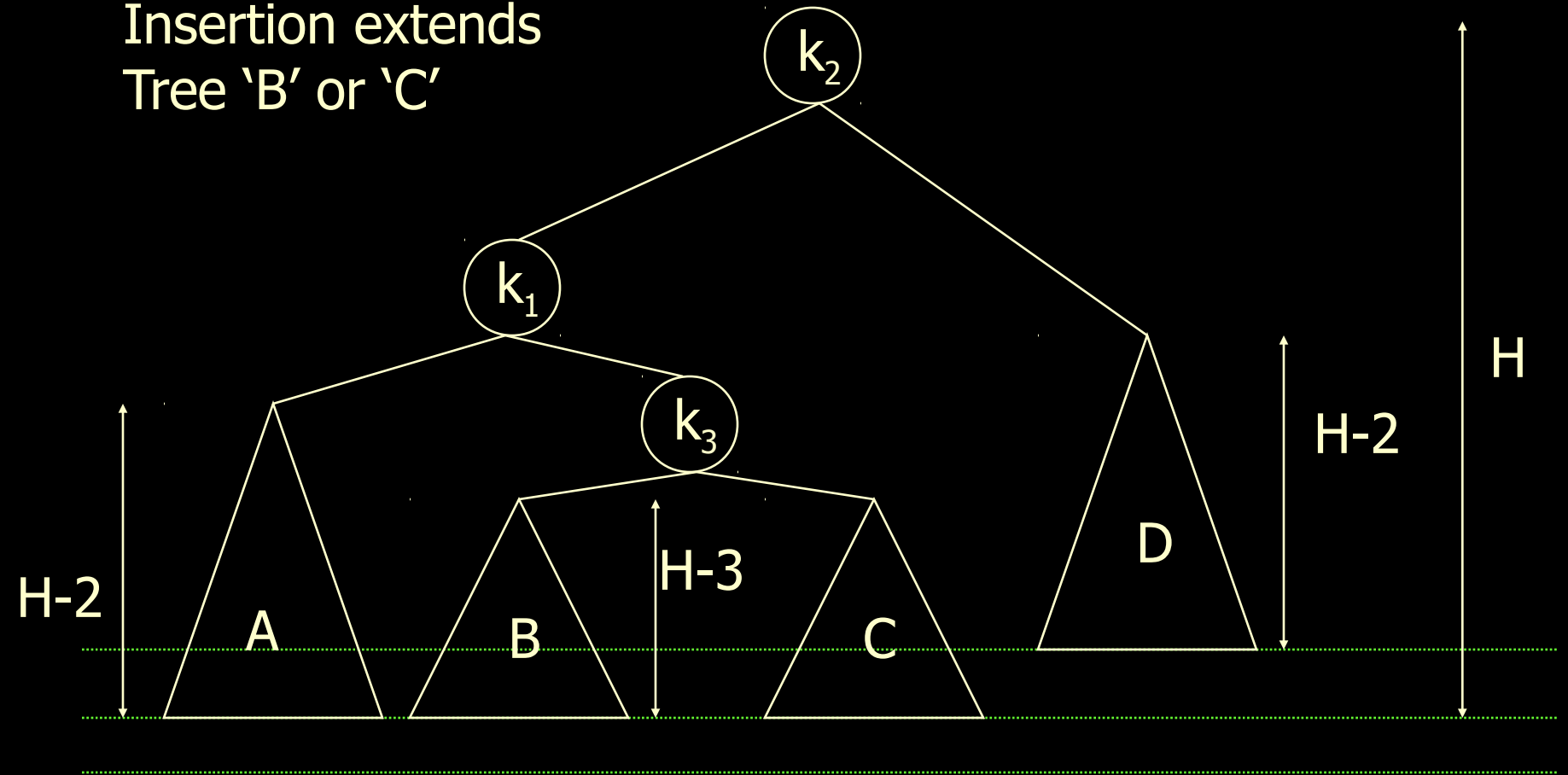
# Cases 2 & 3

Insertion extends  
Tree 'B'



# Cases 2 & 3

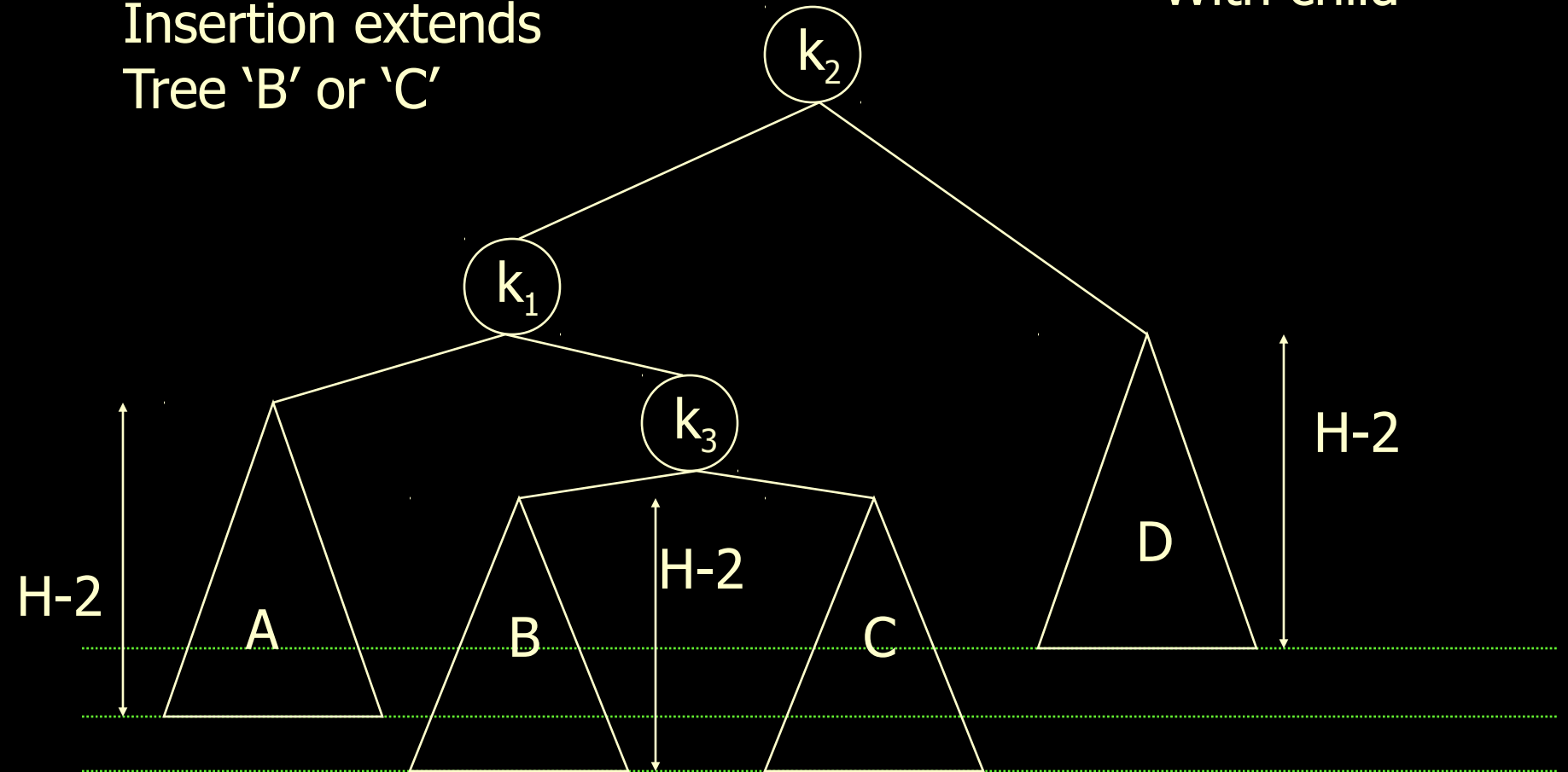
Insertion extends  
Tree 'B' or 'C'



# Cases 2 & 3

Rotate grandchild  
With child

Insertion extends  
Tree 'B' or 'C'



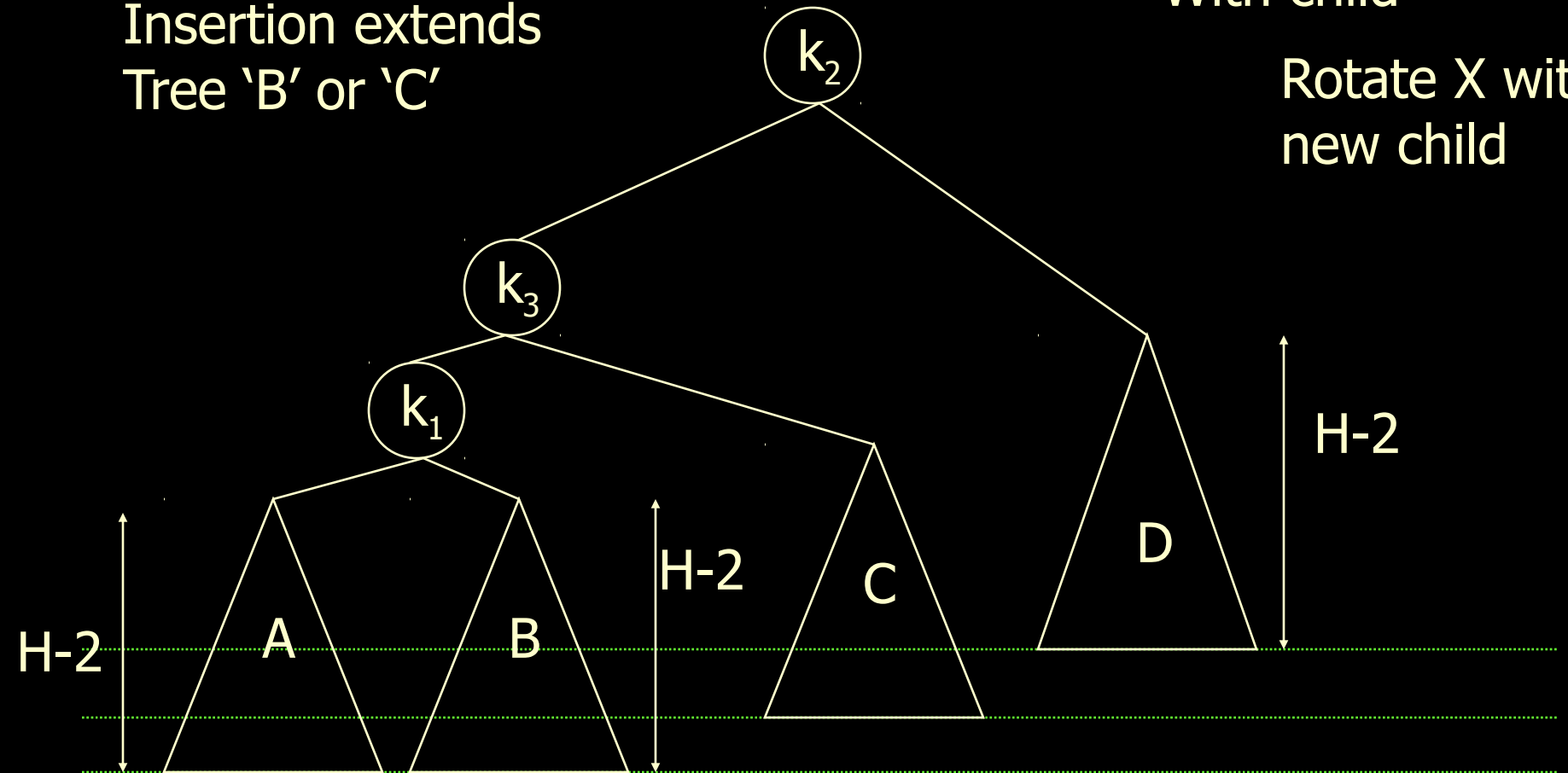


# Cases 2 & 3

Rotate grandchild  
With child

Rotate X with  
new child

Insertion extends  
Tree 'B' or 'C'

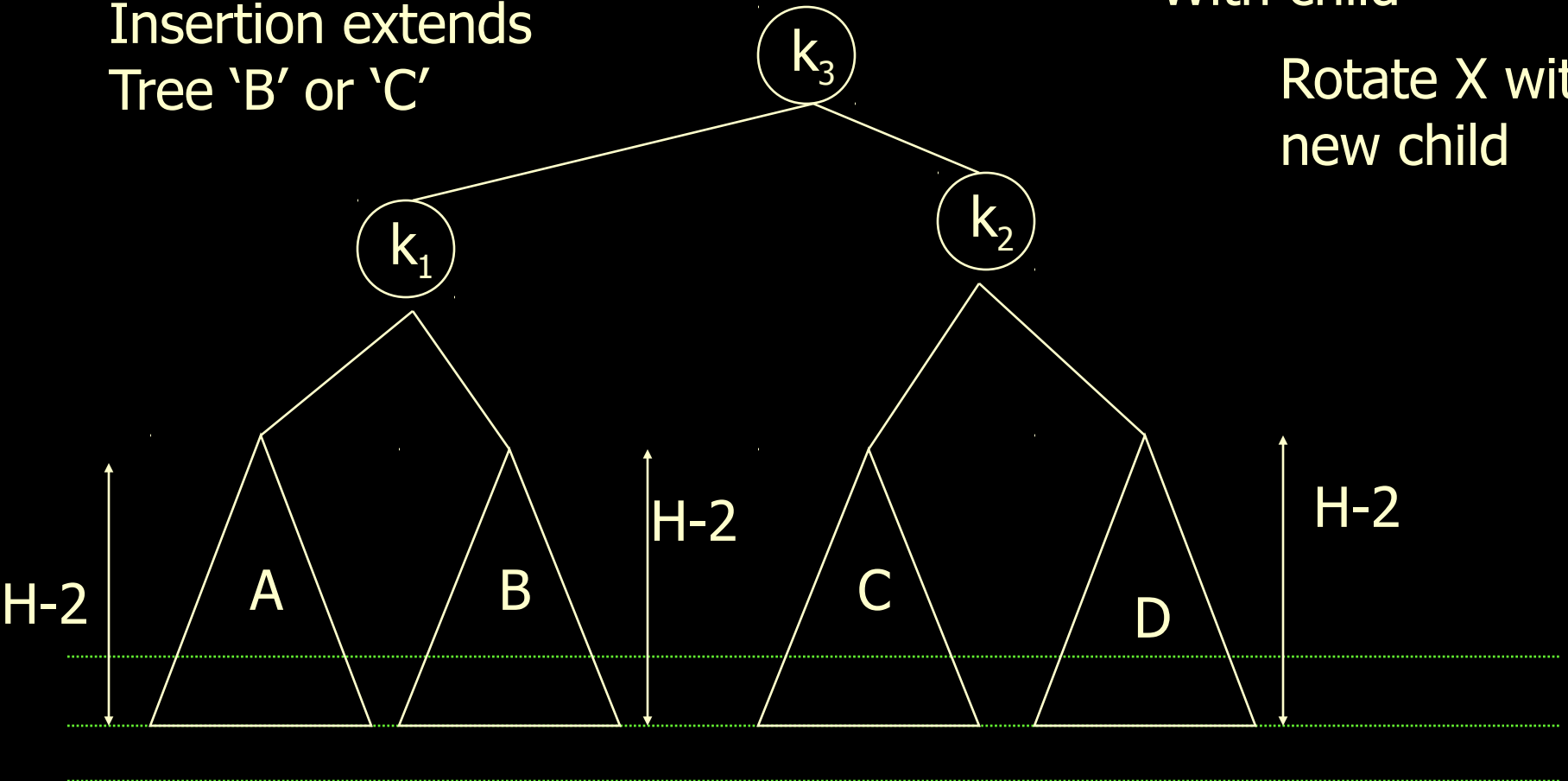


# Cases 2 & 3

Rotate grandchild  
With child

Insertion extends  
Tree 'B' or 'C'

Rotate X with  
new child



# Implementation

## Do this in lab

Insert

Fixup

Rotate

Need to keep track of balance

# Rotation

```
private void rotateLeft(Node p) {  
    Node r = p.right;  
    p.right = r.left;  
    if (r.left != null) r.left.parent = p;  
    r.parent = p.parent;  
    if (p.parent == null) {root = r; r.parent = null;}  
    else if (p.parent.left == p) p.parent.left = r;  
    else p.parent.right = r;  
    r.left = p;  
    p.parent = r;  
}
```

root

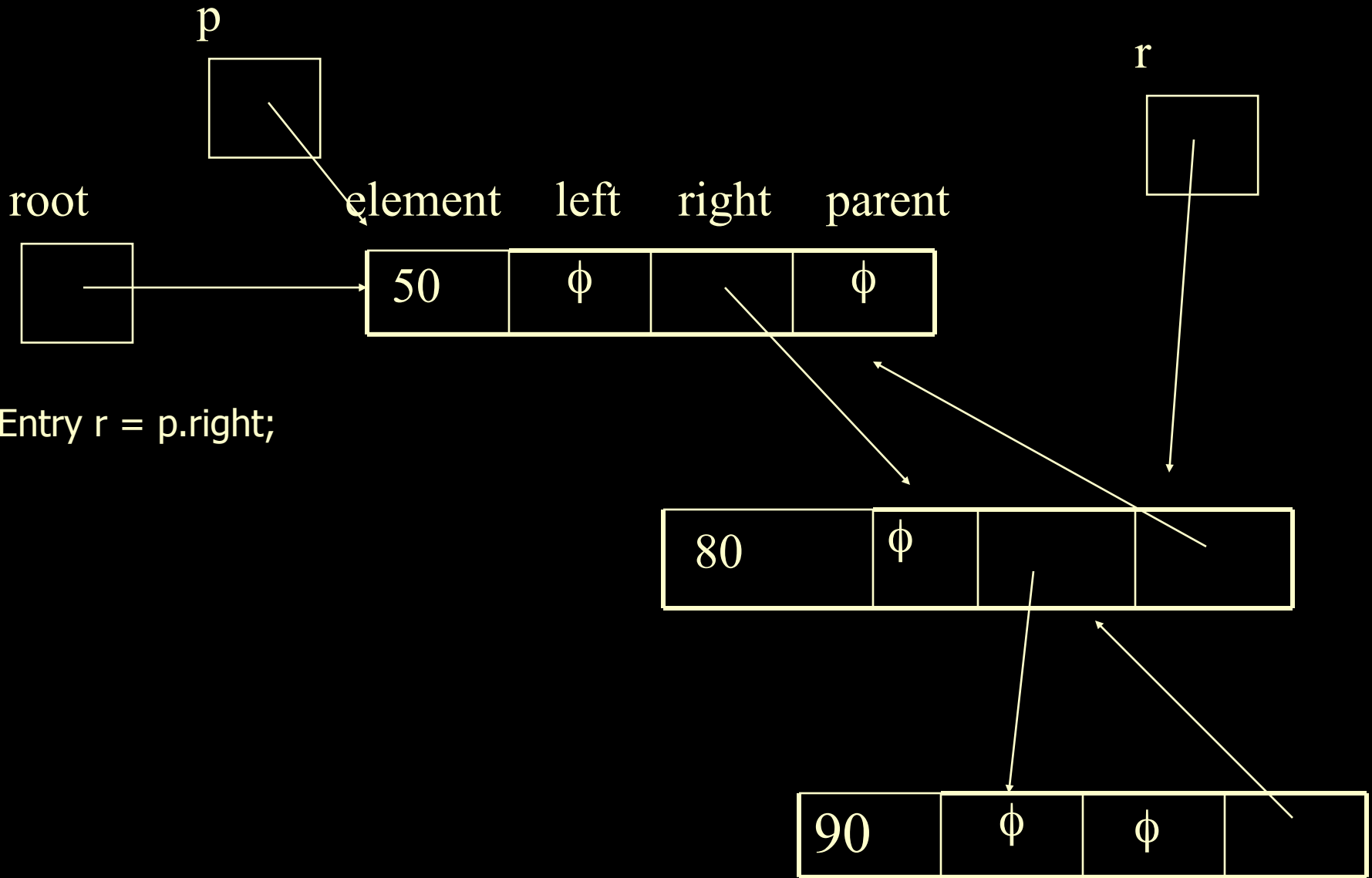
element

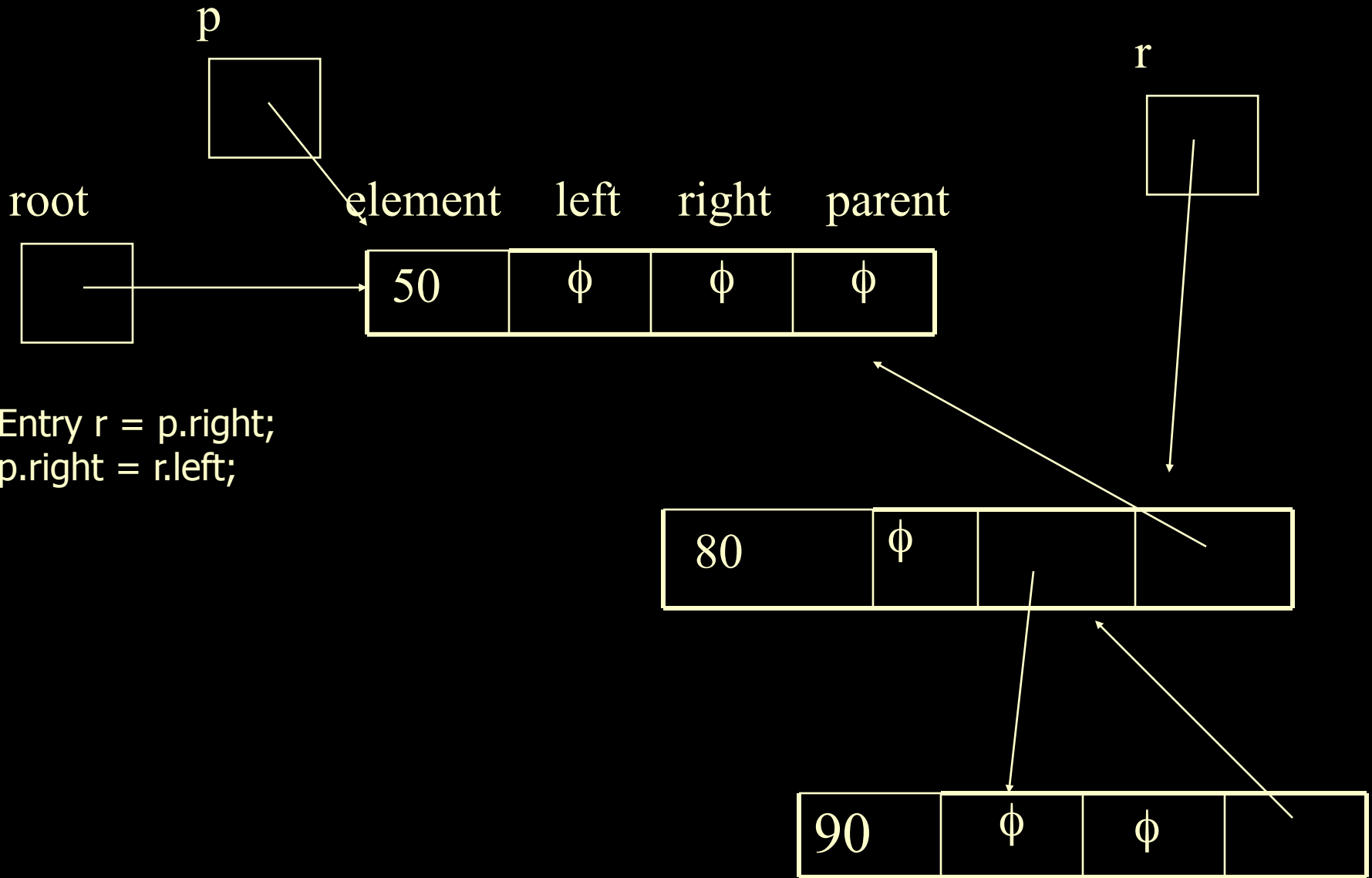
left

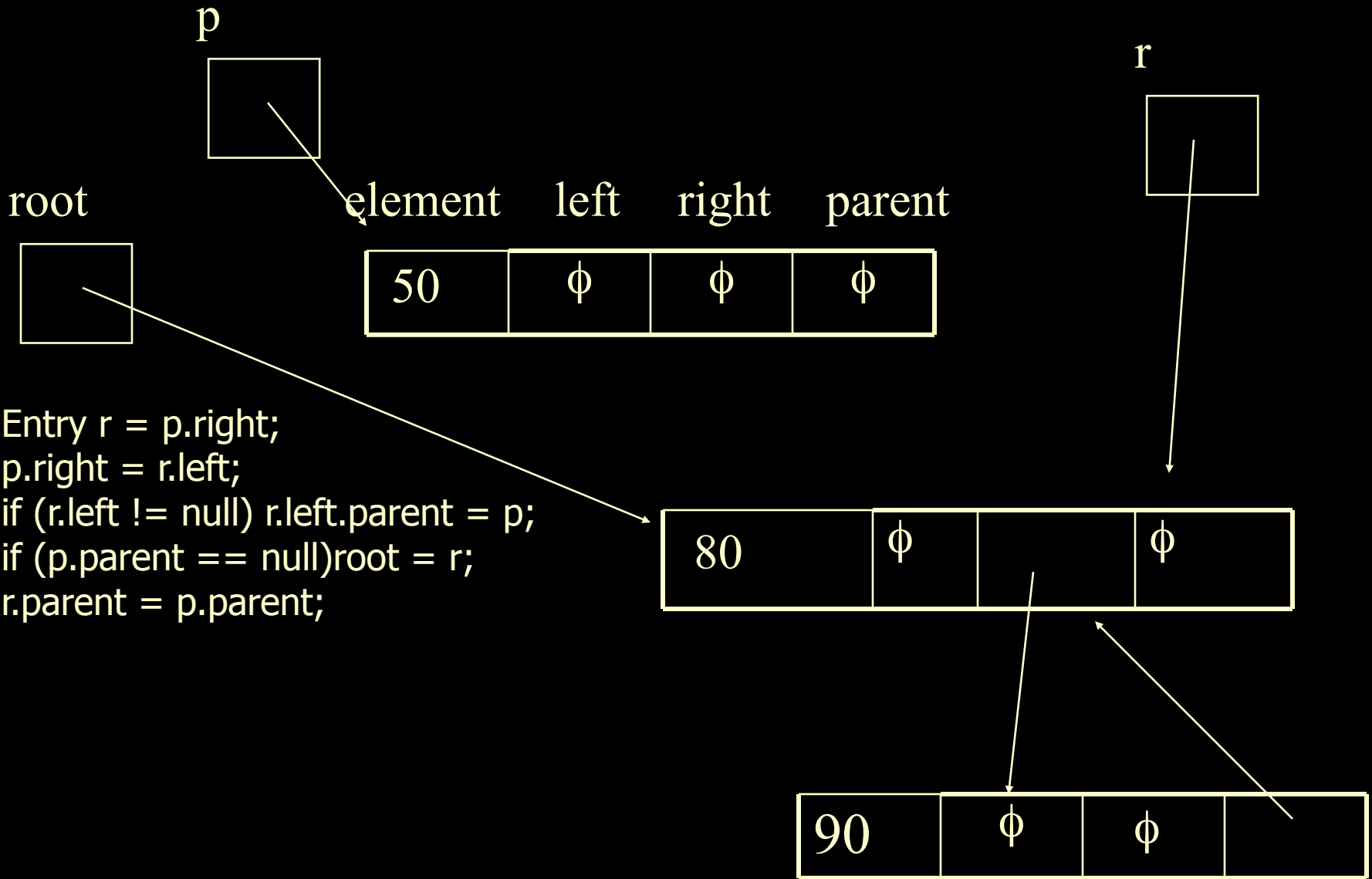
right

parent

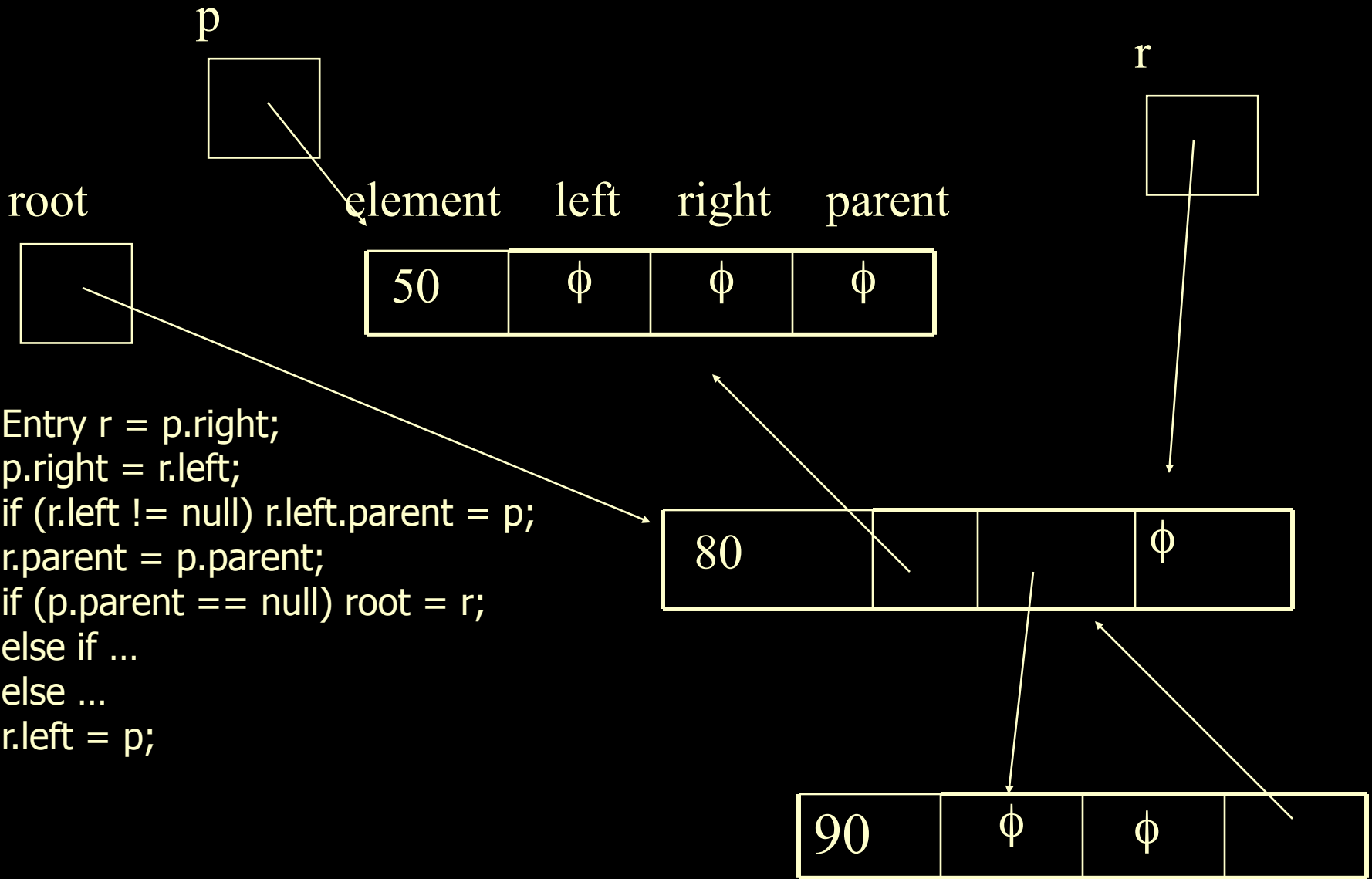


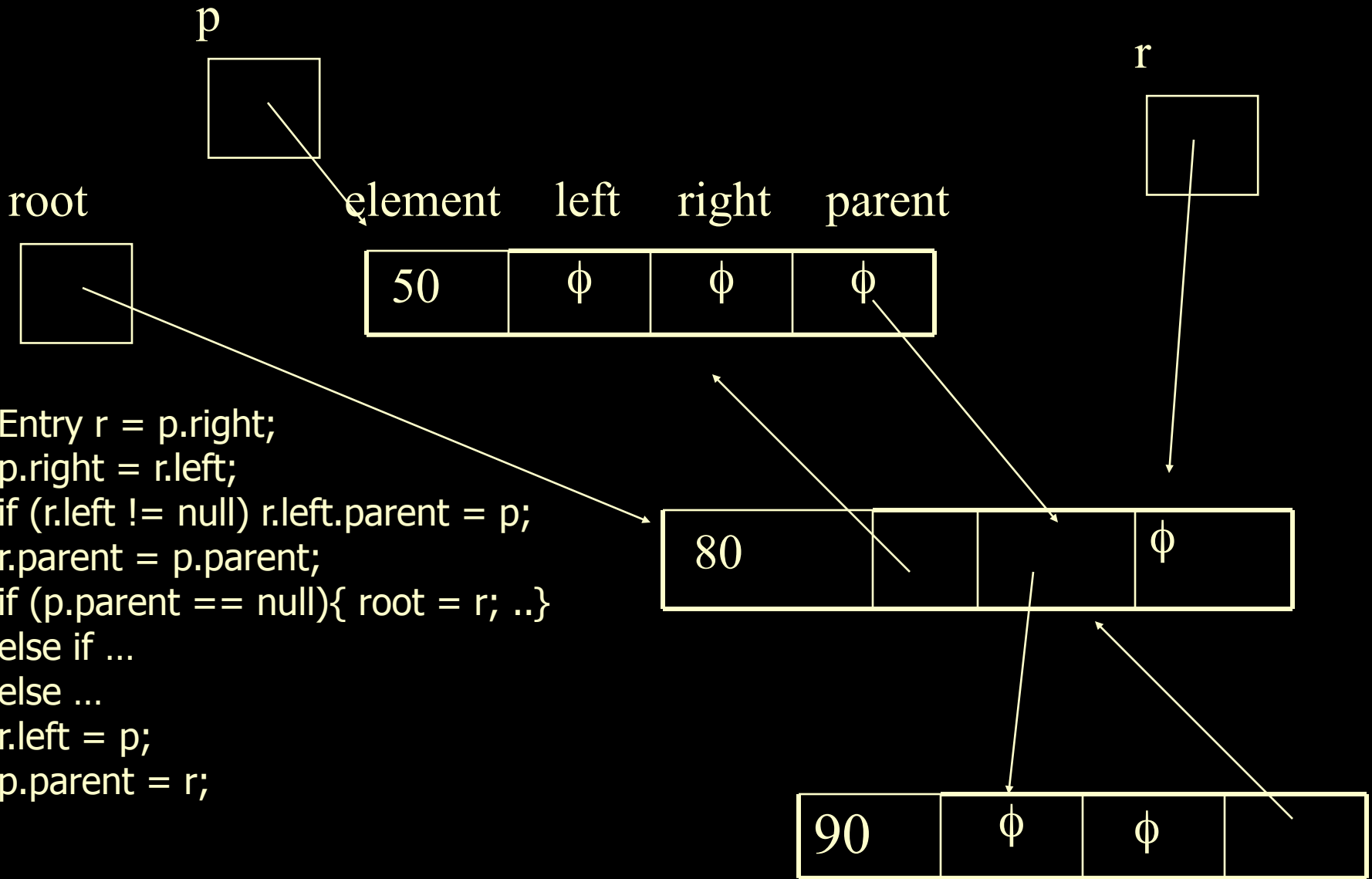












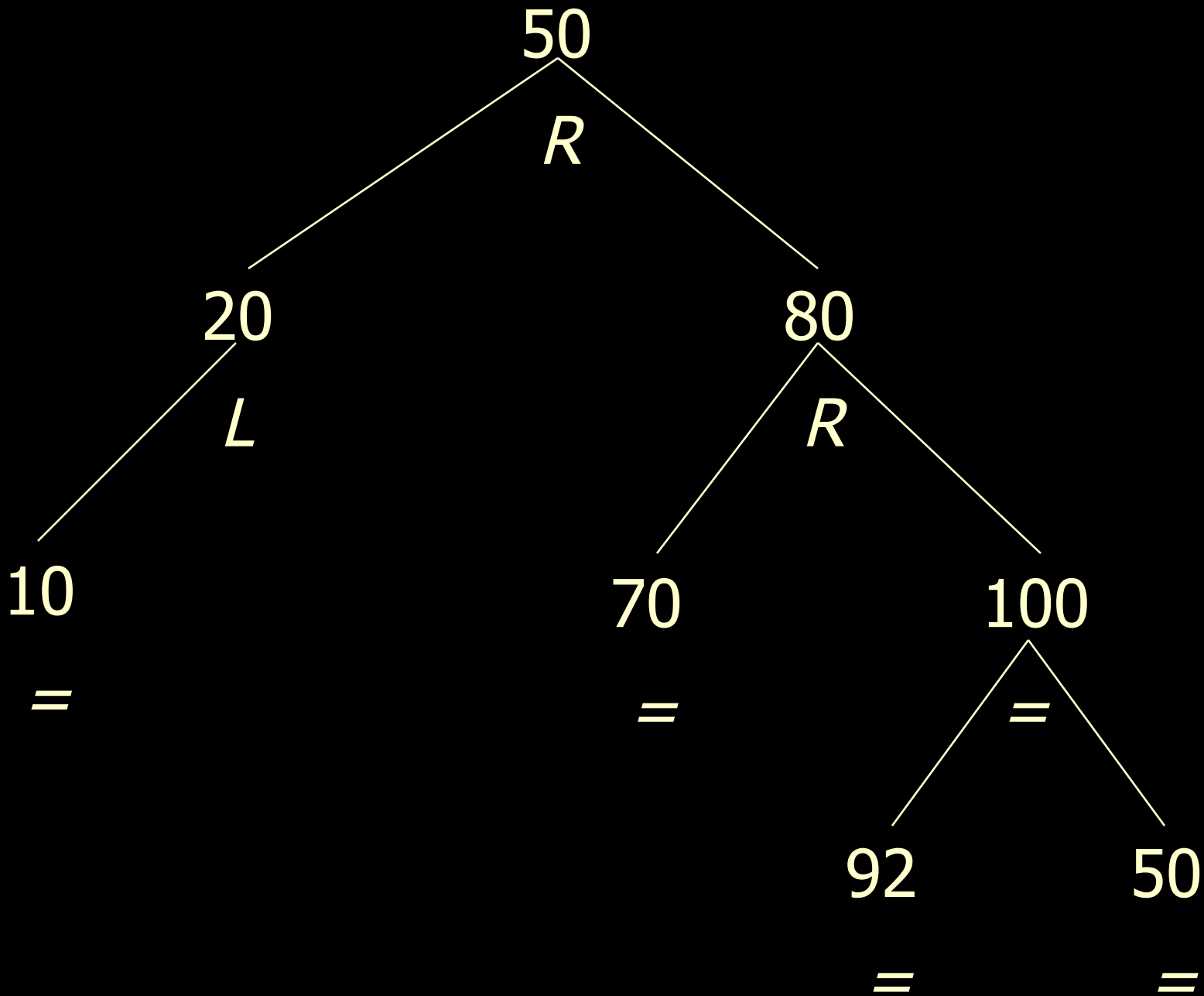
```

Entry r = p.right;
p.right = r.left;
if (r.left != null) r.left.parent = p;
r.parent = p.parent;
if (p.parent == null){ root = r; ..}
else if ...
else ...
r.left = p;
p.parent = r;

```

# Node class

```
private static class Node {  
    Object element;  
    char balanceFactor = '='; // new nodes are balanced  
    // we could set this to R or L indicating child with > height  
    Node left = null, right = null, parent;  
    Node (Object element, Node parent) {  
        this.element = element;  
        this.parent = parent;  
    }  
}
```



```
public boolean add(Object o){
    if (root == null) {
        root = new Node(o,null);
        size++;
        return true;
    } // empty tree
    else {
        Node temp = root,
        ancestor = null; // we keep track of nearest unbalanced ancestor
        int comp;
        while (true) {
            comp = ((Comparable)o).compareTo(temp.element);
            if (comp == 0) return false;
```

```
if (comp < 0) {
    if (temp.balanceFactor != '=') ancestor = temp;
    if (temp.left != null) temp = temp.left;
    else {
        temp.left = new Node(o,temp);
        fixAfterInsertion(ancestor,temp.left);
        size++;
    }
} // comp < 0
```

```
else { // comp > 0
    if (temp.balanceFactor != '=') ancestor = temp;
    if (temp.right != null) temp = temp.right;
    else {
        temp.rig = new Node(o,temp);
        fixAfterInsertion(ancestor,temp.right);
        size++;
    }
}
```

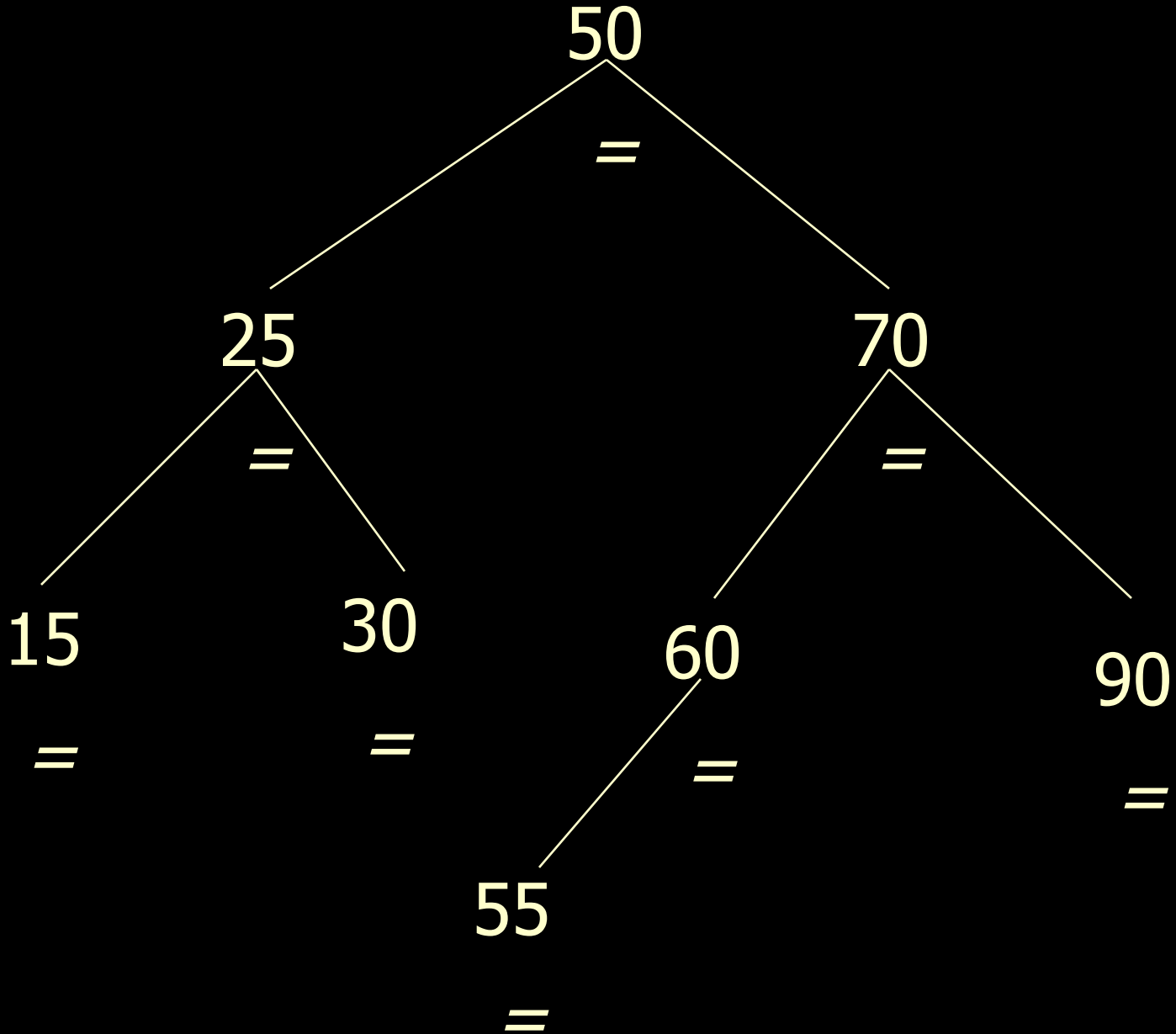
```
}// comp < 0
```

```
}// while
```

```
}// root not null
```

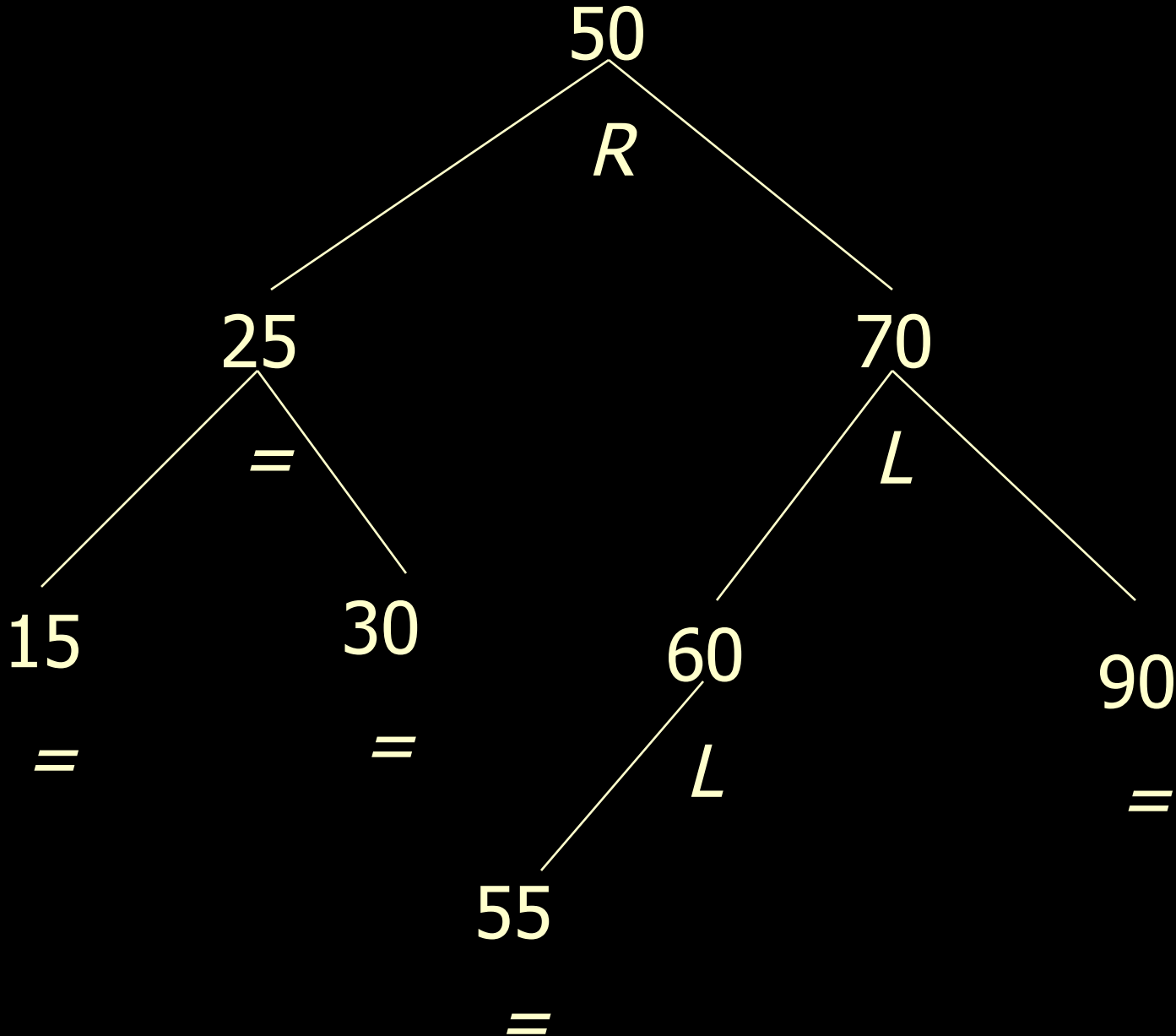
```
}//method add
```

# Adjusting paths





# Adjusting paths



```
protected void adjustPath(Entry to, Entry inserted) {
    Object o = inserted.element;
    Node temp = inserted.parent;
    while (temp != to) {
        if (((Comparable)o).compareTo(temp.element)<0)
            temp.balanceFactor = 'L';
        else
            temp.balanceFactor = 'R';
        temp = temp.parent
    } // while
} //adjust path
```

```
protected void fixAfterInsertion(Node ancestor, Entry inserted) {
    Object o = inserted element;
    if (ancestor == null) {
        if (((Comparable)o).compareTo(root.element)<0)
            root.balanceFactor = 'L';
        else
            root.balanceFactor = 'R';
        adjustPath(root,inserted);
    } // Case 1: all ancestor of inserted element have '=' balanceFactor
```

```
if ((ancestor.balanceFactor == 'L' &&
    ((Comparable)o).compareTo(ancestor.element)>0)||
    (ancestor.balanceFactor == 'R' &&
    ((Comparable)o).compareTo(ancestor.element)<0)) {
    ancestor.balanceFactor = '=';
    adjustPath(ancestor,inserted);
} // Case 2: insertion causes ancestor's balanceFactor to '='
```

```
if ((ancestor.balanceFactor == 'R' &&
    ((Comparable)o).compareTo(ancestor.right.element)>0) {
    ancestor.balanceFactor = '=';
    rotateLeft(ancestor);
    adjustPath(ancestor.parent,inserted);
} // Case 3: ancestor's balance factor = 'R'
//         and o > ancestor's right child
```

```
if ((ancestor.balanceFactor == 'L' &&
((Comparable)o).compareTo(ancestor.left.element)<0) {
    ancestor.balanceFactor = '=';
    rotateRight(ancestor);
    adjustPath(ancestor.parent,inserted);
} // Case 4: ancestor's balance factor = 'L'
//         and o < ancestor's right child
```

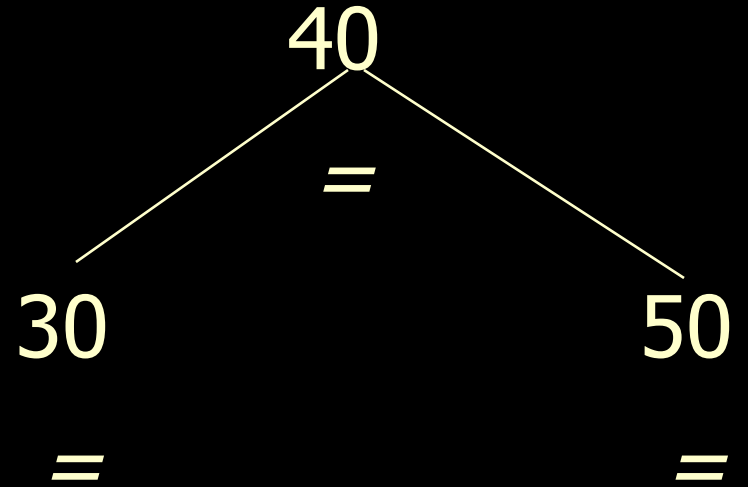
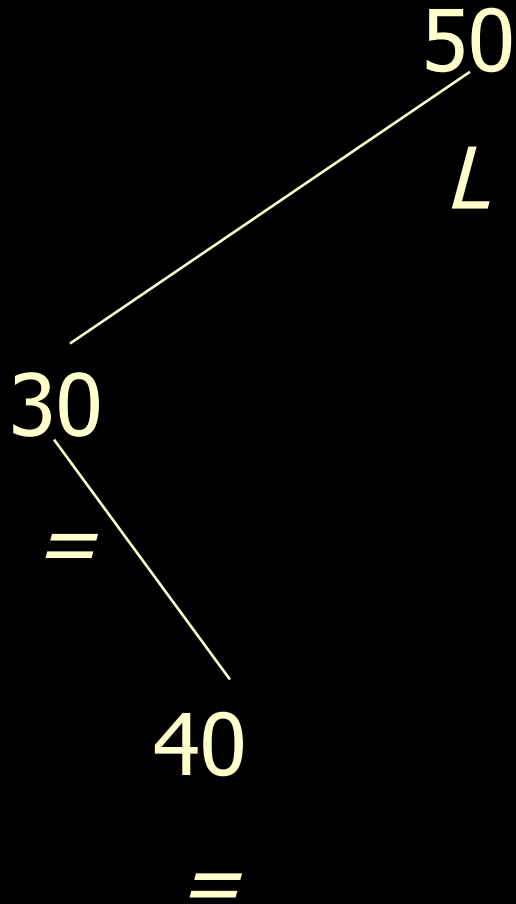
```
if (ancestor.balanceFactor == 'L' &&
    ((Comparable)o).compareTo(ancestor.left.element)>0) {
    rotateLeft(ancestor.left);
    rotateRight(ancestor);
    adjustLeftRight(ancestor,inserted);
} // Case 5: ancestor's balanceFactor = 'L'
//    and o > ancestor's left child
```

```
else {
    rotateRight(ancestor.right);
    rotateLeft(ancestor);
    adjustRightLeft(ancestor,inserted);
} // Case 6: ancestor's balanceFactor = 'R'
    // and o < ancestor's right child
} // fixAfterInsertion
```

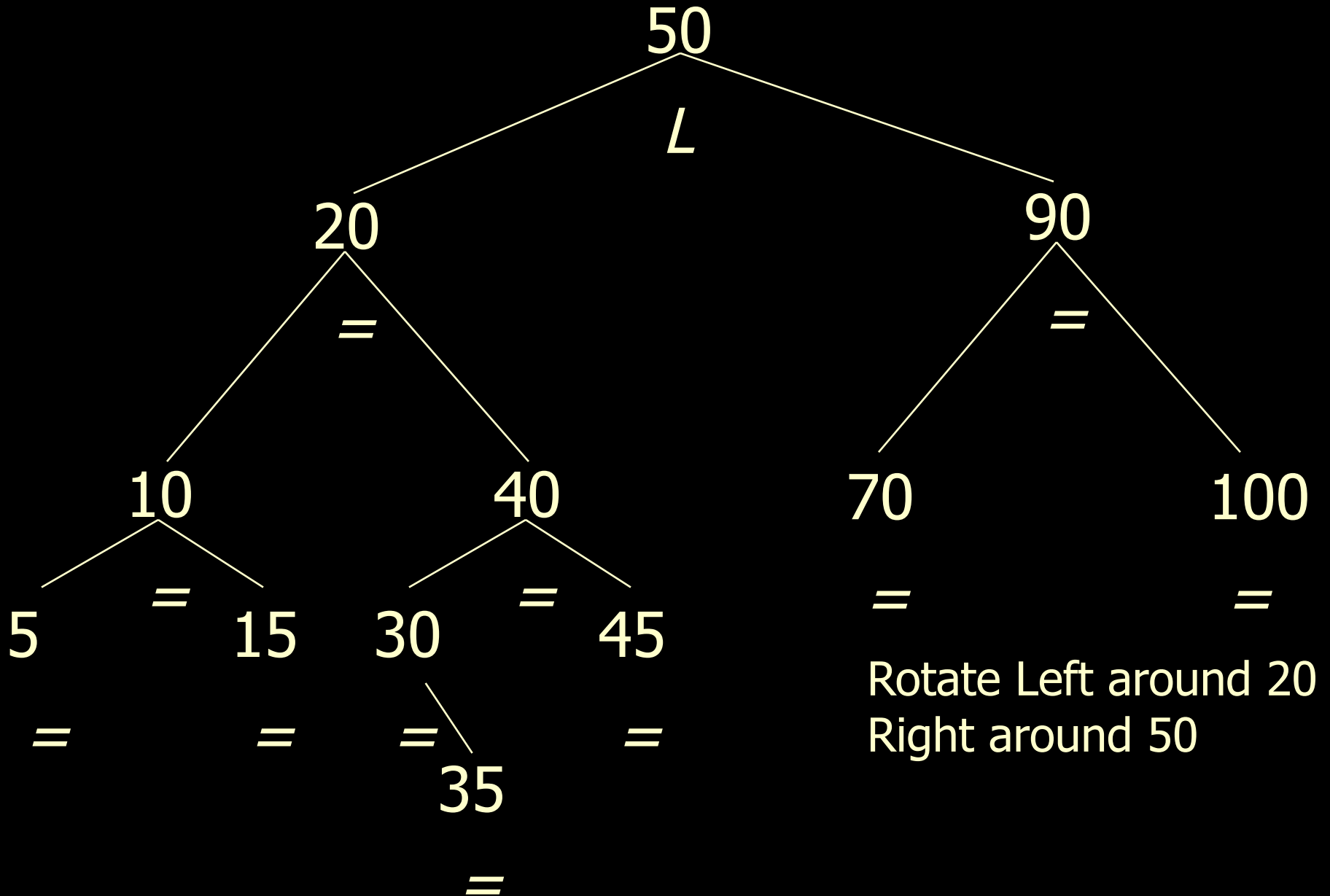


```
protected void adjustLeftRight(Entry ancestor, Entry inserted) {
    Object o = inserted.element;
    if (ancestor.parent == inserted) ancestor.balanceFactor = '=';
    else if (((Comparable)o).compareTo(ancestor.parent.element)<0) {
        ancestor.balanceFactor = 'R';
        adjustPath(ancestor.parent.left,inserted);
    }// o < ancestor's parent
    else {
        ancestor.balanceFactor = '=';
        ancestor.parent.left.balanceFactor = 'L';
        adjustPath(ancestor,inserted);
    }// while
} //adjustLeftRight
```

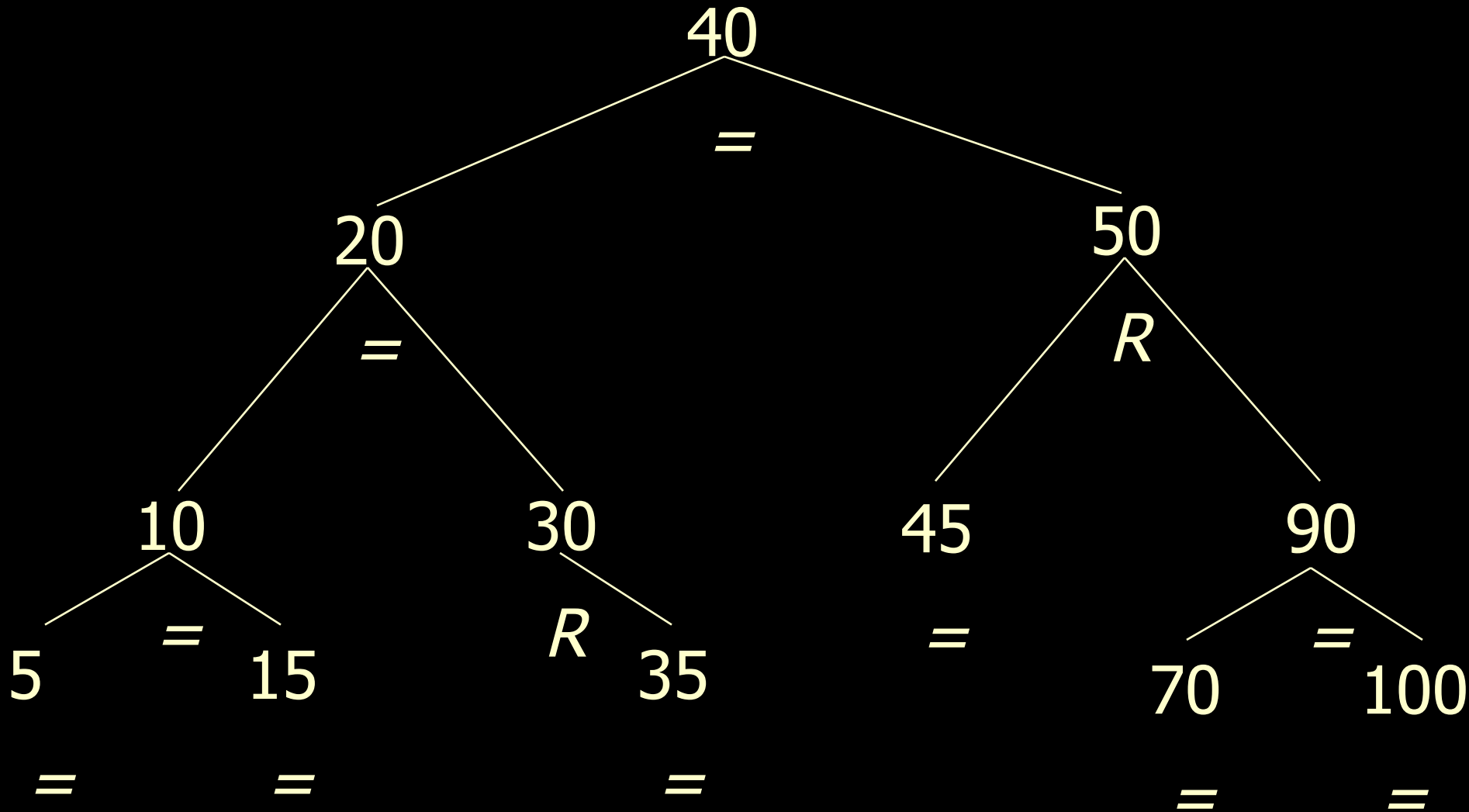
# AdjustLeftRight Case 1



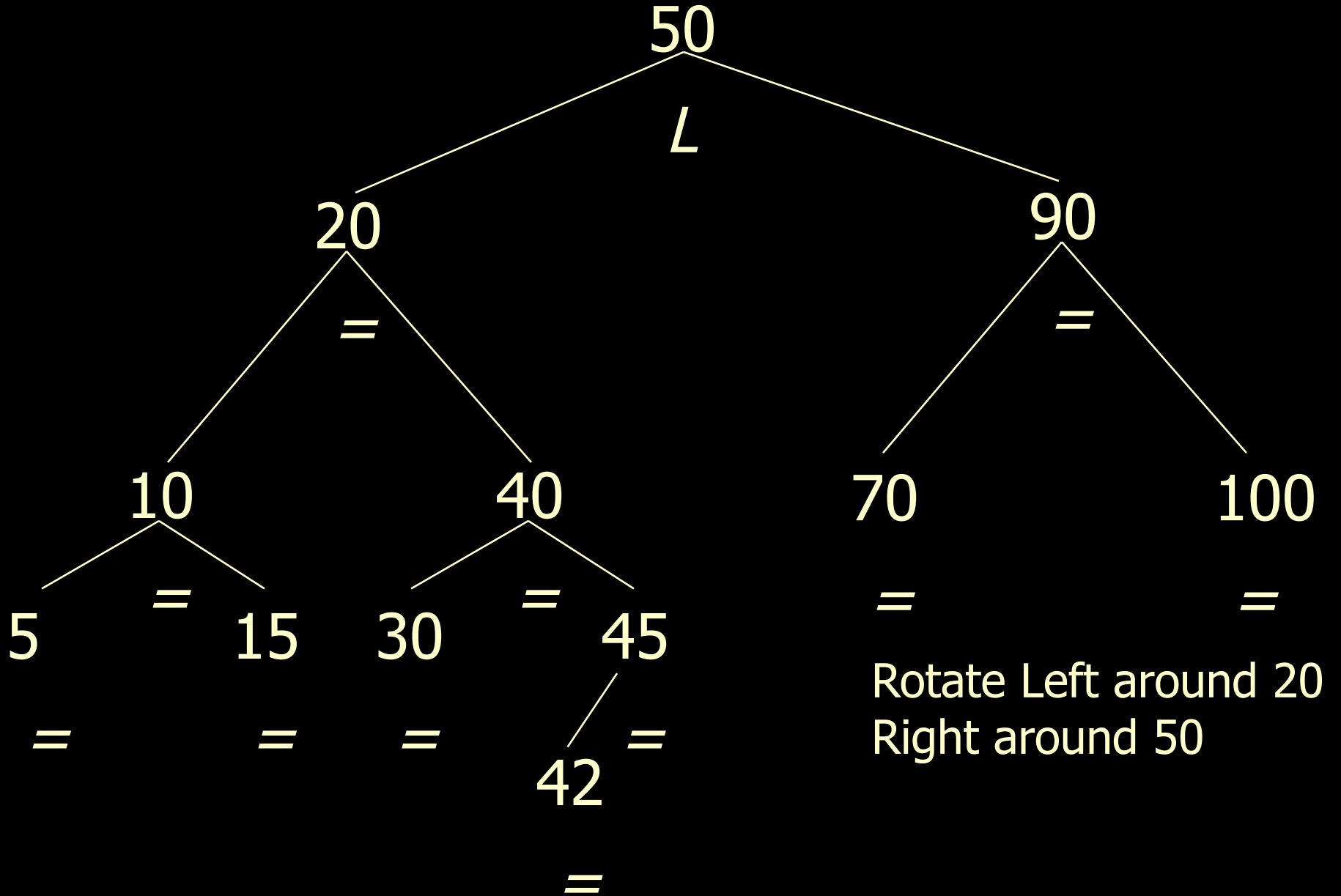
# AdjustLeftRight Case2



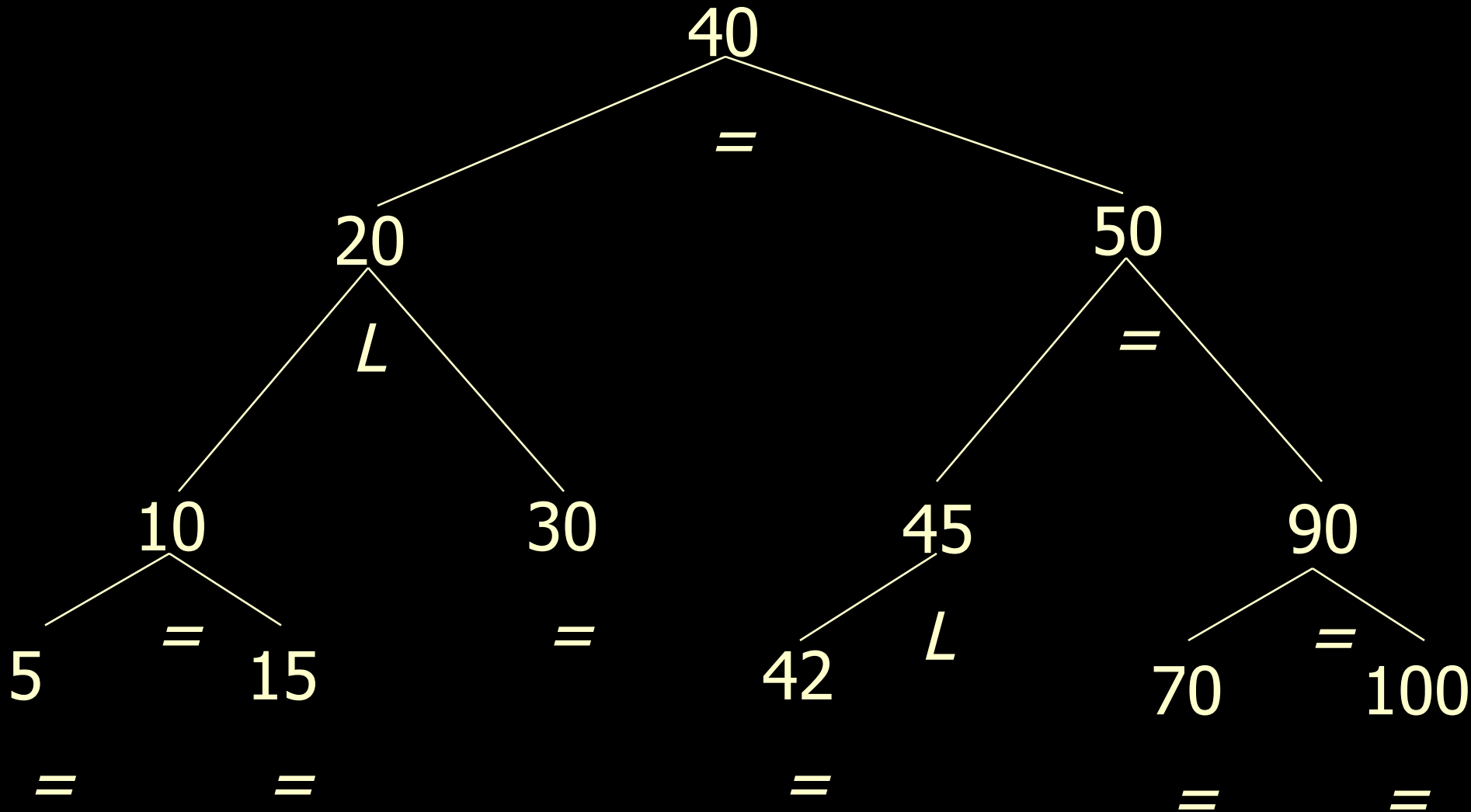
# AdjustLeftRight Case2



# AdjustLeftRight Case3



# AdjustLeftRight Case3



```
protected void adjustRightLeft(Entry ancestor, Entry inserted) {
    Object o = inserted.element;
    if (ancestor.parent == inserted) ancestor.balanceFactor = '=';
    else if (((Comparable)o).compareTo(ancestor.parent.element)>0) {
        ancestor.balanceFactor = 'L';
        adjustPath(ancestor.parent.right,inserted);
    }// o < ancestor's parent
    else {
        ancestor.balanceFactor = '=';
        ancestor.parent.right.balanceFactor = 'R';
        adjustPath(ancestor,inserted);
    }// while
} //adjustRightLeft
```