CSC2/455 Software Analysis and Improvement
Instruction Selection/Instruction Scheduling/Register Allocation

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URCS
Outline

Review

Instruction Selection

Instruction Scheduling

Global Register Allocation

Postscript
Outline

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Postscript
Optimizations

- Analysis
  - Iterative Dataflow Analysis
  - SSA Form
  - Interprocedural Analysis
  - Loop Analysis
  - Type Checking [next class]
- Optimization
  - Dead Code Elimination
  - Partial Redundancy Elimination
- Code Generation
  - Instruction Selection
  - Instruction Scheduling
  - Register Allocation
Outline

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Postscript
The Problem

\[ a = b + c \]
\[ c = c + 1 \]

gets converted to (in some pseudo-assembly language):

\[
    \text{add } a, b, c \\
    \text{inc } c 
\]
$x = x + 1 \Rightarrow \text{inc } x$

- A pattern (or "macro") is specified and when it matches on the 3-address code, instructions are emitted.
- Pattern matching is no longer used to select instructions
  - but still useful for something called *peephole optimization*
Tree Covering

- Pattern matching operates on sequences of instructions
  - Usually one, but could be more
- Tree-based systems operate on *trees* of expressions
  - Limited to basic blocks
See Figure 1 from the Twig Paper
See Table 1 from the Twig Paper
See Page 501 of the Twig Paper.

- Informally, a tree pattern can be viewed as a string formed by a pre-order traversal of the tree.
- All patterns can then be viewed as strings and an automaton constructed to search for strings and substrings.
See the discussion in Section 5 of the Twig Paper.

- Evaluate costs of subtrees
- Use that to evaluate cost of tree
- Use dynamic programming to pick the tree "covering" with the lowest cost.
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The need for instruction scheduling

• Code is in 3-address form
• Instructions have been selected
  • e.g. \( a / 4 \) becomes \( \text{shr} \ a, 2 \) (shift-right) or \( \text{sar} \ a, 2 \) (shift-arithmetic right)
• What (linear) order should we output instructions?
A modern CPU pipeline

Source: Intel 64 and IA-32 Architectures Optimization Manual
Pipeline Details Known to Compiler

- Number of functional units and their types
  - Which instructions can be issued together
- Latencies of operations
  - Note some operations can have variable latencies (e.g. memory operations)
- Forwarding latencies
  - Delays when sending values from one functional unit to another
  - Or from one instruction to another
- And other sundry architecture details
  - See the architecture manuals for processor of interest
Two metrics of interest

- **Throughput**: Instructions completed per cycle
  - Higher is better
- **Latency**: Total cycles for execution
  - Lower is better
Basic block to data dependence graph

\[ a = x + y \]
\[ b = v / w \]
\[ c = a \times b \]
\[ d = z \ll 2 \]
\[ e = c - d \]
## Schedule

Assume the following delays:

- Add, Shifts, Subtraction: 1 cycle
- Multiplication: 2 cycles
- Division: 3 cycles

<table>
<thead>
<tr>
<th>Cycle</th>
<th>ALU1</th>
<th>ALU2</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>$a = x + y$</td>
<td>$b = v / w$</td>
</tr>
<tr>
<td>1</td>
<td>$d = z &lt;&lt; 1$</td>
<td>(busy)</td>
</tr>
<tr>
<td>2</td>
<td>(idle)</td>
<td>(busy)</td>
</tr>
<tr>
<td>3</td>
<td>$c = a * b$</td>
<td>(idle)</td>
</tr>
<tr>
<td>4</td>
<td>(busy)</td>
<td>(idle)</td>
</tr>
<tr>
<td>5</td>
<td>$e = c - d$</td>
<td>(idle)</td>
</tr>
</tbody>
</table>

What determines the total time of the path?
Critical Path

- The path that takes the longest to execute
  - Equivalently, has zero slack
- Delaying operations on critical path will increase total time
Goal: Schedule instructions on every cycle (i.e. build a table like in previous figure)

- Overall structure:
  - Mark operations whose predecessors have completed as ready
  - Pick operations that are ready
  - Schedule them in some order if resources available
  - Proceed to next cycle
  - Repeat until all operations are scheduled

(Figure 12.3 in Cooper and Torczon)
Always prioritize instructions on critical path.

- Can be hard to achieve
- Variable delays
- Multiple functional units
- Other architecture-specific constraints
Other Complications

- Small basic block sizes
  - Loop unrolling
  - Trace scheduling
- Long dependence chains with little parallelism
  - Software Pipelining
The Memory Hierarchy (on CPUs)

- Registers (on chip, few tens to low hundreds)
- L1 cache (tens of KB)
- L2 cache (tens of MB)
- DRAM (tens of GB)
  - off-chip
Numbers every programmer should know

- Register access: 1 cycle
- L1 cache: less than 10 cycles
- L2 cache: less than 100 cycles
- DRAM: hundreds of cycles

Bottomline: Placing frequently used variables in registers can improve performance
The Problem of Global Register Allocation

- Code is in SSA form (with $\phi$-functions removed)
- SSA form uses infinite registers (each temporary is a “register”)
- Must map these logical registers to physical machine registers
Must answer two main questions:

- How many physical registers are needed? ($k$ is the number of physical registers)
  - $< k$, fewer than available – everybody gets a register!
  - $> k$, more than available – results in spill code
  - Spilled registers are stored in memory and reloaded later
- Which physical register is assigned to which variable?
A live range is a set of all definitions and uses of the same variable such that:

- If a use $u$ for variable $i$ is in $LR_i$, then all definitions $d$ that reach $u$ are also in $LR_i$.
- If a definition $d$ for variable $i$ is in $LR_i$, then all uses $u$ reached by $d$ are also in $LR_i$. 
Live range example

- Live ranges
  - \( LR_a = \{a_0\} \)
  - \( LR_b = \{b_0\} \)
  - \( LR_c = \{c_0\} \)
  - \( LR_d = \{d_0, d_1, d_2\} \)
Interference Graphs

- Two live ranges "interfere" if they are both live at an operation
  - Implies that they must be in different registers
- Represented by an interference graph
  - Nodes are live ranges
  - An (undirected) edge between nodes $n$ and $m$ indicates interference
Given an interference graph, the minimal number of physical registers required is equal to the *chromatic number* of the graph (due to Chaitin).

- Chromatic number is the minimum number of colors required such adjacent nodes have different colors.
- NP-complete problem
Three Color Example

![Graph (LR0, LR1, LR2, LR3)]

- LR0
- LR1
- LR2
- LR3

![Graph (LR0, LR1, LR2, LR3)]

- LR0
- LR1
- LR2
- LR3
Parting thoughts

What order should instruction selection, instruction scheduling, and register allocation be performed?
References

- Aho, Ganapathi, and Tjiang, Code Generation Using Tree Matching and Dynamic Programming, ACM TOPLAS
- Blindell, Survey on Instruction Selection: An Extensive and Modern Literature Review, arXiv
- Chapter 12 of Cooper and Torczon (Instruction Scheduling)
- Chapter 13 of Cooper and Torczon (Register Allocation)
- Also recommended: Chapter 8 of the Dragon Book