

# **CSC2/455 Software Analysis and Improvement**

## **Array/Loop Dependence Analysis**

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URCS

# Outline

Loop Analysis and Transformation

Characterizing loop dependences

Identifying Loop Dependences

Distance and Direction Vectors

Postscript

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## Loop Analysis and Transformation

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# Why Loop Transformations

- Potentially lots of computation
  - A few operations execute many times
- Potentially lots of memory accesses
- Array-based data structures show up frequently
  - Matrices, vectors, etc.
- Loops are naturally paired with arrays
- FORTRAN
  - FORMula TRANslator
  - World's first high-level programming language

# Important Applications

- Scientific Computing/Computational Science
  - Simulation of Galaxies, Molecules, etc.
  - Drug Discovery
- Audio/Video Processing
  - Signal Processing
  - Compression
- Machine Learning (specifically Deep Learning)
  - Recognizing cats
  - Showing targeted ads

# Matrix Multiply – IJK

- Multiplying two matrices:
  - $A (m \times n)$
  - $B (n \times k)$
  - $C (m \times k)$  [result]
- Here:  $m = n = k$

```
for(ii = 0; ii < m; ii++)  
  for(jj = 0; jj < n; jj++)  
    for(kk = 0; kk < k; kk++)  
      C[ii * k + kk] += A[ii * n + jj] * B[jj * k + kk];
```

## Matrix Multiply – IKJ

```
for(ii = 0; ii < m; ii++)  
  for(kk = 0; kk < k; kk++)  
    for(jj = 0; jj < n; jj++)  
      C[ii * k + kk] += A[ii * n + jj] * B[jj * k + kk];
```

## Performance of the two versions?

- on 1024x1024 matrices of `ints`
- which is faster?
- by how much?



## Performance of the two versions

- on 1024x1024 matrices
- Time for IJK:  $0.554 \text{ s} \pm 0.003\text{s}$  (95% CI)
- Time for IKJ:  $6.618 \text{ s} \pm 0.032\text{s}$  (95% CI)

# What caused the nearly 12X slowdown?

- Matrix Multiply has a large number of arithmetic operations
  - But the number of operations did not change
- Matrix Multiply also refers to a large number of array elements
  - Order in which they access elements changed
  - Memory hierarchy/Caches affect cost of each access

# Loop Analysis

Which *iterations* of a loop depend on other iterations?

- Once dependences are identified
  - can change order
  - can vectorize
  - can redistribute across processors

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# Dependences

- Definition: Two *dynamic* statements have a dependence if:
  - Both access same location (memory or register)
  - And one of the accesses is a write
- *Dynamic* required, since we're talking about loops (examples later)

## Dependence types

$S_1$  occurs “earlier” than  $S_2$  (in the dynamic trace)

- True dependence
  - $S_1 \delta S_2$
  - $S_1$  writes,  $S_2$  reads
- Anti-dependence
  - $S_1 \delta^{-1} S_2$
  - $S_1$  reads,  $S_2$  writes
- Output dependence
  - $S_1 \delta^o S_2$
  - Both  $S_1$  and  $S_2$  write

## Loop-independent dependence

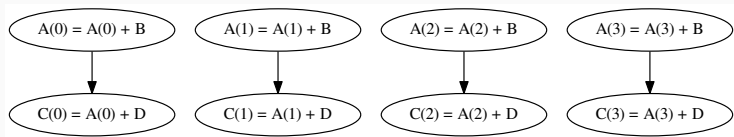
- What are the dependences in the loop body below?
- Can you change the order of the statements in the loop body?

```
DO I = 0, 9
  A(I) = A(I) + B
  C(I) = A(I) + D
ENDDO
```

- Can you change the (execution) order of loop iterations?

Note: FORTRAN uses parentheses in array references: e.g.,  $A(I)$ . FORTRAN arrays usually start at 1, but for this lecture, we will assume they start at 0.

## Loop-independent dependences visualized



NOTE: Only dependences from first four iterations visualized.



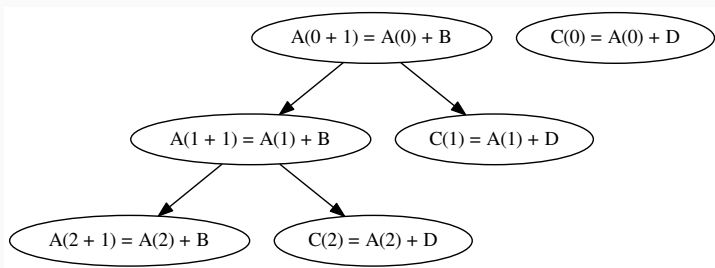
## Loop-carried dependences

- What are the dependences in the loop body below?
- Can you change the order of the statements in the loop body?

```
DO I = 0, 9
  A(I + 1) = A(I) + B
  C(I) = A(I) + D
ENDDO
```

- Can you change the (execution) order of loop iterations?

## Loop-carried dependences visualized



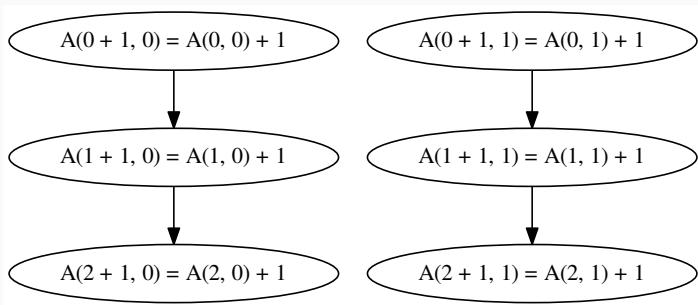
NOTE: Only dependences from first three iterations visualized.

## Dependence Level for Loop-Carried Dependences

```
DO I = 0, 9
  DO J = 0, 1
    A(I + 1, J) = A(I, J) + 1
  ENDDO
ENDDO
```

- Can you change the order of inner loop?
- Can you change the order of the outer loop?

## Dependencies Visualized



NOTE: Only dependencies from first three iterations visualized.

# Loop Dependences

- Loop-independent dependence
  - In same iteration, independent of loops
- Loop-carried dependence
  - Across different iterations of atleast one loop
- Dependence Level of a Loop-carried Dependence
  - The nesting level  $k$  of loop that carries the dependence
  - $S_1 \delta_k S_2$

# Iteration Spaces

```
DO I = 1, 2
  DO J = 1, 2
    S
  ENDDO
ENDDO
```

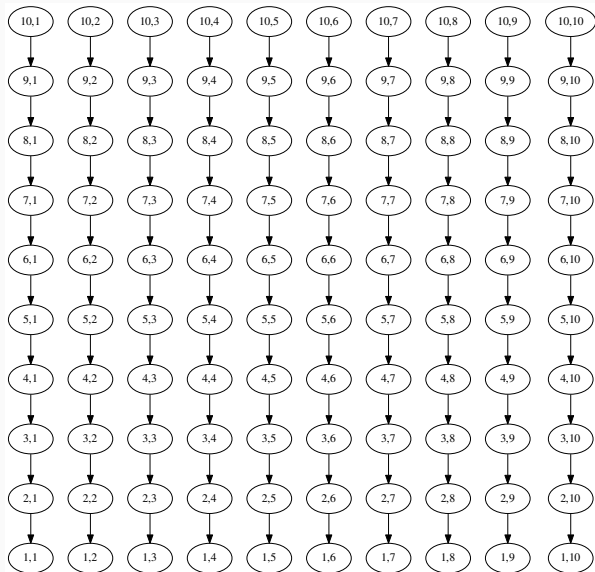
- S has four instances  $(I, J)$ :  $(1, 1)$ ,  $(1, 2)$ ,  $(2, 1)$ ,  $(2, 2)$
- Each of these values represents an *iteration vector*
  - Particular values of loop indices
  - Ordered from outermost loop to innermost loop

# Iteration Space Example

```
DO J = 1, 10
  DO I = 1, 10
    A(I+1, J) = A(I, J) + X
  ENDDO
ENDDO
```

Assuming  $A$  starts from 1. FORTRAN allows you to change the "origin" of arrays.

# Iteration Space Figure





## Iteration Vector Ordering (Definition)

For two vectors  $i = (i_1, i_2, \dots, i_n)$  and  $j = (j_1, j_2, \dots, j_n)$ , each containing  $n$  elements,  $i < j$  if there exists  $m \in [0, n)$ , such that:

- $i_x = j_x$  for  $x < m$
- $i_m < j_m$

## Iteration Vector Ordering (Code)

For two vectors  $i$  and  $j$ , each containing  $n$  elements,  $i < j$  is defined as:

```
def lessthan(i, j, n):
    if n == 1:
        return i[0] < j[0]

    # test prefix for elementwise-equality
    if i[0:n-1] == j[0:n-1]:
        return i[n-1] < j[n-1]
    else:
        return lessthan(i, j, n-1)
```

Can similarly define other order relations.

# Loop dependence

Dependence from Statement S1 (source) to statement S2 (sink) if:

- There exist iteration vectors  $i$  and  $j$  such that  $i < j$  or  $i = j$
- There is a path from S1 to S2 in the loop
- S1 accesses memory location M in iteration  $i$
- S2 accesses memory location M in iteration  $j$
- and one of the accesses is a write

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## Generalizing Loop Indices

```
DO I_1 = ...
  DO I_2 = ...
    ...
    DO I_N = ...
      A(f1, f2, f3, ..., fM) = ...
      ... = A(g1, g2, g3, ..., gM)
    ENDDO
  ENDDO
ENDDO
```

where  $A$  is  $M$ -dimensional array, and  $fX$  and  $gX$  are *index functions* of the form

- $fX(I_1, I_2, \dots, I_N)$
- $gX(I_1, I_2, \dots, I_N)$
- $1 \leq X \leq M$

# Dependence using Iteration Vectors

Let  $\alpha$  and  $\beta$  be iteration vectors:

- $\alpha = (i_1, i_2, i_3, \dots, i_N)$
- $\beta = (i'_1, i'_2, i'_3, \dots, i'_N)$

Then a dependence exists if:

- (vectors)  $\alpha < \beta$
- $fX(\alpha) = gX(\beta)$ , for  $1 \leq X \leq M$

## Example

```
DO J = 0, 9
  DO I = 0, 9
    A(I+1, J) = A(I, J) + X
  ENDDO
ENDDO
```

- $f1(J, I) = I + 1, f2(J, I) = J$
- $g1(J, I) = I, g2(J, I) = J$
- For  $\alpha = (0, 0)$  (i.e.  $J = 0, I = 0$ ) and  $\beta = (0, 1)$  (i.e.  $J = 0, I = 1$ ):
  - $f1(\alpha) = g1(\beta)$ , i.e.  $1 = 1$
  - $f2(\alpha) = g2(\beta)$ , i.e.  $0 = 0$
  - Many other values of  $\alpha$  and  $\beta$  also satisfy these equations.

# Dependence Testing

Do iteration vectors  $\alpha$  and  $\beta$  exist such that:

- (vectors)  $\alpha < \beta$
- $fX(\alpha) = gX(\beta)$ , for  $1 \leq X \leq M$

How can we find  $\alpha$  and  $\beta$  if they exist?



## Restrictions on Index functions

- $fX$  and  $gX$  must be decidable (i.e. computable)
- $fX$  and  $gX$  must be "analyzable"
  - to avoid brute force search

# Affine Index Functions

- Let  $fX$  and  $gX$  must be affine functions of loop indices:
  - i.e. for  $fX(i_1, i_2, i_3, \dots, i_n)$
  - $fX = a_1i_1 + a_2i_2 + \dots + a_ni_n + e$
  - $e$  is optional loop invariant calculation (i.e. constant for the loop)

# Dependence Testing on Restricted Index Functions

- Given that  $fX$  and  $gX$  are affine functions of loop indices
- Do iteration vectors  $\alpha$  and  $\beta$  exist such that:
  - (vectors)  $\alpha < \beta$
  - $fX(\alpha) = gX(\beta)$ , for  $1 \leq X \leq M$

How can we find  $\alpha$  and  $\beta$  if they exist?

What is this problem better known as? Hint: an affine function is a linear function plus constant.

# Dependence Testing

- Integer Linear Programming is NP-complete
- Lots of heuristics invented
  - Profitable to know if no solution exists since it implies no dependence!
  - See Chapter 3 of AK
  - Or Chapter 11 of the Dragon Book
  - We will cover this in a later class

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## Representing Dependences

Do we need to track all the iterations that have a dependence explicitly (e.g. in a list)?

# Distance Vectors

$$d(i, j)_k = j_k - i_k$$

- Where  $i, j, d(i, j)$  are  $n$ -element vectors
- $i_k$  indicates  $k$ -th element of  $i$

Example distance vector:  $(0, 1)$

# Direction Vectors

$$D(i, j)_k =$$

- " $<$ ", if  $d(i, j)_k > 0$
- " $=$ ", if  $d(i, j)_k = 0$
- " $>$ ", if  $d(i, j)_k < 0$

Example direction vector for  $(0, 1)$ :  $(=, <)$



## Information we need to track

For every pair of memory references:

- Iteration Vectors  $i$  and  $j$  which have a dependence, or
- Unique Distance Vectors  $d(i, j)$ , or
- Unique Direction Vectors  $D(i, j)$

- Which of these indicates a loop-independent dependence?
  - $(=, =)$
  - $(=, <)$
- Of the loop-carried dependence in example above, what level is the loop-carried dependence?

# Theorems

WARNING: Informal language

- Direction Vector Transform (Theorem 2.3 in AK)
  - If a transformation reorders loop iterations, and preserves the leftmost non-"=" component as "<", all dependences are preserved.
- Theorem 2.4 in AK
  - If a level- $k$  dependence exists, and a transformation reorders loop iterations while not reordering the level- $k$  loop
  - And does not move loops inside  $k$  outside the loop and vice versa
  - It preserves all level- $k$  dependences.
- Iteration Reordering (Theorem 2.6 in AK)
  - Iterations of a level  $k$  loop can be reordered if there is no level  $k$  dependence.

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## References

- Much of this lecture is based on Allen and Kennedy, Optimizing Compilers for Modern Architectures, Chapter 2.
- Chapter 11 of the Dragon Book also presents this information, but differently.