Parallel Programming

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Parallel Programming Steps

- Converting a sequential application to a parallel one
  - Decomposition into tasks
  - Assign tasks to processors
  - Orchestrate data access and synchronization

Dependences

- Two tasks are dependent if they must follow their order in the sequential program so that the program execution isn’t changed.
- Three types of data dependences:
  - Read-after-write
  - Write-after-read
  - Write-after-write
- Only “read-after-write” is called true dependence

Task Decomposition

- Decomposition:
  - over different functions
  - over different data segments (over loop iterations)
- How good is the decomposition?
  - Assume dependent tasks must run serially. We can build a DAG of task dependencies and critical path length indicates lower bound of parallel execution time.
- Tradeoff on task granularity:
  - smaller tasks may offer more parallelism/concurrency
  - smaller tasks require more management/programming overhead
Task Assignment

Goals:
- Load balance
- Minimize inter-processor communication/synchronization
  ⇒ maximize locality

Ways:
- Static assignment
- Dynamic assignment

Dynamic Task Assignment

- How does it work?
  - Maintain a centralized queue of ready tasks, protected by mutex lock
  - Each thread grabs a task at the beginning; grabs another task after completing the current one
  - New tasks may be generated on the fly and added to queue

- Advantage: good load balancing

- Disadvantages with dynamic task assignment
  - Locality may be lost in the interests of load balancing
  - Synchronization contention on manipulating the task queue

- Fixable through distributed queues with work stealing
  - Each thread has an exclusive task queue with good locality and no contention.
  - When out of work (load imbalance or synchronization), steal some task from another queue with ready tasks

Orchestration

- Access shared data
  - Shared-memory parallel platform
  - Distributed-memory parallel platform

- Synchronization
  - mutex lock, condition variables, barrier, ...
  - Enforce dependences

- Fine-grained synchronization between tasks
  - Require tasks to run simultaneously
  - Work better with static task assignment
Parallel Programming Example: Successive Over Relaxation

- SOR implements a mathematical model for many natural phenomena, e.g., heat dissipation, ocean currents.
- Given a 2D grid of data, for some number of iterations:
  - For each internal grid point, compute average of its four neighbors.

```c
for (i=1; i<n; i++)
    for (j=1; j<n; j++)
        temp[i][j] = 0.25 * (grid[i-1][j]+grid[i+1][j]+grid[i][j-1]+grid[i][j+1]);
for (i=1; i<n; i++)
    for (j=1; j<n; j++)
        grid[i][j] = temp[i][j];
```

- Task decomposition:
  - 1D partitioning: each task manages some columns or rows
  - 2D partitioning: each task manages a 2D block of the grid
- Impact on communications in distributed memory parallel computing?

Parallel Programming Example: Successive Over Relaxation

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```

- Dependences:
  - First (i,j) loop nest?
  - Second (i,j) loop nest?
  - Between the two loop nests?
  - Between two iterations?

Parallel Programming Example: Gaussian Elimination

- Solving a system of linear equations
- Reduce an equation matrix into an equivalent upper-diagonal matrix
- Partial pivoting to maintain numerical stability

```
\[
p* A_1 + A_2 \quad X \quad R_1 \quad r_1 + p*r_1
\]```
Assignment #1
Shared Memory Parallel Programming
- All should have accounts in CS research network
- Familiarize with the parallel machines

Pre-assignment
- We provide you sequential/parallel SOR code (a different version, called red-black ordering)
- You read, understand, and run it

Main assignment
- Parallel Gaussian Elimination
- Different settings (input matrices, machines, proc #’s)
- Analysis and comparison
- Grading based on your written report!

Pthreads Synchronization Primitives
- Mutex lock (mutual exclusion)
  - pthread_mutex_lock
  - pthread_mutex_unlock
- Condition variable (waiting for a condition)
  - pthread_cond_wait
  - pthread_cond_signal
  - pthread_cond_broadcast

Barrier Synchronization
- A barrier is set for all threads
- A thread can proceed beyond its barrier if and only if all threads have reached respective barrier points
- Implement the barrier in pthreads?
  - Count the number of arrivals at the barrier
  - When a thread arrives, wait if this is not the last one, otherwise unlock everyone else and proceed
  - Mutex lock to protect the arrival number, condition variable to implement wait and broadcast