



Using LL/SC to Simplify Word-based Software Transactional Memory

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LL/SC as a restricted KCSS (*K-Compare-Single-Swap*)

Here $K = 2$. Already-read value of B is 4.

```
old = LL(&A)
if (B == 4) // compare
    if (SC(&A,new) == false) // compare & swap
        return false
    else
        return true
else
    return false
```

NB: requires that A and B always play the same roles



Software Transactional Memory

- A nonblocking synchronization construct
- Makes concurrent programming (almost) as simple as sequential programming
- Gives disjoint access parallel concurrency

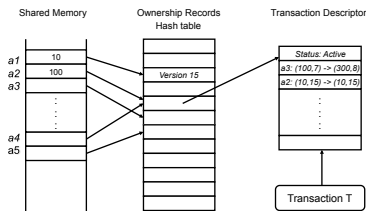


Word-based Software Transactional Memory (WSTM)

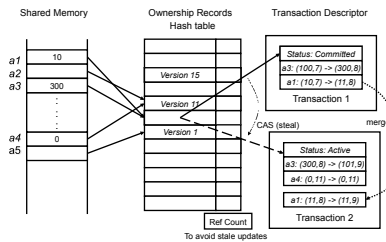
- by Harris and Fraser [OOPSLA'03]
- API for accessing shared memory
 - STMStart
 - STMRead(addr), STMWrite(addr,value)
 - STMCommit(), STMAbort()
 - STMValidate()



WSTM Data Structures



WSTM Transactions



WSTM Drawbacks

- Bounded memory blow-up problem – due to merging and false conflicts
- Complexity – merging, stealing, redo, reference counts
- False merging – merge even when not necessary



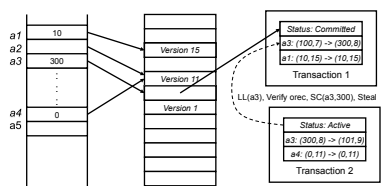
2CSS Application in Atomic Copyback

```
old = orec
...
LL(addr)
if (orec == old) // compare
    if (SC(addr,newval)) // compare and swap
        return true
    else
        return false
```

NB: meets the requirement that orec and addr always play the same roles



Proposed Simplification



Tradeoffs of Proposed Simplification

- + Simplicity – No merging, no stealing, no redo, just selective helping
- + No reference counts
- + No false merging
- Too many LL/SCs (equal to number of locations updated)
- Future Work – Empirical evaluation of proposed simplification

