CSC2/455 Software Analysis and Improvement Intermediate Representations (IRs)

Sreepathi Pai February 3, 2021

URCS

Outline

Introduction

 ${\sf Midend}$

Miscellaneous

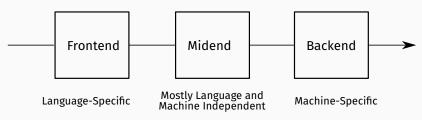
Outline

Introduction

Midena

Miscellaneous

Classic Compiler Architecture



Recommended reading: Chris Lattner, LLVM, The Architecture of Open Source Applications

Outline

Introduction

Midend

Miscellaneous

What does the midend do?

- Mostly language and machine independent analyses
- Majority of analyses run in this stage
- Multiple intermediate representations used
 - Starts from abstract syntax tree
 - Usually stops before instruction scheduling/register allocation
 - Examples: AST, CFG, DDG, PDG, etc.
- There is no one Intermediate Representation (IR)
 - although people have tried ...

Mid-end structure

- Organized as a set of *passes*
- Each pass usually performs one task
 - Some specific analysis of the IR
 - Some transformation of the IR
- Input to each pass is the IR and output is also the IR
 - And usually analysis results, etc.

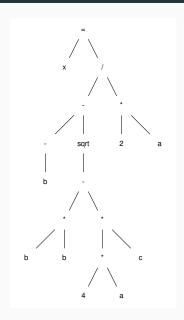
Running example

$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

In Python (assume math.sqrt is sqrt):

$$x = (-b - sqrt(b*b - 4 * a * c)) / (2*a)$$

Abstract Syntax Tree



AST as a list

```
ast = ['='],
       ['x',
          ['-',
           ['neg', 'b'],
           ['sqrt',
             ['*', 'b', 'b'],
             ['*',
['*', '4', 'a'],
```

• What does this remind you of?

Slightly less LISPy

```
class Node(object):
    operator = None
    left = None
    right = None
```

- Actual ASTs are not binary trees!
 - Will usually have list of descendants instead of left and right
 - descendants may be more specific, while may have condition and body

Linear Forms

- ASTs imply treewalking
 - Works best when manipulating source code
 - e.g. Source-to-source compilers
 - Or when control flow is not important
- Other lower level forms are "closer to machine"
 - Stack machines
 - 3 address code

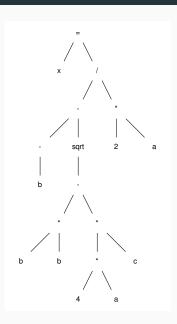
Stack Machine

```
push a
push 2
mul
push c
push a
push 4
mııl
mul
push b
push b
mııl
sub
sqrt
push b
neg
sub
div
pop x
```

How do you produce stack machine code from an AST?

Generating Stack Machine Code

push a push 2 mul push c push a push 4 mul mul push b push b mul sub sqrt push b neg sub div рор х



Stack machines

- Compact in size
 - Operands are implicit on top of stack.
- Easy to execute
- BUT, fixed order of execution
 - Bad for parallelism
- Hard to analyze
- Nevertheless, widely used:
 - Java bytecode
 - Python bytecode
 - WebAssembly

3 address code

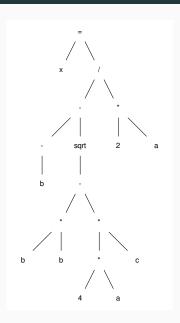
- 3 "addresses"
 - Two source operands
 - One destination operand
 - One operation
- Addresses are actually names generated by compiler
 - Or refer directly to variables

Our example in 3-address code

```
t1 <- -b
    t2 < -b * b
    t3 < -4 * a
    t4 < -t3 * c
    t5 <- t2 - t4
    t6 <- sqrt(t5)
    t7 <- t1 - t6
    t8 <- 2 * a
    t9 <- t7 / t8
    x < - t.9
(Here '<-' is \leftarrow, signifying assignment)
```

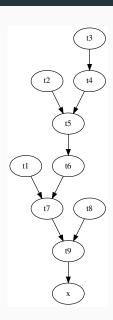
How do we produce 3-address code from the AST?

Producing 3-address code



Data Dependence Graphs (DDGs)

- DDGs track "data flow" as an acyclic graph
- Strict (partial) order in which operations must be performed
 - Can't use a value that has not been calculated yet!
- But multiple orders may be allowed!
 - Topological sort
- Will revisit DDGs when we discuss instruction scheduling



DDG Example

```
t1 <- -b

t2 <- b * b

t3 <- 4 * a

t4 <- t3 * c

t5 <- t2 - t4

t6 <- sqrt(t5)

t7 <- t1 - t6

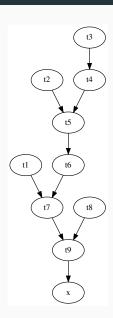
t8 <- 2 * a

t9 <- t7 / t8

x <- t9
```

DDG Example, reordered by dependence

```
# group 1
t1 <- -b
t2 < -b * b
t3 < -4 * a
t8 <- 2 * a
# must wait for t3
t4 < -t3 * c
# must wait for t4 and t2
t5 <- t2 - t4
# ...
t6 <- sqrt(t5)
t7 <- t1 - t6
t9 <- t7 / t8
x < - t9
```

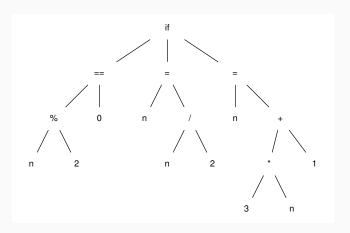


Control structures in 3-address code

What should the 3 address code for the code below look like?

```
if n % 2 == 0:
    n = n / 2
else:
    n = 3 * n + 1
```

The AST for if



 An if AST node has a condition, true-code, and false-code

3-Address Code for if

```
t1 <- n % 2

tc <- t1 == 0

if (tc == 0) goto L1

t2 <- n / 2

n <- t2

goto L2

L1:

t3 <- 3 * n

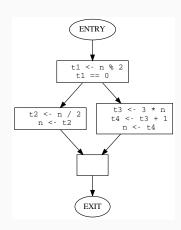
t4 <- t3 + 1

n <- t4
```

- 3-address code can contain:
 - conditional branches, usually just a comparison to zero
 - · unconditional branches
 - labels

Control Flow Graphs (CFGs)

- "Hybrid" representation
 - Linear code + Graph structure
- Each node in the CFG is a "basic block"
 - Linear code
 - Single entry, single exit
 - "Straight-line code"
- Most common form for analysis



In this course

- We will write midend + some bits of a backend
- Input language: C
- Output language: C in 3-address form
 - Not assembly (maybe extra credit?)
- Using Python library pycparser

For assignments, make sure to review

- Basic data structures
 - lists
 - trees
 - graphs
- Basic data structure traversals
 - Infix, prefix, postfix
 - Depth-first, breadth-first
- And how to implement them in Python
 - Using Python standard libraries is fine

Outline

Introduction

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Miscellaneous

Not-so-classic 'Compiler' Architectures

What is this code from TensorFlow doing?

```
a = tf.constant(2)
b = tf.constant(3)
with tf.Session() as sess:
    print("a=2, b=3")
    print("Addition with constants: %i" % sess.run(a+b))
    print("Multiplication with constants: %i" % sess.run(a*b))
```

https://github.com/aymericdamien/TensorFlow-Examples/blob/master/examples/1_Introduction/basic_operations.py

Metaprogramming

- TensorFlow API builds a graph
 - directed, acyclic
 - similar to the DDG
 - very common technique
- When sess.run is called, graph is compiled and executed
- Advantages:
 - No syntax, no parsing!
- Disadvantages:
 - 1

References

- Chapter 5 of Cooper and Turczon
 - Up to 5.4 in this lecture, but we will ultimately study the whole chapter