## CSC 284/484 Advanced Algorithms - homework 1 applied homework due: February 6th, 11:59pm EST theoretical homework due: February 6th, 12:30pm EST

Do not search for a solution online, do not use any written material when writing any part of the code (for example, no copy-paste, no open textbook when writing code, no reediting of an old source file from an old project, etc). You can discuss your solution(s) or problems/issues you are facing with your classmates. Do not share code.

- Sufficient: Implement any *s*-*t*-max-flow algorithm.
  - Test your algorithm on the attached test data.
  - Write a little note indicating:
    - \* Whether your code passed all the tests (if not where is the problem).
    - \* Timing information (how long your code runs on each test).
- Better: Implement two *s*-*t*-max-flow algorithms.
  - Test your algorithms on the attached test data.
  - Write a little note indicating:
    - \* Whether your code passed all the tests (if not where is the problem).
    - \* Timing information (how long your code runs on each test).
    - \* Compare the two algorithms.
- Even Better: Try all kinds of heuristics for the algorithms... see which one is the best on the test data.
- Your note should be typeset (LaTeX, Word, etc); send your code and note to the instructor and the TA (breber@cs.rochester.edu) (subject line: CSC 284/484 HOMEWORK 1).
- Allowed languages: C, C++, Java, Python (I discourage you from using Python—it is better for you to learn how to code algorithms in more efficient languages).

INPUT FORMAT: The first line contains k, the number of problems. Then descriptions of the problems follow. The first line contains n (the number of vertices) and m (the number of edges). The second line contains  $s, t \in \{1, ..., n\}$  (the source and the sink). Then m lines follow. Each line contains three integers u, v, c, where  $u, v \in \{1, ..., n\}$  and  $1 \le c \le 1,000,000$ ; this means that there is edge (u, v) with capacity c in the graph. We have  $2 \le n \le 1,000$  and  $1 \le m \le 10,000$ .

OUTPUT FORMAT: The output contains one line for each problem—the value of the maximum s-t flow.

EXAMPLE INPUT:

```
2
2 1
1 2 1
2 1
2 1
2 1
2 1
2 1
2 1
2 10
EXAMPLE OUTPUT:
10
0
```

**Definition 0.1** Let G = (V, E) be a directed graph. Let  $s, t \in V$  be two distinct vertices. An *s*-*t* alternative-cut is a pair of disjoint sets A, B such that  $s \in A, t \in B, A \cup B = V$ , every vertex in A is reachable from s using paths contained in A, and every vertex in B can reach t using paths contained in B.

Prove or disprove

**Conjecture 1** Let G = (V, E) be a directed graph with edge capacities  $c : E \to \mathbb{R}_{\geq 0}$ . Let  $s, t \in V$  be two distinct vertices. Assume that every  $v \in V$  is reachable from s. Assume that t can be reached from every  $v \in V$ . We have

$$\max_{s-t \text{ flows } f} |f| = \min_{s-t \text{ alternative-cuts } A, B} c(A, B).$$

Typeset your solution (LaTeX, Word, etc) and send pdf to the instructor by email (subject line: CSC 284/484 HOMEWORK 1). In addition to sending the pdf print your solution and bring it to class of Feb. 6.

We are given a directed graph G = (V, E), two vertices s and t and a positive integer K. We want to find s-t-walks  $P_1, \ldots, P_K$  such that no two walks have the same i-th edge (for any i) and the length of the longest walk is minimized.

(We describe the problem more formally. Let walk  $P_j$  be  $v_{j,0}, \ldots, v_{j,\ell_j}$  where  $v_{j,0} = s$  and  $v_{j,\ell_j} = t$ . We want to minimize  $\max_{j \in [K]} \ell_j$ . The condition on the walks is: for any distinct  $j, h \in [K]$  and any  $i \in [\min\{\ell_j, \ell_h\}]$  we have  $(v_{j,i-1}, v_{j,i}) \neq (v_{h,i-1}, v_{h,i})$ .)

Show how to solve the problem using flows.

Typeset your solution (LaTeX, Word, etc) and send pdf to the instructor by email (subject line: CSC 284/484 HOMEWORK 1). In addition to sending the pdf print your solution and bring it to class of Feb. 6.