CSC2/452 Computer Organization Networking

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URCS

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Outline

Administrivia

Networking

The Internet Protocol (IP)

The Sockets Client Programming Interface

The Sockets Server Programming Interface

An Echo Server and Client
Outline

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An Echo Server and Client
Assignment #4 is out
  - Due date: Tuesday, Nov 26, 7PM
Assignment #5 (last!) will be out Dec 2
  - Due on Dec 11
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Getting different computers to work together requires networking

Many different ways of physically building a network
  - Physical: Copper, Optical fibre, Wireless, etc.
  - Technologies: Ethernet, InfiniBand, GSM, CDMA, 802.11 etc.

Lots of physical (and logical) components
  - Nodes, bridges, routers, etc.

Ultimately, though, software needs to use the network
  - The focus of today’s class
  - See the textbook for details on network organization
  - Or take CSC2/457
Streams and Packets

Stream-oriented communication
- Data is sent in some order and must be received in that same order
- Usually because data bytes are ordered (e.g. contents of a file)
- Communication must be reliable, data bytes must not be lost and errors must be detected
- Usually long-lived

Packet-oriented communication
- Data is sent as “packets”, independent pieces of data
- Packets are not ordered, can be delivered out-of-order
- Packets may be lost, corrupted or duplicated
- Communication usually does not last longer than a single packet
Circuit-switching vs Packet switching

Switching: determining the path taken by data

Circuit-switching
- Establish a circuit from source to destination once at beginning of connection
- Data takes the path established by the circuit
- Beloved of telephone companies
- e.g. ISDN, ATM (not the machines dispensing money)

Packet-switching
- No fixed path from source to destination
- Packets choose their own path, on a hop-by-hop basis
- Foundation of the Internet
- e.g. Internet Protocol
Stream-oriented communication using packets?

ARPA chose packet-switching for the initial design of the Internet because it was highly resilient.

Imagine the network as a graph, with computers as nodes and communication links as edges.

Packet-switching can get the data across even when many links and nodes fail.

Usually by choosing a path around these failures.

But, packets can:

- be reordered
- can be corrupted
- can get lost

How to build an abstraction of stream-oriented communication over a packet-switched network?
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Internet Protocol

- The building block of the Internet
  - two major versions: IPv4 and IPv6
- IP is designed to connect networks (hence internet(working))
  - Typically, machines on a single network can talk directly to each other
  - IP specifies how to route packets across networks
- All nodes in the network have an address
  - IPv4: 32-bit addresses
  - IPv6: 128-bit addresses
- IP usually runs over a local network protocol
  - Usually, Ethernet for wired networks
Ethernet: A local network protocol

- Designed at Xerox PARC by Bob Metcalfe
- Originally designed for a shared broadcast medium
  - All computers share the same communication medium
  - Two (or more) computers cannot talk at the same time
  - Key question: how to allow computers to share this medium?
How Ethernet solves the shared medium problem

- Computers check if a communication is in progress
- If not, they start communicating
- If two computers started communicating at the same time, they’ll interfere with each other
  - Called a *collision*
  - These collisions can be detected
- Both computers will stop and try again after waiting a *random* interval of time
  - This interval of time will increase exponentially on each unsuccessful attempt
  - Called *random (or exponential) backoff*
- Alternatives to Ethernet: Token Ring
How IP works

$ route -n
Kernel IP routing table
Destination Gateway Genmask Flags Metric Ref Use Iface
0.0.0.0 192.168.1.1 0.0.0.0 UG 600 0 0 wlan0
192.168.1.0 0.0.0.0 255.255.255.0 U 600 0 0 wlan0

▶ Each node contains a routing table (you can use the command `route`) to print it out
▶ The routing table contains information on how to route packets
▶ In the example above
  ▶ addresses 192.168.1.0 to 192.168.1.255 are local and packets should be sent directly via interface wlan0
  ▶ all other addresses should be sent to the gateway at 192.168.1.1 which will forward them
▶ These are IPv4 addresses, consisting of four 8-bit numbers separated by periods.
Local IP transmission

- If 192.168.1.5 wants to send packets to 192.168.1.1, it can send them directly.
- However, this requires knowing the MAC address of 192.168.1.1.
  - Media Access Control, i.e. the “Ethernet” address.
- Usually, this is done by broadcasting a query to all local nodes: “who is 192.168.1.1?”
  - This is called the address resolution protocol (or ARP).
  - The result is cached to avoid frequent broadcasts.
- Once a MAC address is obtained, Ethernet can be used to directly send the packet.
Non-local IP transmission

- If 192.168.1.5 wants to send packets to 8.8.8.8, it cannot send them directly.
- For non-local addresses, packets are sent to the gateway.
  - The gateway is on the local network.
- The gateway is connected to at least two networks.
  - If the packet is intended for one of its two networks, it can do local delivery.
  - Otherwise, it forwards the packet to the next gateway.
- The forwarding step can be made faster using “routers.”
  - Gateways that discover and remember the fastest path to a network.
IP Data transmission algorithm

- If address is local
  - find MAC address using ARP (if not known)
  - send packet directly using link-layer protocol (i.e. Ethernet)
- If address is non-local
  - send packet to gateway
IP Packets

▶ TTL: field is decremented by one by each gateway/router
  ▶ Packet discarded when the field reaches 0
▶ Protocol: identifies next-level protocol
  ▶ Ethernet at bottom, followed by IP, followed by ...

(from historical standard RFC 791)
Example IP Packet (or ‘Datagram’)

Protocol 1 is Internet Control Message Protocol (ICMP)

- The protocol used for sending pings
- Or network error messages (“Destination unreachable”)
- BTW, this is not a valid ICMP packet, just an example
TCP and UDP

- Like ICMP, these are next level protocols built on top of IP
- User Datagram Protocol or UDP
  - Provides user-level (think application-level) packet communication
  - Recall, IP is for node-to-node communication
- Transmission Control Protocol or TCP
  - Provides a reliable, stream-oriented communication on top of IP
  - Also user/application-level
UDP extends IP addresses with the notion of a *port*

- A port is a 16-bit unsigned number (0 to 65535)
- An application “listens” on a well-known port
  - e.g. 8.8.8.8:53, where 53 is the port on node 8.8.8.8
  - data addressed to a port is forwarded to the application

UDP packets are encapsulated in IP

- I.e. a UDP packet is the “data” portion of an IP packet
- Protocol is set to 17
TCP

- TCP must:
  - provided stream-like data transfer
  - deliver data in order
  - detect data corruption
  - handle lost packets

- Usually implemented over IP, so TCP/IP
General outline of TCP

- TCP provides stream-like data transfer
  - By using the notion of a *connection* before transferring data
  - Each connection is identified by an address:port pair, one for both sides of a connection
  - e.g. 192.168.1.2:5810 to 8.8.8.8:80

- Delivers data in order
  - Adds sequence numbers which aid ordering

- Detect data corruption
  - Adds checksums

- Handles lost packets
  - Sequence numbers identify lost or late packets
  - Packets are retransmitted after a timeout
  - Also identifies duplicate packets

- Many other details
  - Setting up connections
  - Congestion control
Well-known ports

- TCP requires a connection to be established
  - Need to know addresses of both nodes
  - Need source port and destination port
- Many destination ports are well-known
  - 80 is for a web server using HTTP
  - 443 is for a web server using HTTPS
  - 21 is for file transfer protocol
  - 22 is for secure shell
  - 25 is for SMTP (i.e. sending email)
- Your OS usually assigns a random port to the source side of the connection
Programmer’s Perspective

- Does your application need reliable, ordered communication?
  - Use TCP
- Is it transactional and can handle missing/corrupted packets?
  - Use UDP
- Note sometimes protocols mandate use of TCP or UDP
  - Domain name service (DNS) uses UDP (but also supports TCP)
  - HTTP/1 and HTTP/2 only use TCP
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Sockets

- Original terminology from ARPANET
  - Usage today very different from that definition
- A socket is one end of a connection (i.e. one address:port pair)
- Programming interface for networking applications
  - Developed and popularized by BSD Unix
- Nearly the same interface in any operating system
  - And across programming languages
  - the Plan9 OS is a notable exception
Creating sockets

```c
#include <sys/types.h>
#include <sys/socket.h>

int socket(int domain, int type, int protocol);
```

- **domain** is AF_INET for IPv4
  - also used AF_UNIX for Inter-process Communication
- **type** is SOCK_STREAM or SOCK_DGRAM
  - Roughly correspond to TCP and UDP respectively
- **protocol** is usually 0
- **socket** returns a file descriptor
  - Use `read` and `write` for TCP
  - Use `send` and `recv` for UDP (can also be used for TCP)
Specifying destination

```c
int connect(int sockfd, const struct sockaddr *addr, socklen_t addrlen);
```

- For TCP, `connect` will try to open a connection to `addr`.
- For UDP, `connect` will send any future packets to `addr`.
  - It will also only accept packets sent by `addr`.
- These must be numeric addresses, not names.
  - I.e. 172.217.7.14 not google.com
Aside: Resolving names to addresses

```c
#include <sys/types.h>
#include <sys/socket.h>
#include <netdb.h>

int getaddrinfo(const char *node, const char *service,
    const struct addrinfo *hints,
    struct addrinfo **res);
```

- Names are “converted” into addresses using a resolver
  - Known as name resolution
  - Needs to look up a database of names to addresses
- Nearly always use the Domain Name Server (DNS) protocol
  - But can also use other alternatives like UPnP
- `getaddrinfo` takes a name (node) and a service (“http”) and returns an address
  - Service can be NULL
  - Don’t use `gethostbyname` which is obsolete
Sending data

ssize_t send(int sockfd, const void *buf, size_t len, int flags);

- Sends data to a destination
  - Must have been specified using a previous connect call
Receiving data

ssize_t recv(int sockfd, void *buf, size_t len, int flags);

- Receives data from a socket
  - This must be a connected socket
- `recv` will block until data is received
Blocking vs Nonblocking I/O

- Many functions that perform I/O block by default
  - read, recv, scanf
- Although this simplifies programming, programs may want to check if data is available or not
- Programs can use fcntl to set a socket to non-blocking mode
- In this case, if the operation would normally block, it will instead return EWOULDBLOCK
  - Or EAGAIN
- Allows you to overlap I/O with computation
  - Very useful prior to threading
Asynchronous I/O

- Instead of checking if an I/O operation has completed, you can use asynchronous I/O
- OS will notify your program that I/O is complete
- All operations in asynchronous I/O are non-blocking by default
- See manual page for aio for more information
  - Or see Overlapped I/O in Windows
Overview of Client-side Interface

- Create a socket
- connect it
- send or recv data from it
- close it
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Server-side sockets: bind

```c
int bind(int sockfd, const struct sockaddr *addr,
         socklen_t addrlen);
```

- `sockfd` is obtained from `socket`
- `addr` is *local* address to accept requests on
  - Usually set to 0.0.0.0 for IPv4 to accept requests on all of the node’s IP addresses
int listen(int sockfd, int backlog);

- sockfd is a socket for which bind has been called
- backlog specifies queue depth – connections made after this has been exceeded may be refused
int accept(int sockfd, struct sockaddr *addr, socklen_t *addrlen);

- Waits until a connection is made (unless the socket is non-blocking) to sockfd
  - which must be a socket on which listen has been called
- Returns the address of the other side of the connection in addr
- Returns the file descriptor for an accepted socket, i.e. the connection
  - This is used to send and receive data
  - The original socket continues to listen for more connections
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l = socket(AF_INET, SOCK_STREAM, 0);
if(l == -1) {
    perror("socket");
    exit(1);
}

struct sockaddr_in addr;
memset(&addr, 0, sizeof(addr));
addr.sin_family = AF_INET;
addr.sin_port = htons(5000); /* convert to network byte order */
addr.sin_addr.s_addr = INADDR_ANY;

if(bind(l, (struct sockaddr *) &addr, sizeof(addr)) == -1) {
    perror("bind");
    exit(1);
}
Echo Server: Listening

```c
if(listen(l, 1) == -1) {
    perror("listen");
    exit(1);
}
```
struct sockaddr_in sender;
socklen_t addrlen = sizeof(sender);
int s, r;

char buf[256];

while((s = accept(l, (struct sockaddr *) &sender, &addrlen)) != -1) {
    r = recv(s, buf, 255, 0);
    if(r == -1) {
        perror("recv");
        close(s);
    } else {
        buf[r] = '\0';

        printf("received '%s'\n", buf);

        if(send(s, buf, r, 0) == -1) {
            perror("send");
        }
        close(s);
    }
}
l = socket(AF_INET, SOCK_STREAM, 0);
if(l == -1) {
    perror("socket");
    exit(1);
}
struct sockaddr_in addr;

memset(&addr, 0, sizeof(addr));
addr.sin_family = AF_INET;
addr.sin_port = htons(5000);
addr.sin_addr.s_addr = htonl(0x7f000001); // 127.0.0.1

if(connect(l, (struct sockaddr *) &addr, sizeof(addr)) == -1) {
    perror("connect");
    exit(1);
}

printf("connected\n");
char buf[256];
int r;

r = scanf("%255s", buf);
if(r != 1) {
    perror("scanf");
    exit(1);
}

if(send(l, buf, strlen(buf)+1, 0) == -1) {
    perror("send");
    exit(1);
}

r = recv(l, buf, 255, 0);
if(r == -1) {
    perror("recv");
    exit(1);
} else {
    buf[r] = '\0';
    printf("got '%s'\n", buf);
}
Problems with the Echo Server

- The echo server handles one connection at a time
- It is serial
- What if a client connects but does not send a message?
  - What happens to clients who connected after that client?
References

- Chapter 11 of the textbook
  - And Chapter 12.2 which talks about I/O multiplexing (one solution to one client at a time)
- Read the manual pages for `socket` and pages related to it
  - Or read the GNU libc manual for sockets