Principles of Network Security

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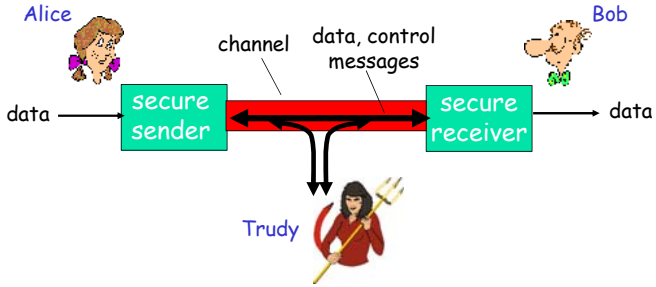
The Network Security Model

- Bob and Alice want to communicate "securely".
- Trudy (the adversary) has access to the channel.



Who might Bob and Alice be?

- Web browser/server for electronic transactions (e.g., on-line purchases/banking)
- DNS servers
- routers exchanging routing table updates
- … well, *real-life* Bobs and Alices!

What can an adversary do?

- eavesdrop: understand the content of messages
- actively changing messages
- impersonation: fake (spoof) identity
- denial of service: prevent service from being used by others (e.g., by overloading resources)

What is Network Security?

Confidentiality: only sender, intended receiver should "understand" message contents

Authentication: sender, receiver want to confirm identity of each other

Message Integrity: sender, receiver want to ensure message not altered (in transit, or afterwards)

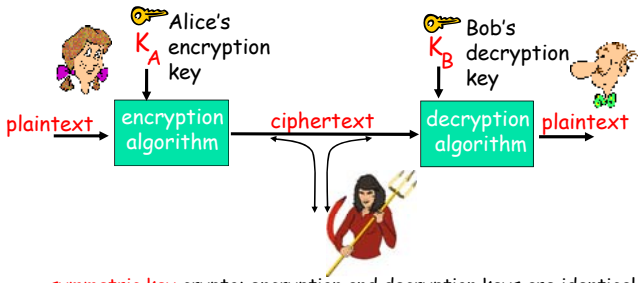
Access and Availability: services must be accessible and available to (and only to) legitimate users.

Principles of Network Security

- Confidentiality: cryptography
- Authentication
- Integrity

The Language of Cryptography

First goal of cryptography: confidentiality.



- symmetric key crypto: encryption and decryption keys are identical. (both are secret)
- public key crypto: encryption key is public, decryption key is secret.

Symmetric Key Cryptography: Monoalphabetic Cipher

Monoalphabetic cipher: substitute one letter for another.

```
plaintext:  abcdefghijklmnopqrstuvwxyz
            |                       |
ciphertext: mnbvcxzasdfghjklpoiuytrewq
```

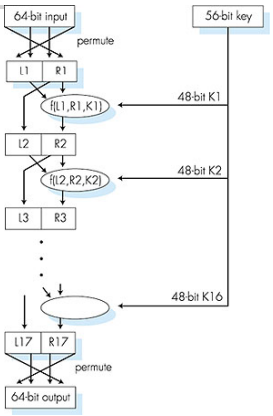
Example: `Plaintext: bob. i love you. alice`
`ciphertext: nkn. s gktc wky. mgsbc`

Q1: How hard to break this simple cipher?
- brute force?
- other?

Q2: How to make it more difficult to break?

Symmetric Key Cryptography: DES

- DES: Data Encryption Standard
  - US encryption standard [NIST 1993]
  - 56-bit symmetric key, 64-bit plaintext input
  - encryption: initial permutation ⇒ 16 "rounds", each using different 48 bits of key ⇒ final permutation
  - decryption: reverse operation using the same key
- How secure is DES?
  - DES Challenge (1999): 56-bit-key-encrypted phrase decrypted (brute force) in 22 hours 15 minutes
- Making DES more secure:
  - use three keys sequentially (3-DES)
  - use more bits



AES: Advanced Encryption Standard

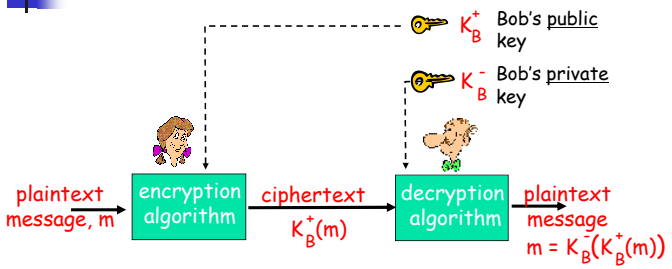
- new (Nov. 2001) symmetric-key NIST standard, replacing DES
- processes data in 128 bit blocks
- 128, 192, or 256 bit keys
- brute force decryption (try each key) taking 1 sec on DES, takes 149 trillion years for 128-bit AES

Public Key Cryptography

symmetric key cryptography
- requires sender, receiver know shared secret key
- Q: how to agree on key in first place? (particularly difficult if Trudy is eavesdropping on all communication)

public key cryptography
- encryption key is different from decryption key
- encryption key is public, known to everyone, also called public key
- decryption key is secret, known only to receiver, also called private key

Public Key Cryptography



Principle for choosing the public/private key pair:
One should not be able to derive the private key from the public key.

Public Key Cryptography: RSA
(Ron Rivest, Adi Shamir and Len Adleman)

- Choosing keys:
  - Choose two large prime numbers p, q. (e.g., 1024 bits each)
  - Compute n = pq, z = (p-1)(q-1)
  - Choose e (with e<n) that has no common factors with z.
  - Choose d such that ed-1 is exactly divisible by z.
  - Public key is (n,e). Private key is (n,d).
- To encrypt a message, m (<n): do $c = m^e \bmod n$
- To decrypt a received ciphertext, c: do $m = c^d \bmod n$
- Reason: for any m (relatively prime with n)
  - $m^z \bmod n = 1$; therefore $m^{ed-1} \bmod n = 1$
- Another property: $(m^d \bmod n)^e \bmod n = m$
- RSA is much slower than the symmetric key cryptos.

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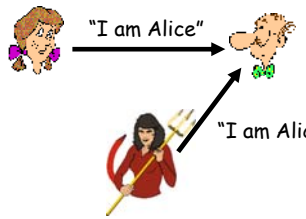
Authentication: version 1.0

Authentication: Bob wants Alice to "prove" her identity to him.

Protocol ap1.0: Alice says "I am Alice".

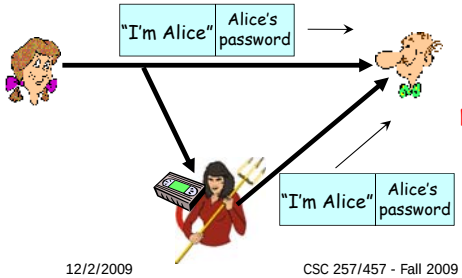


Failure scenario??

Trudy can simply declare herself to be Alice

Authentication: version 2.0

Protocol ap2.0: Alice says "I am Alice" and sends her secret password to "prove" it.



Failure scenario??

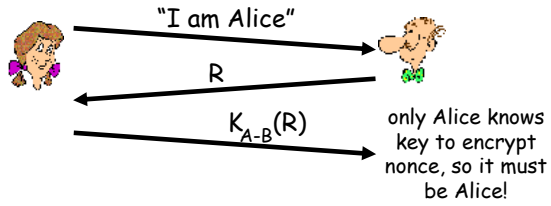
playback attack: Trudy records Alice's packet and later plays it back to Bob

Authentication: version 3.0

Goal: avoid playback attack

Nonce: number (R) used only *once-in-a-lifetime*

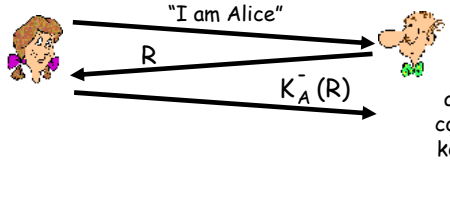
ap3.0: Bob sends Alice a nonce, R. Alice must return R, encrypted with shared secret key



Authentication: version 4.0

ap3.0 requires shared symmetric key. Key distribution can be a problem.

ap4.0: use nonce, public key cryptography.



Principles of Network Security

- Confidentiality: cryptography
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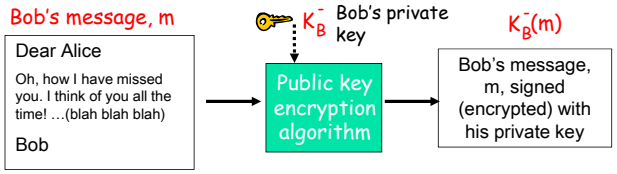
Integrity

- Digital Signatures:
  - Cryptographic technique to ensure document integrity.
  - analogous to hand-written signatures.
- sender (Bob) digitally signs document, establishing he is document owner/creator.
- the recipient (Alice) receives the document and the digital signatures.
- the recipient can be sure that the document is
  - verifiable: Bob signed the document.
  - nonforgeable: the document hasn't been changed since Bob signed it.

Digital Signatures

- Bob signs m by encrypting with his private key, creating a digital signature $K_B^-(m)$

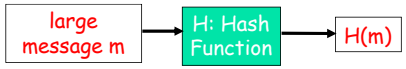


- Suppose Alice receives msg m and its digital signature $K_B^-(m)$
- Alice applies Bob's public key K_B^+ to $K_B^-(m)$ then checks whether $K_B^+(K_B^-(m)) = m$.
- If so, whoever signed m must have used Bob's private key.

Problem: computationally expensive to public-key-encrypt long messages.

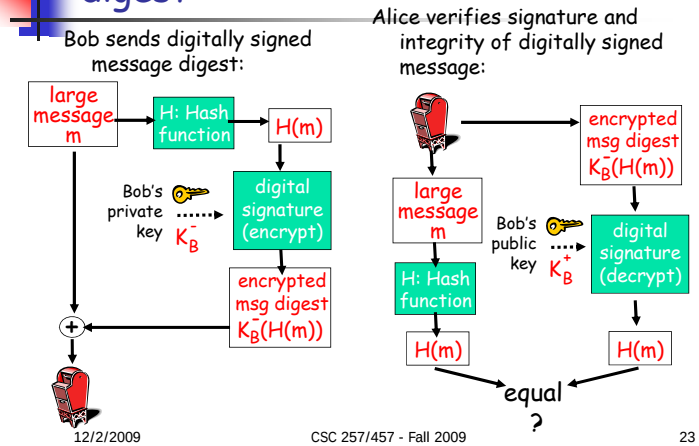
Message Digests

- apply a hash function H to m, get a much smaller message digest H(m).



- public-key-encrypt the message digest to generate the digital signature $K_B^-(H(m))$.
- example hash functions?

Digital signature = signed message digest



Message Digests: good/bad hash function

- apply a hash function H to m, get a much smaller message digest H(m).
- public-key-encrypt the message digest to generate the digital signature $K_B^-(H(m))$.

Good/bad has functions?

- Note: given a hash function, it is possible for many messages sharing the same digest.

Internet Checksum: Poor Hash Function for Generating Message Digests

Given a message and its Internet checksum, it is easy to find another message with same checksum.

message	ASCII format
I O U 1	49 4F 55 31
0 0 . 9	30 30 2E 39
9 B O B	39 42 D2 42
	B2 C1 D2 AC

message	ASCII format
I O U 9	49 4F 55 39
0 0 . 1	30 30 2E 31
9 B O B	39 42 D2 42
	B2 C1 D2 AC

different messages but identical checksums!

Hash function property: given digest x for message m, computationally infeasible to find another message m' that shares the same digest.

Good Hash Functions for Generating Message Digests

- MD5 hash function widely used
  - computes 128-bit message digest in 4-step process.
  - appears difficult to construct message m whose MD5 hash is equal to x.
- SHA-1 is also used.
  - [NIST, FIPS PUB 180-1]
  - 160-bit message digest

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