# CSC2/455 Software Analysis and Improvement Dead Code Elimination

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**URCS** 

#### **Outline**

Review

Dead Code Elimination

Postscript

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#### So far

- Source code
- Three-address form
- Control-flow graphs
- SSA form
- Data flow analyses

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#### **Definitions**

- Dead code
  - Useless operation: Not externally visible
  - Unreachable code: Cannot be executed
- Critical operation: (Direct) "Useful operation"
  - Operation that computes return value
  - Operation that stores to memory (i.e. is externally visible)
  - Operation that performs I/O
  - ...

#### Two Steps: Step 1

- Find all directly useful operations and mark them
- Find all indirectly useful operations and mark them
  - I.e. those that feed into directly useful operations
- Iterate until all operations that ultimately feed into directly useful operations have been found and marked

## Two Steps: Step 2

• Remove all operations that remain unmarked

### Example #1

```
void swap(int *x, int *y) {
  int t;

  t = *x;
  *x = *y;
  *y = t;
}
```

# Example #2

```
int min(int x, int y) {
   int r;

   if (x > y) {
     r = y;
   } else {
     r = x;
   }

   return r;
}
```

### Example #2: 3AC

```
int min(int x, int y) {
  int r;
  int t;
  t = x > y;
  if(t == 0) goto L1;
  r = y;
  goto L2;
L1:
r = x;
L2:
 return r;
```

# Example #2: With useless operations removed

```
int min(int x, int y) {
  int r;

  r = y;
  r = x;

  return r;
}
```

- Marking and removing useless operations uses only dataflow information
- Must also preserve control flow (i.e. control dependences)
  - How to identify useful branches?

## Handling Control Flow

- Assume all "jumps" (unconditional branches) are useful
  - i.e. goto Lx
- What about conditional branches?

#### **Conditional Branches: Example**

```
int first_N_sum(int N) {
   int s = 0;

for(int i = 1; i <= N; i++)
      s = s + i;

return N * (N + 1) / 2;
}</pre>
```

#### 3AC code for conditional branches

```
int first_N_sum(int N) {
   int s = 0;
   int i, t;
   i = 1;
L1:
   t = i \le N;
   if(t == 0) goto L2;
   s = s + i;
   i++;
   goto L1;
I.2:
   return N * (N + 1) / 2;
```

How do we recognize that the conditional branch is useless in this case?

# GCC 8.2 for x86-64 (-O0)

```
first_N_sum(int):
        push
                rbp
        mov
                rbp, rsp
                DWORD PTR [rbp-20], edi
        mov
                DWORD PTR [rbp-4], 0
        mov
                                            ; s = 0
                DWORD PTR [rbp-8], 1
                                             ; i = 1
        mov
 .L3:
                eax, DWORD PTR [rbp-8]
        mov
                eax, DWORD PTR [rbp-20]
        cmp
        jg
                 .L2
        : s = s + i
            eax, DWORD PTR [rbp-8]
        mov
                DWORD PTR [rbp-4], eax
        add
                DWORD PTR [rbp-8], 1
        add
                 .L3
        jmp
 .L2:
                eax, DWORD PTR [rbp-20]
        mov
        add
                eax, 1
                eax, DWORD PTR [rbp-20]
        imul
        mov
                edx, eax
        shr
                edx, 31
        add
                eax, edx
        sar
                eax
        pop
                rbp
        ret.
```

# GCC 8.2 for x86-64 (-O1)

```
first_N_sum(int):
              edi, edi
       test
       jle .L2
       lea edx, [rdi+1]
       mov eax, 1
                             ; i = 1
 .L3:
       add
            eax, 1
                             ; i = i + 1
               eax, edx
       cmp
       jne
              .L3
 .L2:
       lea
              eax, [rdi+1]
       imul
            edi, eax
             eax, edi
       mov
       shr
              eax, 31
       add
               eax, edi
       sar
               eax
       ret
```

# GCC 8.2 for x86-64 (-O2)

```
first_N_sum(int):
    lea    eax, [rdi+1]
    imul    edi, eax
    mov    eax, edi
    shr    eax, 31
    add    eax, edi
    sar    eax
    ret
```

All compiler output examples obtained using the Compiler Explorer.

#### **Conditional Branches**

- A conditional branch is useful only if:
  - A useful operation depends on it
- Control dependence
  - (informal) an operation O is dependent on a branch B if the direction of the branch B affects if O is executed
  - CFG property

#### **Example of control dependence**

```
t = x > y
if(t == 0) goto L1

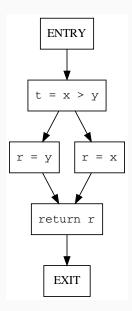
r = y;
goto L2;

L1:
r = x;

L2:
return r;
```

The assignments to r are dependent on if(t == 0), but return r is not

# Control dependence in the CFG

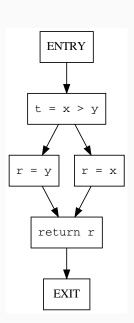


#### **Control Dependence: Formal Definition**

- Postdominance
  - A node n postdominates m if it occurs on all paths from m to EXIT
- A node *k* is control dependent on *i* if:
  - For a path  $i \to j_0 \to j_1 \to ... \to k$ , k postdominates all  $j_x$
  - k does not strictly postdominate i

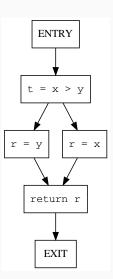
# Control Dependence: Example #1

- Consider k: r = y
- Is it control dependent on i:t = x > y?
- Only one path  $i \to k$ 
  - r = y post-dominates r= y
  - r = y does not strictly postdominate *i*
  - Because it is not a post-dominator of i, and k ≠ i
- So k is control-dependent on i



# Control Dependence: Example #2

- Now, consider k: return r
  - i is still t = x > y
- Two paths, first path
  - $i \rightarrow j_0 \rightarrow k$ 
    - $j_0$  is r = y
    - return r post-dominatesr = y and itself
    - return r strictly postdominates i
    - Because it is a post-dominator of i, and k ≠ i
- So k is not control-dependent on i

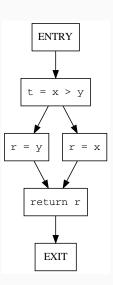


# Path #2 of Example #2

- ullet Second path is  $i o j_1 o k$ 
  - $j_1$  is r = x
- ullet Similar arguments show that k is control-dependent on i

# Using Reverse Dominance Frontiers (RDF)

- Given that return r is useful, so are r = x and r = y
- We can see that t = x > y
   is in the reverse dominance
   frontier (RDF) of r = x
   and r = y
  - RDF is DF on edge-reversed CFG.
- Indeed, RDFs identify control dependences



# Marking unconditional branches useful

- If node k contains useful operations,
- And if *k* is control-dependent on node *i*,
- Then the (conditional) branch in *i* is useful.
- Operationalized as:
  - If block k contains useful operations
  - Mark all conditional branches in k's reverse dominance frontier RDF(k) as useful
  - RDF computed as DF on edge-reversed CFG

# Dead Code Elimination: High-level algorithm

- Mark all directly useful operations
- Repeat until convergence
  - Mark all indirectly useful operations
  - Mark all conditional branches in RDFs of useful operations as useful
- Remove all unmarked operations
- Remove empty nodes in CFG / remove all useless control flow

See algorithms in Figure 10.1 and 10.2 in Turczon and Cooper.

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#### References

- Chapter 10 of Torczon and Cooper
  - Section 10.2