

Assemblers

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Assembler

Translation takes 2 steps:

1. associate memory locations with labels
2. translate each statement by combining the numeric representation of
 - the opcode,
 - the register specifiers, and
 - the label/immediate values

Why? - Forward References

The label is used before it is defined

```
        bne    $4, $8, L1
...
L1:     addu   $4, $4, 5
```

When the assembler sees **bne** the first time, it doesn't know **L1**'s location.

Two solutions — time versus space trade-off

1. 2 step translation
2. *backpatching* — maintain a table of forward references; update each reference (instruction) when the label is defined.

Step 1

Record the name and position of each label in the *symbol table*.

To determine the position, the assembler must determine how many words each instruction or data declaration occupies.

- position is relative from the start of text, data, or bss.
- variable length instructions or data declarations complicate the calculation.
 - SPARC/MIPS instructions are easy, but data declarations such as **.ascii** are harder.

Symbol Table

Often a hash table

Each entry contains

- the string representing the symbol
- the segment (and a bit for private versus global)
 - one of undefined, absolute, text, data, or bss
- the offset from the start of the segment

Step 2

Produce machine code using an *opcode table*.

The opcode table describes how to combine

- the opcode,
- the register specifiers, and
- the label/immediate values

into the representation for the instruction or data declaration.

References (to external symbols in another file) are unresolved.

The symbol table is appended to the object file for use by the linker and debugger.

Opcode Table

For a simple assembler

```
int size_align();
int size_ascii();
int size_1();
int size_4();

void output_align();
void output_ascii();
void output_byte();
void output_instr_3opnd();

struct Opcode opcodes[] = {
    { ".align", 0x00000000, size_align, output_align },
    { ".ascii", 0x00000000, size_ascii, output_ascii },
    { ".byte", 0x00000000, size_1, output_byte },
    { "save", 0x81E00000, size_4, output_instr_3opnd },
    { "sub", 0x80200000, size_4, output_instr_3opnd },
}
```

Object File Format

Six distinct sections

1. *object file header* – describes the size and position of the other pieces of the file
2. *text segment*
3. *data segment*
4. *relocation information* – identifies instructions and data words that depend on absolute addresses
5. *symbol table* – see `/usr/include/nlist.h`
6. *debugging information*

Object File Header

The first 32 bytes of every object file

```
/usr/include/sys/exechnr.h
/*
 * format of the exec header
 * known by kernel and by user programs
 */
struct exec {
    unsigned long a_magic; /* magic number */
    unsigned long a_text; /* size of text segment */
    unsigned long a_data; /* size of initialized data */
    unsigned long a_bss; /* size of uninitialized data */
    unsigned long a_syms; /* size of symbol table */
    unsigned long a_entry; /* entry point */
    unsigned long a_trsize; /* size of text relocation */
    unsigned long a_drsize; /* size of data relocation */
};
```

Linker/Separate Compilation

Search libraries to find undefined labels

Determine memory locations for each module and relocate instructions by adjusting absolute addresses

Resolve references among files - external reference to unresolved reference with the same name

