CSC2/458 Parallel and Distributed Systems Distribute Computing – Other Programming Models

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URCS

Spark's Abstractions

The Spark Runtime

Spark's Abstractions

The Spark Runtime

Abstractions for Shared Memory Programming

- Shared address space
 - One Address, One Value
- Shared memory
 - Disk
 - RAM
- Coherent Caches
 - Disk Cache (OS)
 - Processor Cache (Processor)
- Locking
 - Libraries
 - Hardware Atomics
- Resilience
 - ECC (Hardware)

• Distributed Name Space

• e.g. ?

- Distributed Shared Memory
- Distributed File Systems
- Distributed Caching
 - e.g. memcached
- Distributed Concurrency Control
 - i.e. locking and consistency
- Data Distribution and Marshalling
 - e.g. ntoh, hton
- Distributed Execution
- Resilience
 - e.g. ?

Provided by most OS



Provided by MPI

- Distributed Name Space (Ranks)
- Send/Recv
- Communication Primitives
- Distributed Execution (SPMD)
- No:
 - distributed shared memory
 - distributed file system
 - caching
 - locking and consistency
 - resilience
 - marshalling

What is Erlang?

Maybe, now, we should ask what is Elixir?

Verbatim from the Erlang FAQ: Introduction

- Erlang provides a simple and powerful model for error containment and fault tolerance (supervised processes).
- Concurrency and message passing are a fundamental to the language. Applications written in Erlang are often composed of hundreds or thousands of lightweight processes. Context switching between Erlang processes is typically one or two orders of magnitude cheaper than switching between threads in a C program.

Verbatim from the Erlang FAQ: Introduction

- Writing applications which are made of parts which execute on different machines (i.e. distributed applications) is easy. Erlang's distribution mechanisms are transparent: programs need not be aware that they are distributed.
- The Erlang runtime environment (a virtual machine, much like the Java virtual machine) means that code compiled on one architecture runs anywhere. The runtime system also allows code in a running system to be updated without interrupting the program.

What is Hadoop?

What does Hadoop give us?

- Distributed File System
 - HDFS
- MapReduce programming model
 - distributed execution
 - marshalling
 - caching
- Resilience
 - Always writes to stable storage
 - Reruns failed jobs
- No need for:
 - distributed name space
 - distributed shared memory
 - concurrency control

MapReduce

?

Erlang/Elixir

MPI

Sockets

Spark's Abstractions

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A Spark of an idea

Resilient Distributed Datasets: A Fault-Tolerant Abstraction for In-Memory Cluster Computing

Matei Zaharia, Mosharaf Chowdhury, Tathagata Das, Ankur Dave, Justin Ma, Murphy McCauley, Michael J. Franklin, Scott Shenker, Ion Stoica University of California, Berkeley

- Apache Spark
 - "fast and general engine for large-scale data processing."
- Translation: Write more than MapReduce programs easily
 - Compared to MPI
- 100x faster than Hadoop
 - "In-memory"
- You can write your own data processing engine on top of Spark

What Spark Provides

- Distributed File system
 - Reuses HDFS
- Distributed Execution
- Data Partitioning
- Marshalling
- Resilience
- Distributed Caching
 - No coherence required (or supported)
- No:
 - Distributed Concurrency Control (not supported)
 - Distributed Shared Memory (fine-grained)

Spark Programming Model: 10000ft overview

- Spark is a limited programming model
- Built on observation that:
 - "Many parallel applications naturally apply the same operations to multiple data items"
 - i.e. data-parallel model, e.g. SIMD, SPMD, etc.
- Provides a distributed data structure
 - Resilient Distributed Datasets (RDDs)
 - Like a huge table, but could be anything really
- Programs (you write) operate on RDDs in a *coarse-grained* fashion
 - They always operate on entire RDDs
 - I.e. on *all* elements in a RDD
 - Constrast with DSM which allows fine-grained accesses
- Java/Python/Scala

Resilient Distributed Datasets (RDD)

- A RDD is a
 - "read-only, partitioned set of records"
- Can only be built by "deterministic operations" (*transformations*) on:
 - data in stable storage
 - or other RDDs
- RDDs "remember" the operations that were used to create them
 - paper calls this "lineage"
- RDDs exist "lazily" in memory
 - the operations are only applied when needed (database lingo: materialized)
 - can be stored on disk too
- Why are these properties important?

Aspect	RDDs	Distr. Shared Mem.		
Reads	Coarse- or fine-grained	Fine-grained		
Writes	Coarse-grained	Fine-grained		
Consistency	Trivial (immutable)	Up to app / runtime		
Fault recovery	Fine-grained and low- overhead using lineage	Requires checkpoints and program rollback		
Straggler mitigation	Possible using backup tasks	Difficult		
Work placement	Automatic based on data locality	Up to app (runtimes aim for transparency)		
Behavior if not enough RAM	Similar to existing data flow systems	Poor performance (swapping?)		

Table 1: Comparison of RDDs with distributed shared memory.

```
lines = spark.textFile("hdfs://...")
errors = lines.filter(_.startsWith("ERROR"))
errors.persist()
```

```
errors.count()
```

$map(f:T \Rightarrow U)$:	$RDD[T] \Rightarrow RDD[U]$
$filter(f: T \Rightarrow Bool)$:	$RDD[T] \Rightarrow RDD[T]$
$flatMap(f: T \Rightarrow Seq[U])$:	$RDD[T] \Rightarrow RDD[U]$
sample(fraction : Float)	:	$RDD[T] \Rightarrow RDD[T]$ (Deterministic sampling)
groupByKey()	:	$RDD[(K, V)] \Rightarrow RDD[(K, Seq[V])]$
$reduceByKey(f:(V,V) \Rightarrow V)$:	$RDD[(K, V)] \Rightarrow RDD[(K, V)]$
union()	:	$(RDD[T], RDD[T]) \Rightarrow RDD[T]$
join()	:	$(RDD[(K, V)], RDD[(K, W)]) \Rightarrow RDD[(K, (V, W))]$
cogroup()	:	$(RDD[(K, V)], RDD[(K, W)]) \Rightarrow RDD[(K, (Seq[V], Seq[W]))]$
crossProduct()	:	$(RDD[T], RDD[U]) \Rightarrow RDD[(T, U)]$
$mapValues(f : V \Rightarrow W)$:	$RDD[(K, V)] \Rightarrow RDD[(K, W)]$ (Preserves partitioning)
sort(c: Comparator[K])	:	$RDD[(K, V)] \Rightarrow RDD[(K, V)]$
partitionBy(p: Partitioner[K])	:	$RDD[(K, V)] \Rightarrow RDD[(K, V)]$

count()	:	$RDD[T] \Rightarrow Long$
collect()	:	$RDD[T] \Rightarrow Seq[T]$
$reduce(f:(T,T) \Rightarrow T)$:	$RDD[T] \Rightarrow T$
lookup(k: K)	:	$RDD[(K, V)] \Rightarrow Seq[V]$ (On hash/range partitioned RDDs)
save(path : String)	:	Outputs RDD to a storage system, e.g., HDFS

Spark's Abstractions

The Spark Runtime

- Spark Programs are DAGS/dependence graphs
 - directed acyclic graphs
 - nodes are RDDs
 - edges are operations
- Scheduler executes in order of dependencies
 - prioritizes edges whose inputs already in memory



```
• Narrow operations pipelined
```

```
• i.e. loop coalescing
```

- Operations scheduled on machines based on locality
 - similar to "owner-computes"

- Each RDD knows how to recreate itself
 - Ultimately from stable storage
- Recreation may use RDDs from other machines
 - "Wide" operations
 - e.g. join
- Or from the same machine
 - "Narrow" operations
 - e.g. map
- Can run in parallel
 - RDDs are immutable

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- RDD.map()
- RDD.reduceByKey()

- RDD.select()
- RDD.groupby()
- etc.

- Google Pregel is a graph query engine
 - operates on graphs: vertices and edges
- Each operation is applied to a vertex in parallel
- Each vertex can send messages to other vertices
- Example Pregel in Spark:
 - RDD.flatMap()
 - RDD.join()

- You need to implement transformations
- And actions
- Spark will take care of the rest...

- Spark provides a somewhat general distributed computing programming model
- Operations on immutable, partitioned datasets
- Partitioning, scheduling, marshalling, resilience, etc. for free
- Immensely popular programming model